

# Sorting Integers in the RAM Model

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Joint work with Djamel Belazzougui and Jesper Sindal Nielsen  
(to be presented at the 14th Scandinavian Workshop on Algorithm Theory)

# Sorting (Integers)

Input **58** **13** **138** **55** **141** **137** **11** **53** **10** **54** *n integers*

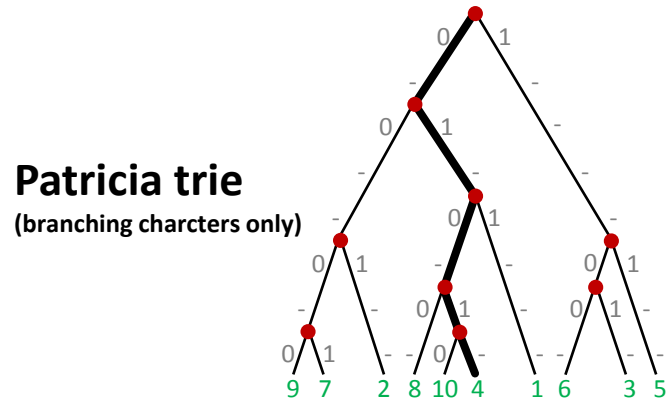
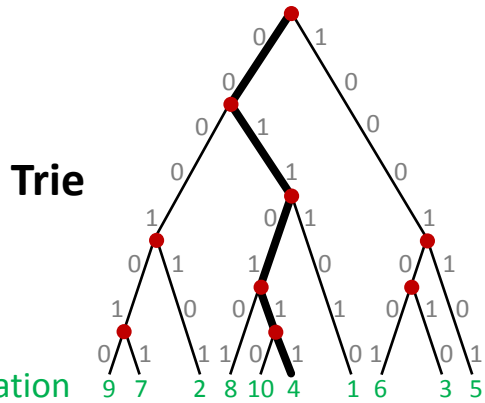
00111010 00001101 10001010 00110111 10001101 10001001 00001010 00110101 00001010 00110110

*w bits*



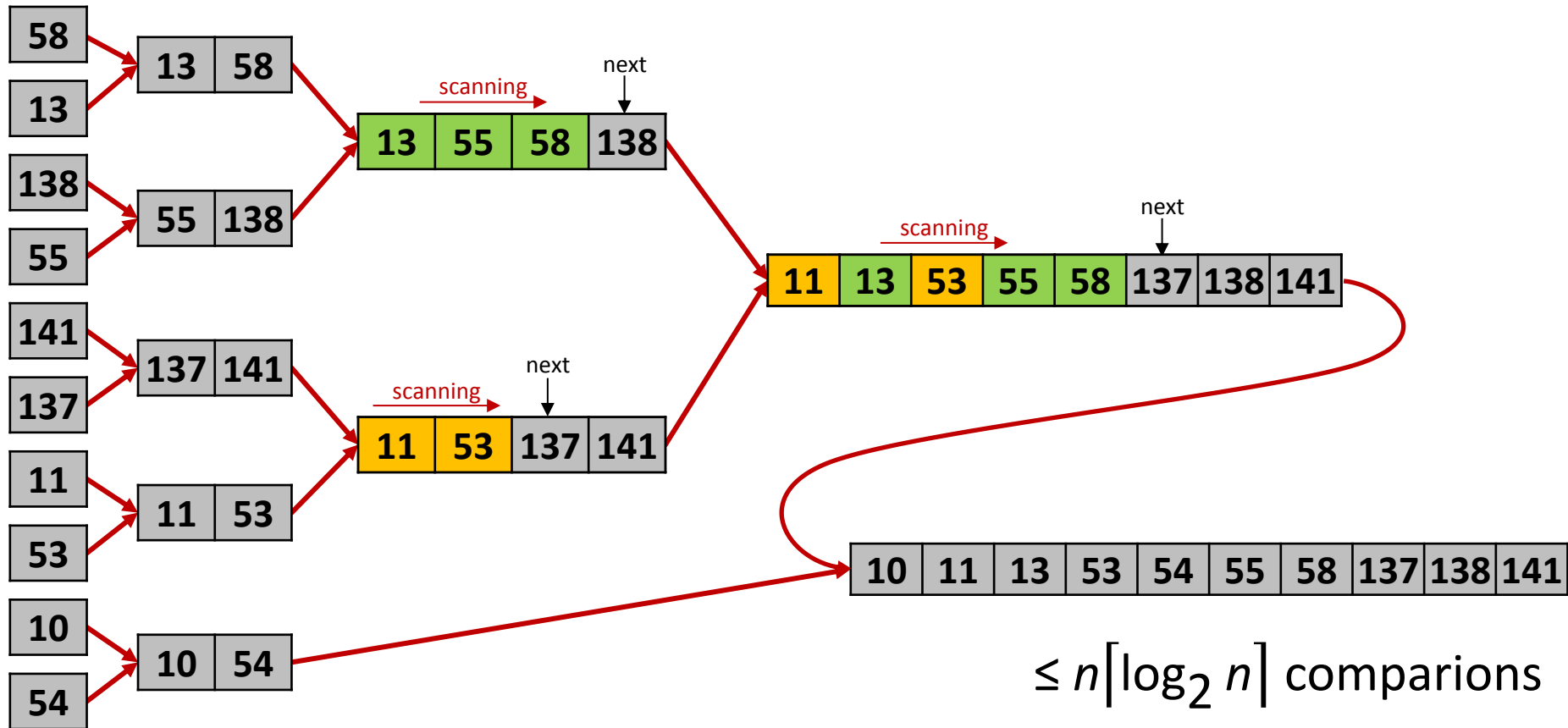
Output **10** **11** **13** **53** **54** **55** **58** **137** **138** **141**

00001010 00001010 00001101 00110101 00110110 00110111 00111010 10001001 10001010 10001101



# Comparison Based Sorting

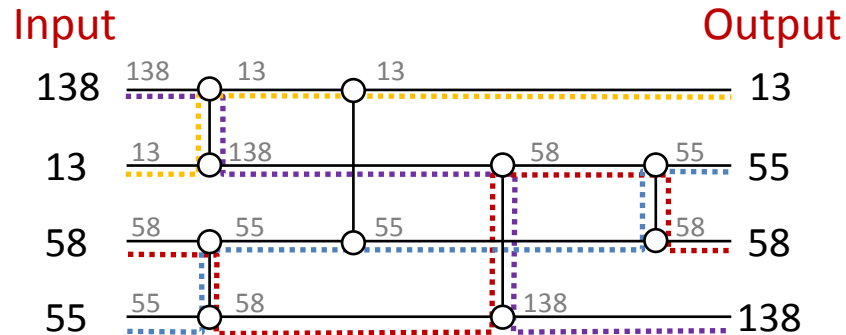
- Merge sort: von Neumann 1945



$\leq n \lceil \log_2 n \rceil$  comparisons

- Ford et al. 1959: Lower bound  $\geq n \log_2 n - \frac{n}{\ln 2}$  comparisons

# Sorting Networks

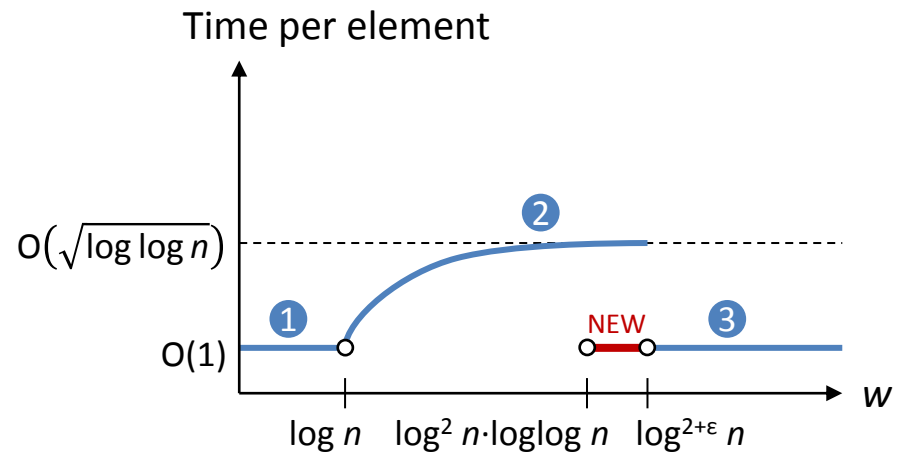


- Original motivated by hardware implementations
- Renewed interest since *data-oblivious*
  - applications in privacy preserving computations
  - allows for parallel computations on bit-level and GPU

$O(n \cdot \log^2 n)$  size : Batcher 1968

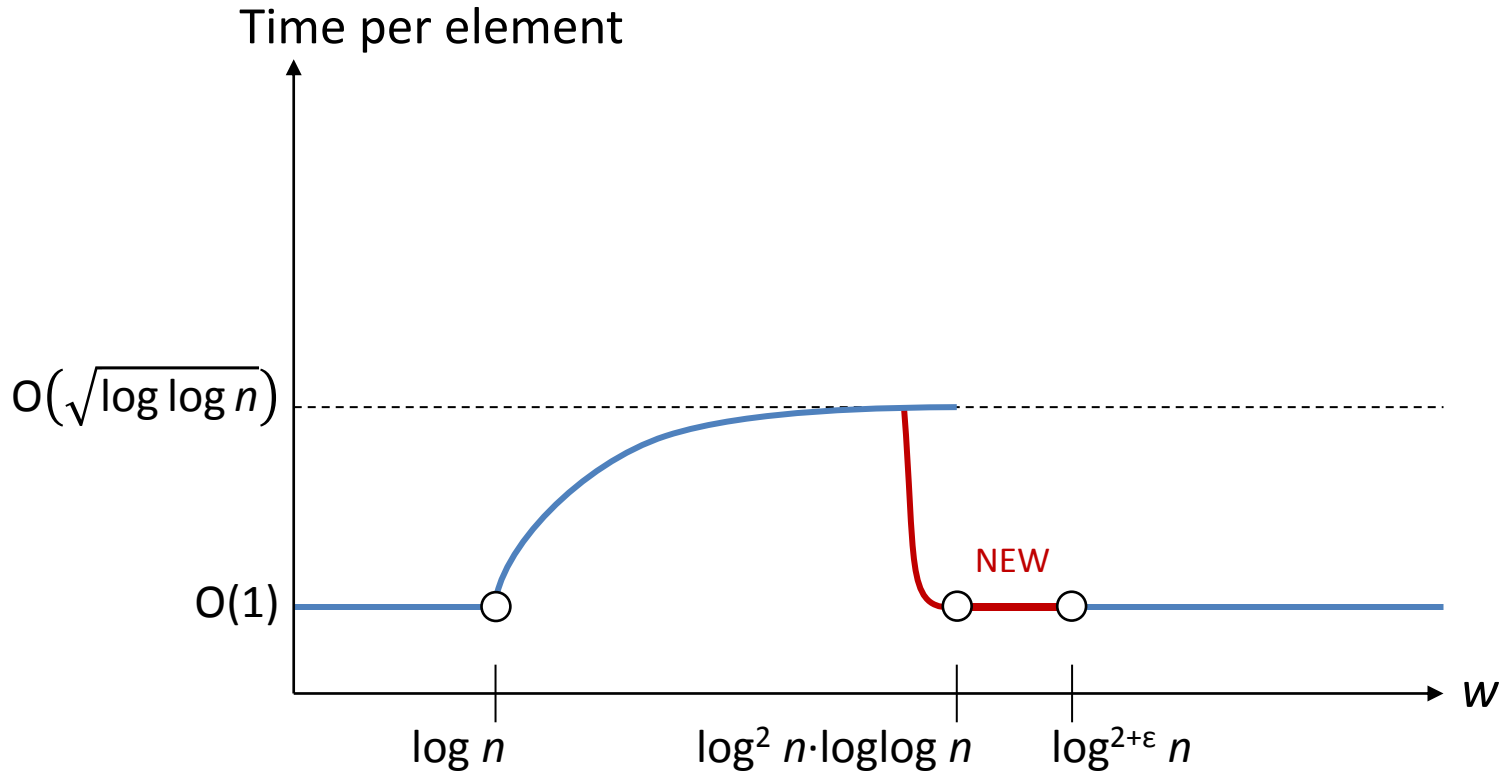
$O(n \cdot \log n)$  size : Ajtai, Komlós, Szemerédi 1983; Paterson 1990; Goodrich 2014

# Results



1	Bucket sort	$O(n+2^w)$	
	Radix sort; Hollerith 1887	$O\left(n \frac{w}{\log n}\right)$	
	van Emde Boas 1975 Willard 1983	$O(n \log w)$	superlinear space expected
	Kirkpatrick and Reicsh 1983	$O\left(n \log \frac{w}{\log n}\right)$	
	Merge sort: von Neumann 1945	$O(n \log n)$	comparison based optimal
2	Thorup and Han 2002	$O(n \sqrt{\log(w/\log n)})$ $O(n \sqrt{\log \log n})$	expected
3	Andersson et al. 1998	$O(n)$	expected, $w \geq \Omega(\log^{2+\epsilon} n)$
	<b>NEW</b>	$O(n)$	expected, $w \geq \Omega(\log^2 n \cdot \log \log n)$

# Open Question



Sorting in linear time for all  $w$ ?