




BRICS Research Activities

Algorithms

Gerth Stølting Brodal

Outline of Talk

- The Algorithms Group
- Courses
- Algorithm Events
- ALCOM  FT
- Expertise within BRICS

Examples

- Algorithms for Evolutionary Trees
- Cache Oblivious Algorithms

The Algorithms Group



Erik Meineche Schmidt
Algorithms, Complexity Theory



Sven Skyum
Algorithms, Complexity Theory



Peter Bro Miltersen
Complexity Theory, Data Structures



Gudmund Skovbjerg Frandsen
Algebraic Algorithms, Dynamic Algorithms



Christian Nørgaard Storm Pedersen
Bioinformatics, String Algorithms



Gerth Stølting Brodal
Data Structures, External Memory



Rolf Fagerberg
Data structures, External Memory

PhD students



Alex Rune Berg
Graph Theory



Jesper Makholm Byskov
Algorithms for NP-hard problems



Bolette Ammitzbøll Madsen
Algorithms for NP-hard problems



Bjarke Skjernaa
Algorithms for NP-hard problems

Kristoffer Arnsfelt Hansen
Complexity

Courses

- dProg (Introduction to Programming) *Frandsen*
- dADS (Algorithms and Data Structures) *Brodal, Schmidt*
- dSøgOpt (Searching and Optimization) *Skyum, Miltersen*
- Computer Architecture and Operating Systems *Pedersen*
- Computational Geometry *Skyum*
- Complexity *Miltersen*
- Randomized, Parallel and Dynamic Algorithms *Frandsen*
- Algorithms (ph.d. course) *Brodal, Fagerberg*
- External Memory Algorithms and Data Structures —
- Algorithms for Web Indexing and Searching —
- Algorithms in Bioinformatics *Pedersen*
- Genome Analysis —

Algorithm Events

Upcoming

18th IEEE Conference on Computational Complexity
(June 2003, organizing chair Peter Bro Miltersen)

Ongoing

Alcom seminar

Recent

EEF Summer School on Massive Data Sets (June 2002)

ALGO 2001 (August 2001)

- ESA 2001 - 9th Annual European Symposium on Algorithms
- WAE 2001 - 5th Workshop on Algorithm Engineering
- WABI 2001 - 1st Workshop on Algorithms in Bioinformatics



Algorithms and Complexity – Future Technologies

The ALCOM-FT project is a joint effort between eleven of the leading groups in algorithms research in Europe. The aim of the project is to discover **new algorithmic concepts**, identify **key algorithmic problems** in important applications, and contribute to the accelerated **transfer of advanced algorithmic techniques** into commercial systems.

The project takes place from June 2000 to June 2003.

- ALCOM-FT (continuation of ALCOM, ALCOM-II, ALCOM-IT)
- BRICS is the coordinator of ALCOM-FT (Erik Meineche Schmidt, Rolf Fagerberg)
- 336 technical reports since July 2000

BRICS

Barcelona

Cologne

INRIA Rocquencourt

Max-Planck-Institut für Informatik

Paderborn

Patras

Rome “La Sapienza”

Utrecht

Warwick

Cyprus

Erik Meineche Schmidt

Josep Díaz

Michael Jünger

Philippe Flajolet

Kurt Mehlhorn

Burkhard Monien

Friedhelm Meyer auf der Heide

Paul Spirakis

Giorgio Ausiello

Jan van Leeuwen

Mike Paterson

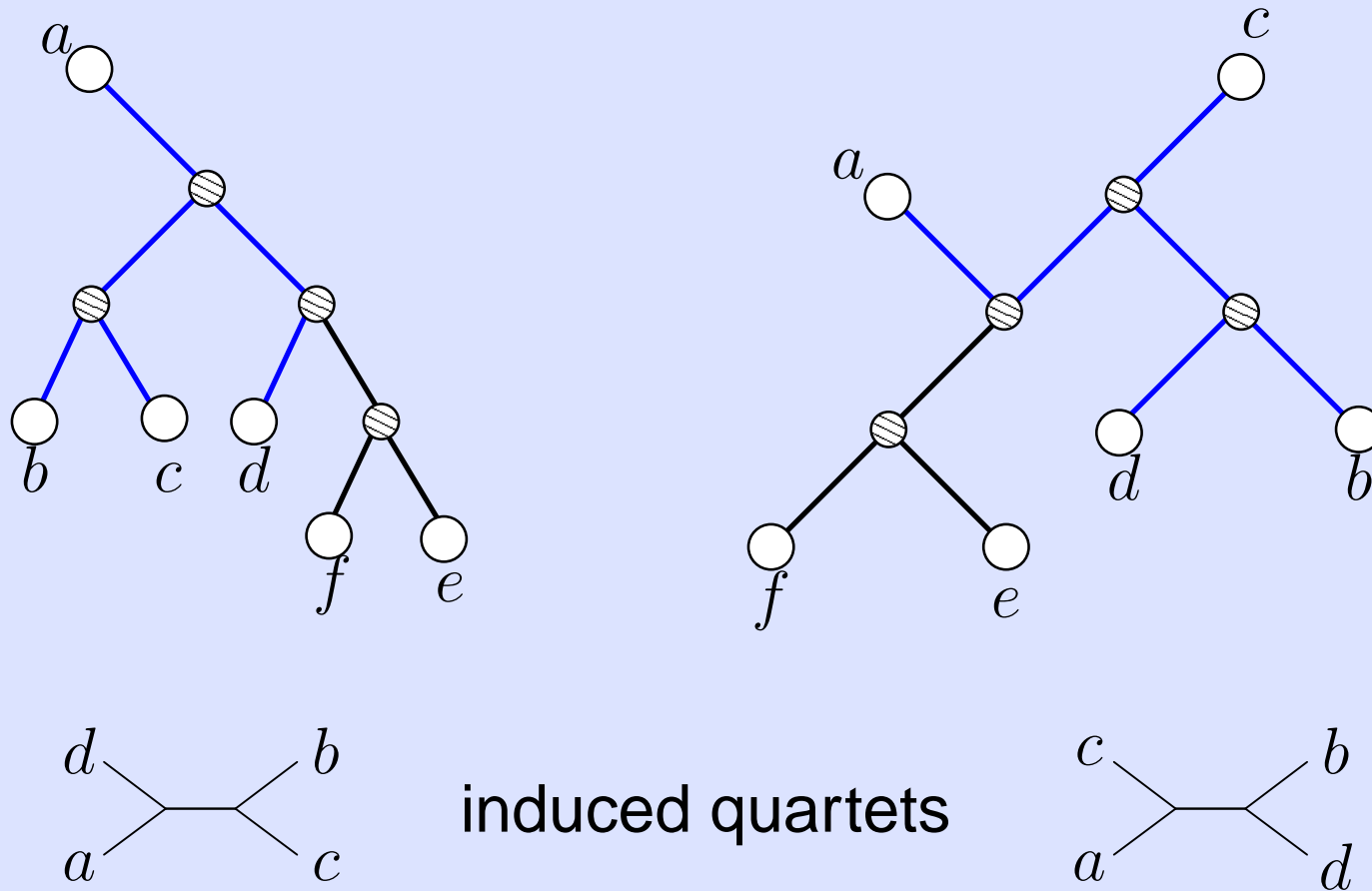
Marios Mavronicolas

Algorithm Expertise within BRICS

- Algorithms in general
- Complexity
- Data structures
- Dynamic algorithms
- External memory algorithms
- Algorithms in Bioinformatics
- Algorithm engineering / experimental algorithmics

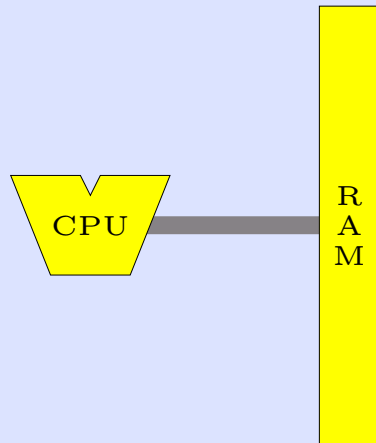
Examples

Algorithms for Evolutionary Trees



- Reconstruct evolutionary trees, e.g. from quartets
- Compare evolutionary trees, e.g. number of common quartets

The Classic RAM Model

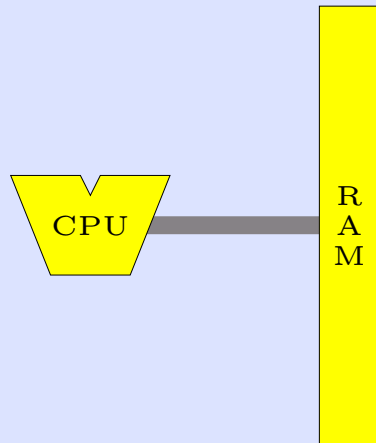


Add: $O(1)$

Mult: $O(1)$

Mem access: $O(1)$

The Classic RAM Model

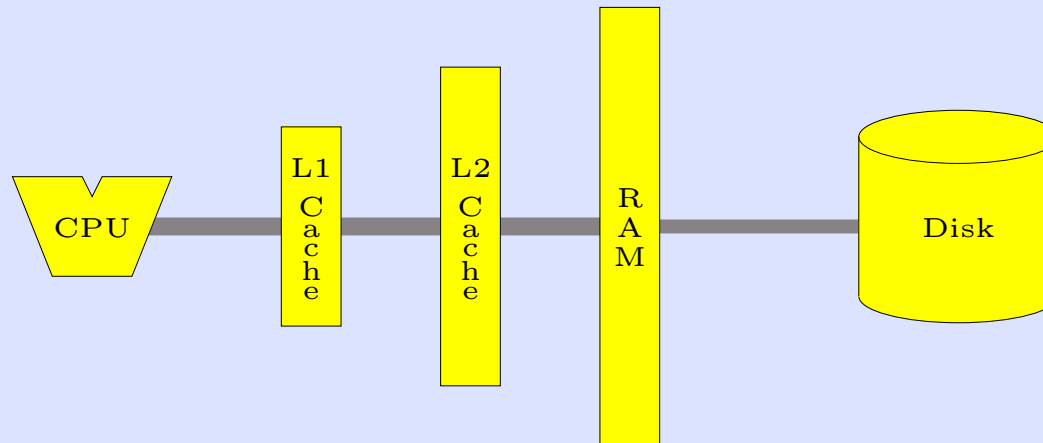


Add: $O(1)$

Mult: $O(1)$

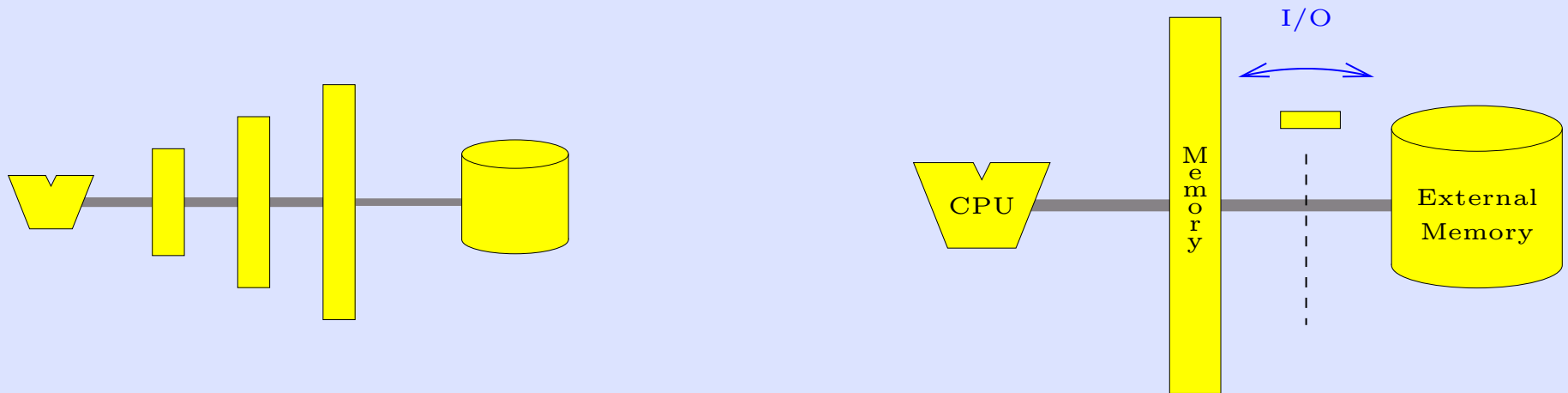
~~Mem access: $O(1)$~~

Real life



Bottleneck: transfer between two highest memory levels in use

The I/O Model



N = problem size

M = memory size

B = I/O block size

Aggarwal and Vitter 1988

- One I/O moves B consecutive records from/to disk
- **Cost:** number of I/Os
- E.g., sorting requires $O\left(\frac{N}{B} \log_{M/B} \frac{N}{B}\right)$ I/Os.

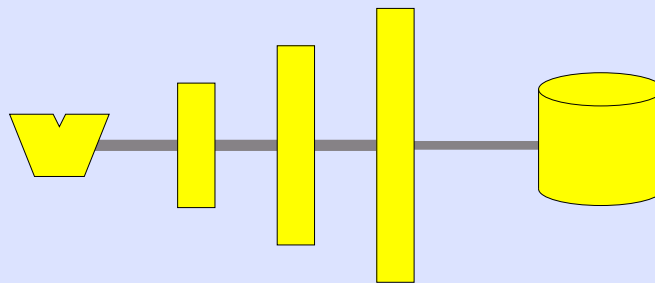
Cache-Oblivious Model

Frigo, Leiserson, Prokop, Ramachandran, FOCS'99

- Program in the RAM model
- Analyze in the I/O model (for arbitrary B and M).

Advantages

- Optimal on arbitrary level \Rightarrow optimal on **all levels**.
- B and M not hard-wired into algorithm.

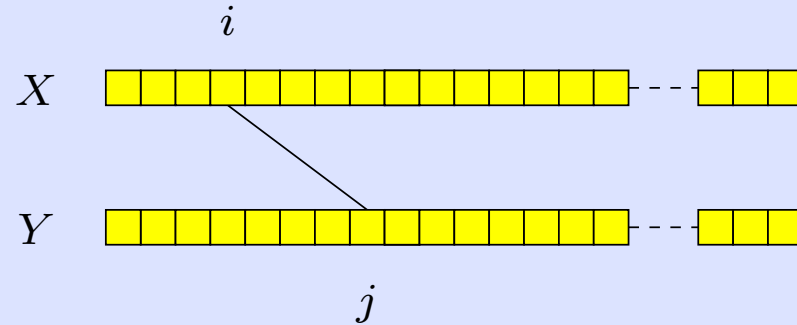


Cache-Oblivious Results

- Scanning \Rightarrow stack, queue, median finding,
- Sorting, matrix multiplication, FFT
Frigo, Leiserson, Prokop, Ramachandran, FOCS'99
- Cache oblivious search trees
Prokop 99
Bender, Demaine, Farach-Colton, FOCS'00
Rahman, Cole, Raman, WAE'01
Bender, Duan, Iacono, Wu and Brodal, Fagerberg, Jacob, SODA'02
- Priority queue and graph algorithms
Arge, Bender, Demaine, Holland-Minkley, Munro, STOC'02
Brodal, Fagerberg, ISAAC'02
- Computational geometry
Bender, Cole, Raman, ICALP'02
Brodal, Fagerberg, ICALP'02
- Scanning dynamic sets
Bender, Cole, Demaine, Farach-Colton, ESA'02

Double for-loop

X, Y arrays of length n :



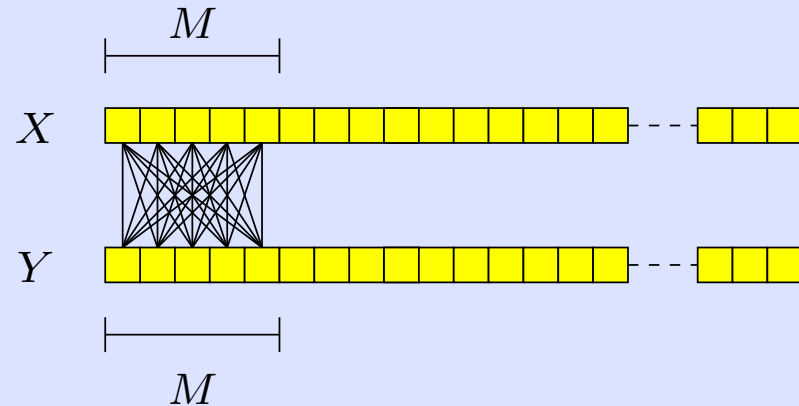
```
for(i=0; i<n; i++)  
  for(j=0; j<n; j++)  
    f(X[i], Y[j])
```

I/O complexity:

$$n \times \frac{n}{B} = \frac{n^2}{B}$$

Double for-loop

More I/O-efficient version:



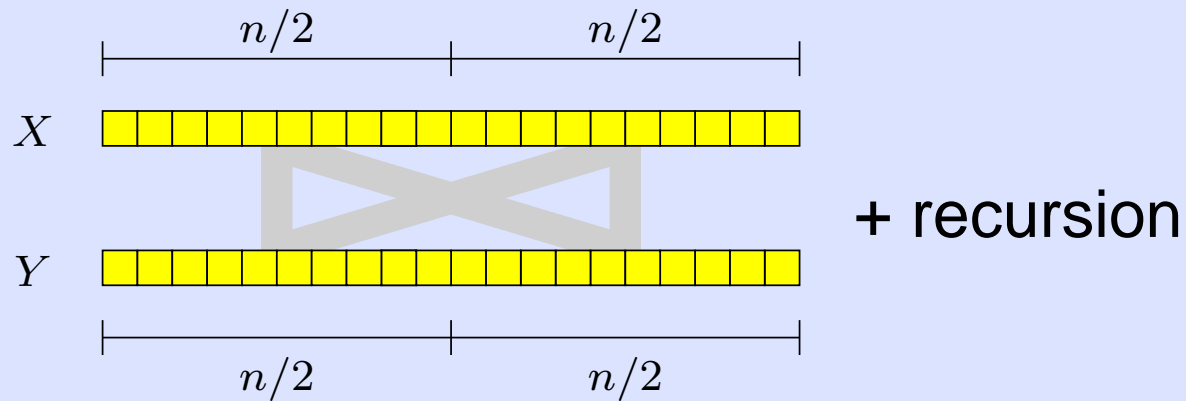
I/O complexity:

$$\frac{n}{M} \times \frac{n}{M} \times \frac{M}{B} = \frac{n^2}{MB}$$

```
for(i=0; i<n; i=i+M)
  for(j=0; j<n; j++)
    for(k=i; k<i+M; k++)
      f(X[k], Y[j])
```

Double for-loop

Cache-oblivious version:



I/O complexity:

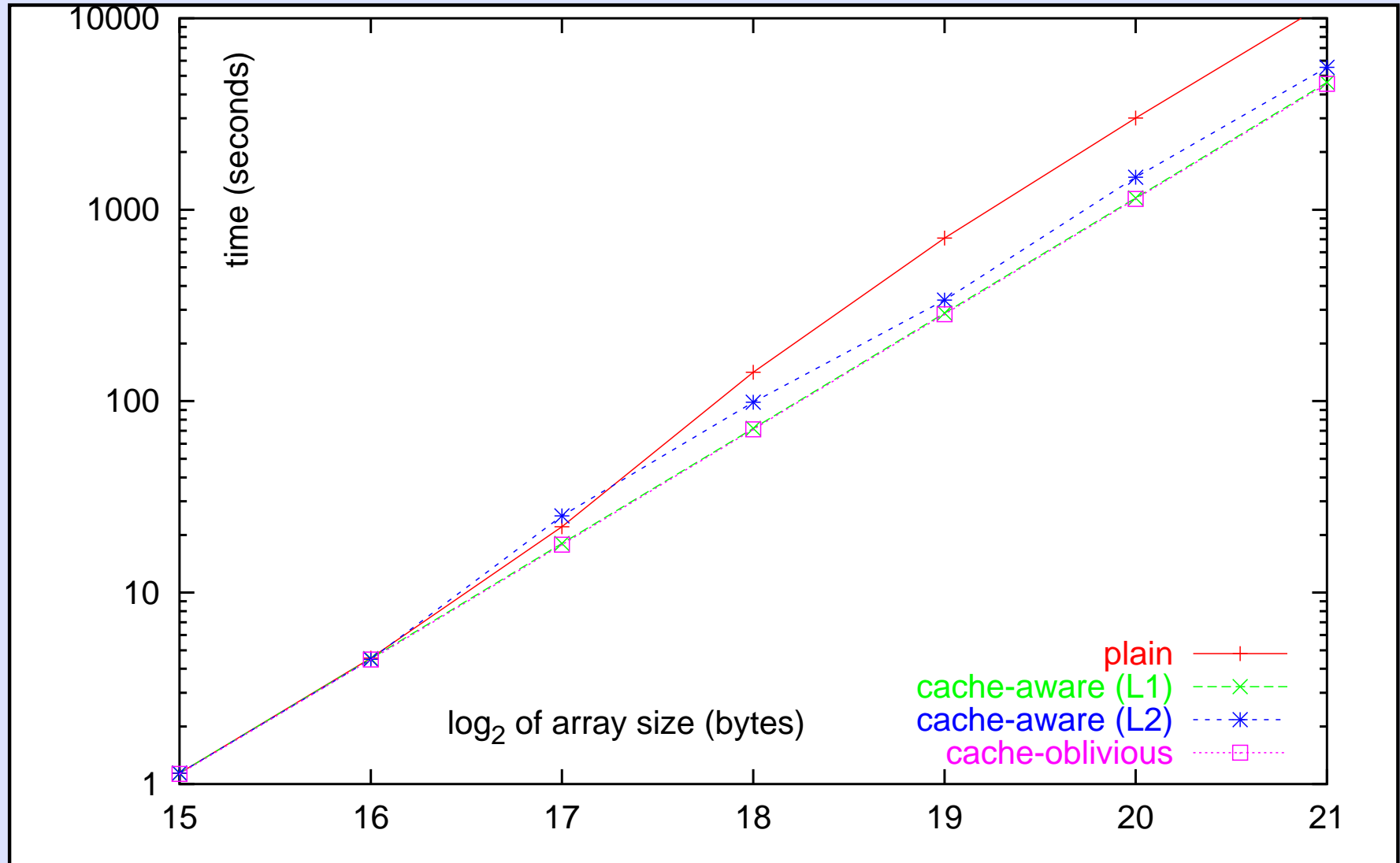
Again $\frac{n^2}{MB}$

Double for-loop

Cache-oblivious version

```
doubleLoop(i, j, length)
  if (length == 1)
    f(X[i], Y[j])
  else
    doubleLoop(i, j, length/2)
    doubleLoop(i, j+length/2, length/2)
    doubleLoop(i+length/2, j, length/2)
    doubleLoop(i+length/2, j+length/2, length/2)
```

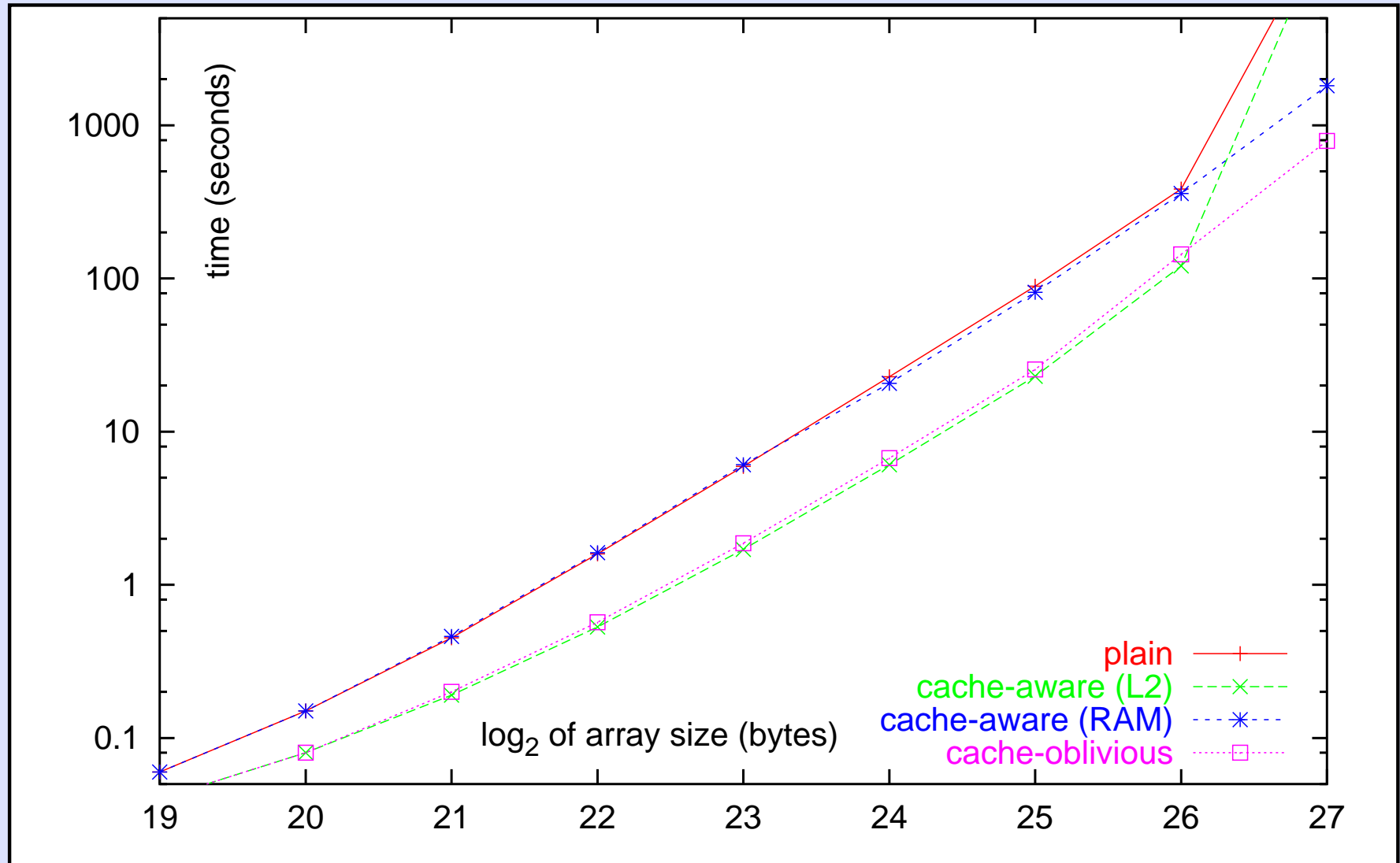
Double for-loop – Experiments



Sizes within RAM (element size 4 bytes)

366 MHz Pentium II, 128 MB RAM, 256 KB Cache, gcc -O3, Linux

Double for-loop – Experiments



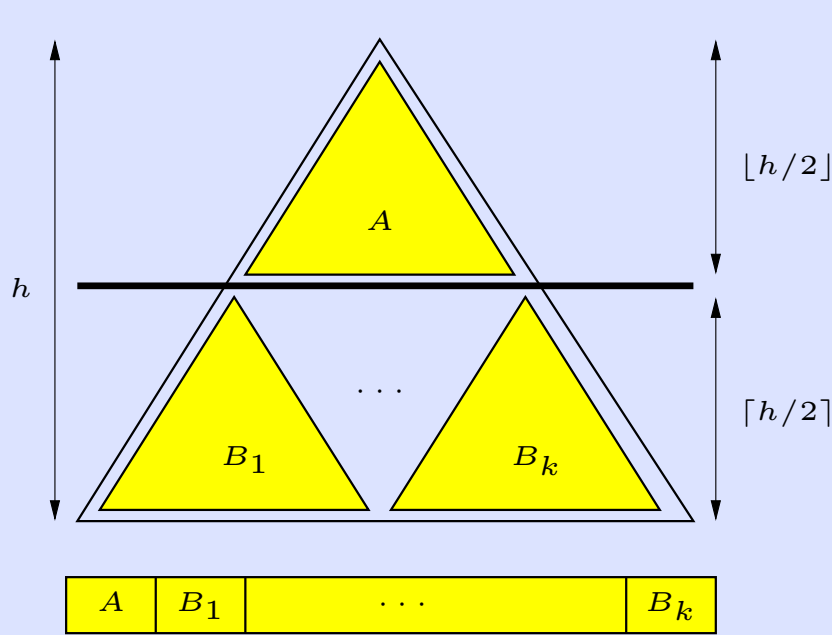
Sizes exceeding RAM (element size 1 KB)

366 MHz Pentium II, 128 MB RAM, 256 KB Cache, gcc -O3, Linux

Cache-Oblivious Search Trees

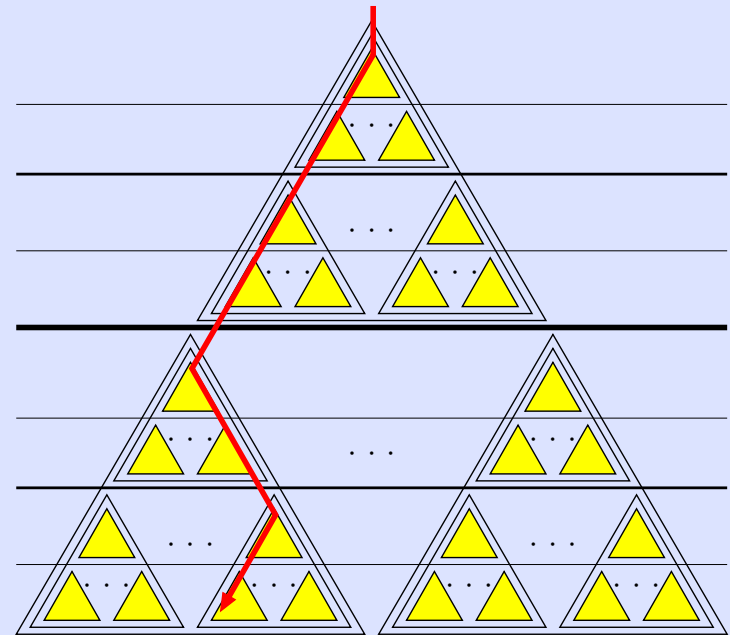
Recursive memory layout (van Emde Boas layout)

Prokop 1999



Binary tree

Dynamization



Searches use $O(\log_B N)$ I/Os

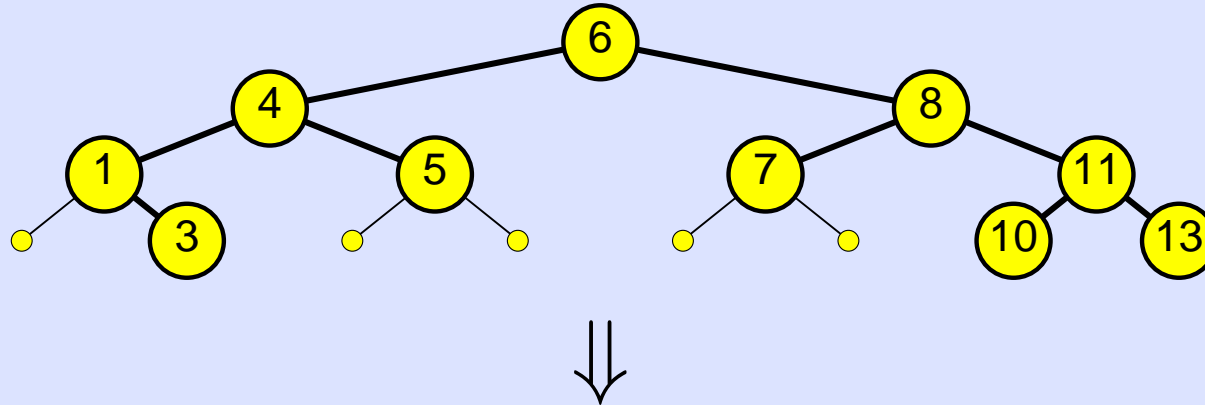
Bender, Demaine, Farach-Colton, FOCS'00

Rahman, Cole, Raman, WAE'01

Bender, Duan, Iacono, Wu, SODA 02

Brodal, Fagerberg, Jacob, SODA'02

Cache-Oblivious Search Trees



6	4	8	1	—	3	5	—	—	7	—	—	11	10	13
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Search	$O(\log_B N)$
Range Reporting	$O\left(\log_B N + \frac{k}{B}\right)$
Updates	$O\left(\log_B N + \frac{\log^2 N}{B}\right)$

Current Work...

$$M \geq B^2 ?$$

ID 5402

mit $\log(B!)$ N/B $B = \frac{N}{\log M}$
 Given N, M, B_1, B_2 $B_1 < B_2$
 Hint $O(\text{Sort}_{M, B_1}(N))$ IOs for B_1
 $\Omega(\text{Sort}_{M, B_2}(N))$ IOs for B_2

$$\frac{N}{B} \cdot B \cdot \log B = N \cdot \log B$$

$$t \cdot \log(k) = t \cdot k \cdot \log \frac{B}{k}$$

Heap sort $\log \log \log$
 Merge $\log \log$
 regular/ABSZIF $T \cdot k \cdot \log B$
 $\sum_{k=1}^N \frac{1}{k} \log \left(\frac{N+1-k}{k} \right) = \sum_{k=1}^N \frac{1}{k} \log \frac{N+1}{k} = O\left(\frac{1}{B} \cdot \log \frac{N}{B} \cdot k\right) = O\left(\frac{1}{B} \cdot \log \frac{N}{B} \cdot k\right)$
 $t \cdot k \cdot \log \frac{N}{k} \geq k \cdot \log \frac{N}{k}$
 $\Theta = \frac{N}{B} \cdot \log \frac{N}{B}$

Bounded Range Queries (BRQ)
 Diagram
 Triangulation
 Endpoint-Dense
 Interval decomposition
 location Monotone Sub
 Counting 2D

$\text{sort}(N) = N \cdot M \cdot B$
 $N = B^2 \quad M = 2B \quad B = B$ - Parallel
 $O(N) \quad N \quad M \quad B = 1$ - Parallel
 Input: $1 \ 2 \ B_1 \ B_2 \ 1 \ B$

$\sum_{k=1}^N \frac{1}{k} \log \frac{N+1-k}{k} = \sum_{k=1}^N \frac{1}{k} \log \frac{N+1}{k} = O\left(\frac{1}{B} \cdot \log \frac{N}{B} \cdot k\right)$
 $a + b^2 \geq ab = (a+b)^2$
 $\Phi(x) = \dots$