

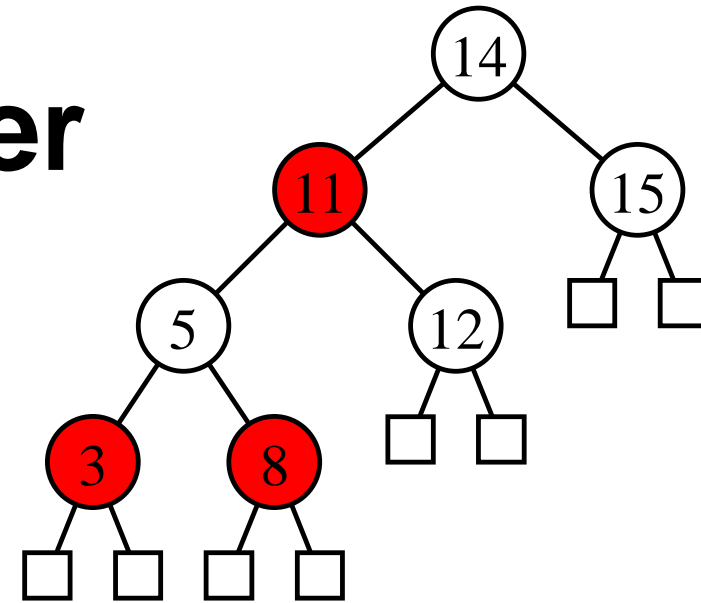
# Algoritmer og Datastrukturer

Balancerede søgetræer: Rød-sorter søgetræer  
[CLRS, kapitel 13]

# Rød-Sorte Søgetræer

## Mål

Søgetræer med dybde  $O(\log n)$



## Invarianter

- Hver knude er enten **RØD** eller **SORT**
- Hver **RØD** knude har en **SORT** far
- Alle stier fra roden til et eksternt blad har **samme antal sorte knuder**

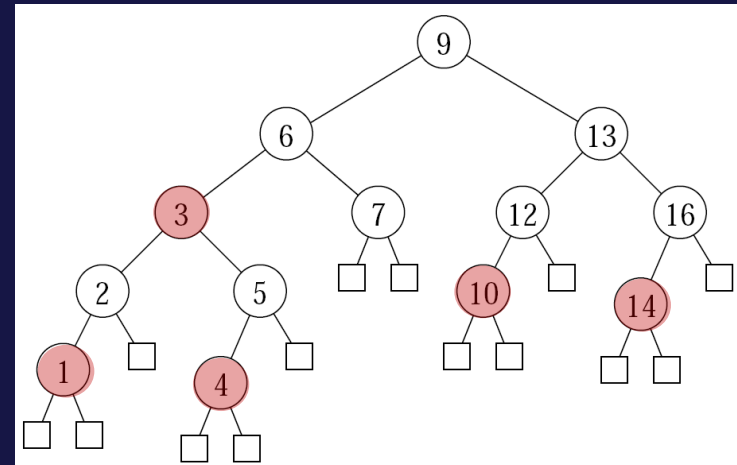
## Sætning

Et rød-sort træ har højde  $\leq 2 \cdot \log(n+1)$

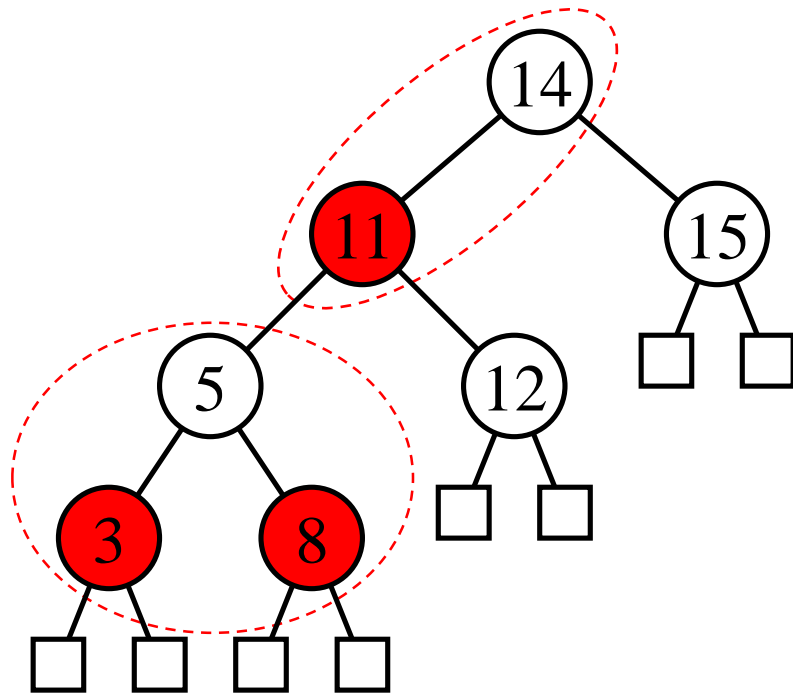
# Hvilke knuder skal farves røde for at træet bliver et rød-sort søgetræ?



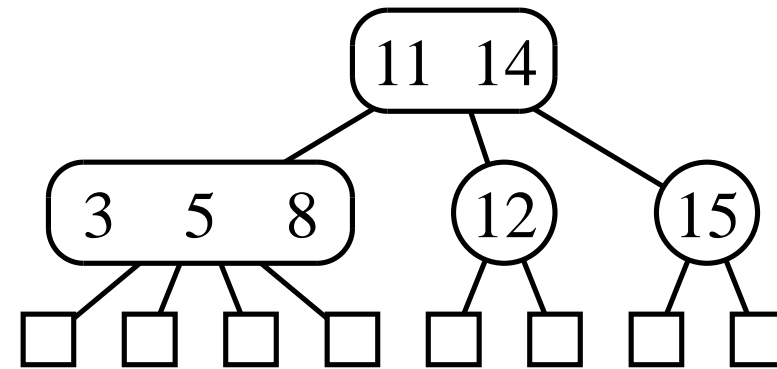
- a) 1, 4, 10, 14
- b) 2, 5, 10, 14
- c) 1, 3, 4, 7, 12, 16
- d) 1, 3, 4, 10, 14
- e) ingen
- f) ved ikke



# Rød-Sorte vs (2-4)-træer



Rød-sort træ

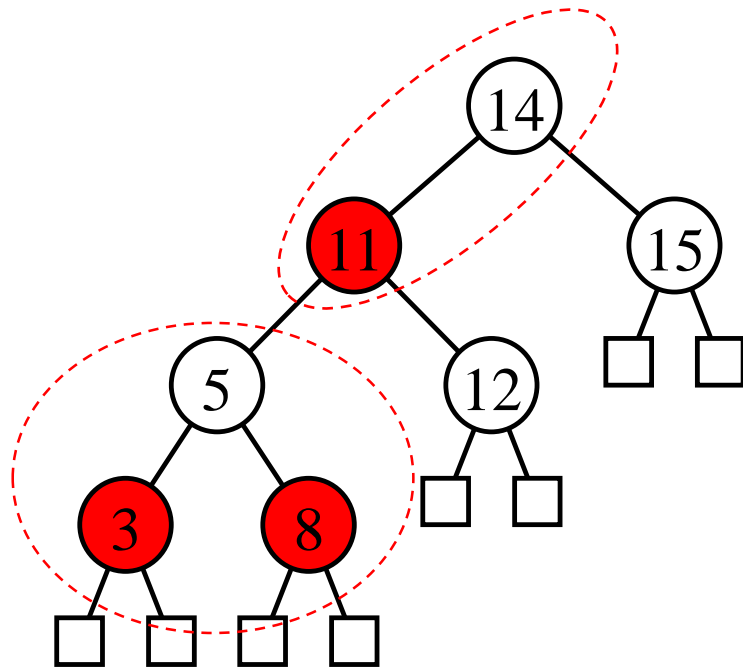


(2,4)-træ

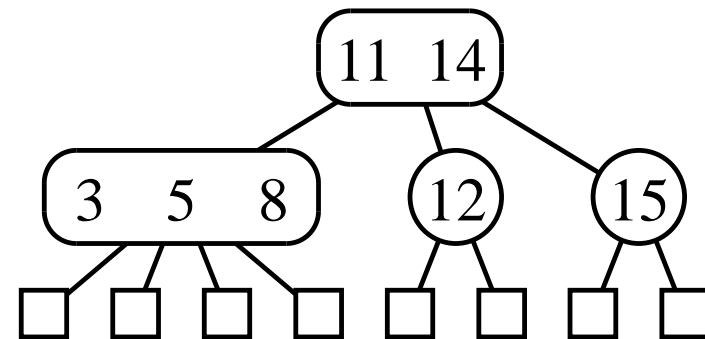
**(2,4)-træ**  $\equiv$  **rød-sort træ** hvor alle røde knuder er slået sammen med faderen

# Egenskaber ved (2,4)-træer

- Alle interne knuder har mellem **2 og 4 børn**
- Knude med  $i$  **børn** gemmer  $i-1$  **elementer**
- Stier fra roden til et eksternt blad er lige lange

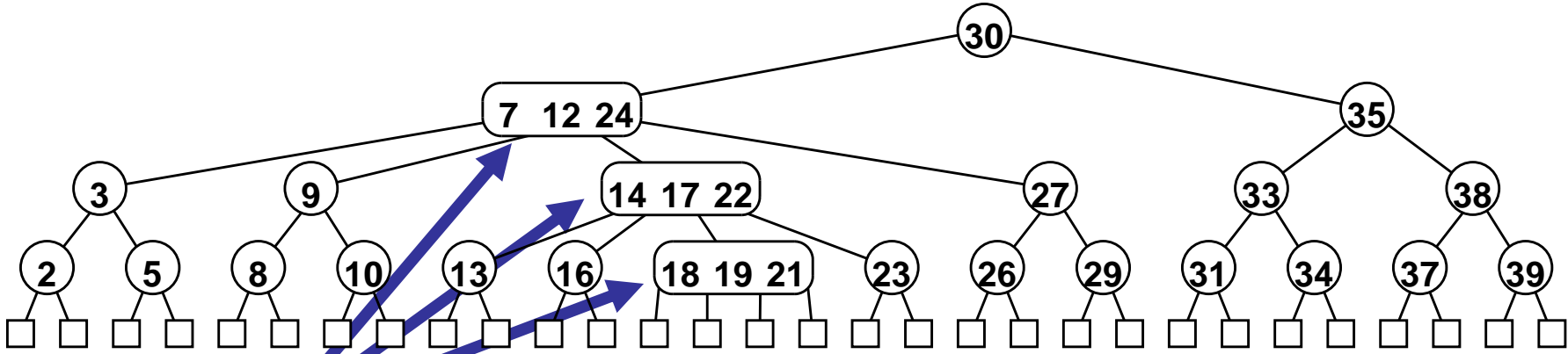


Rød-sort træ

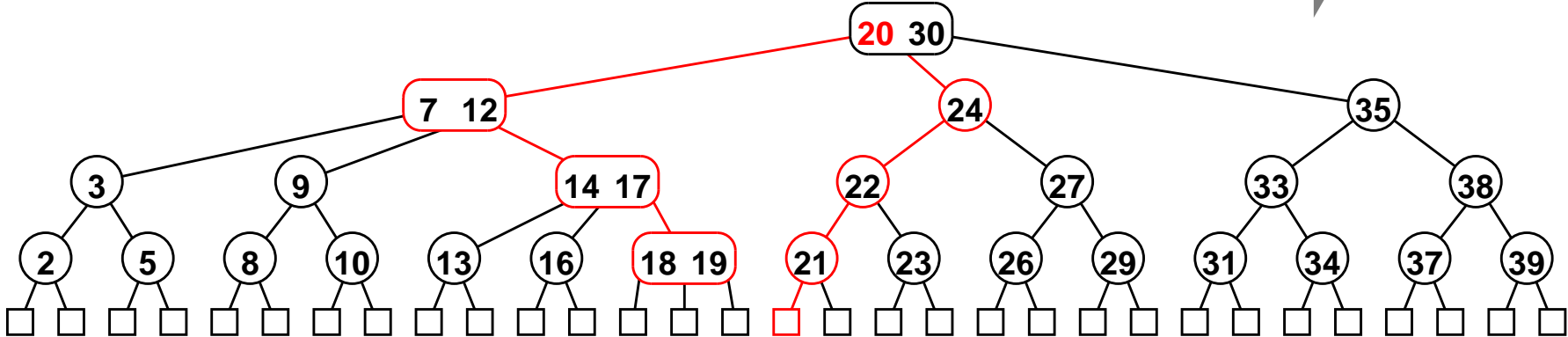
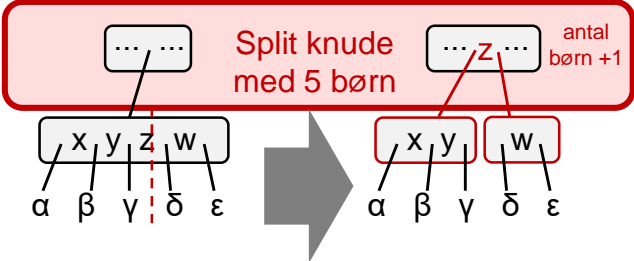


(2,4)-træ

# Indsættelse i et (2,4)-træ

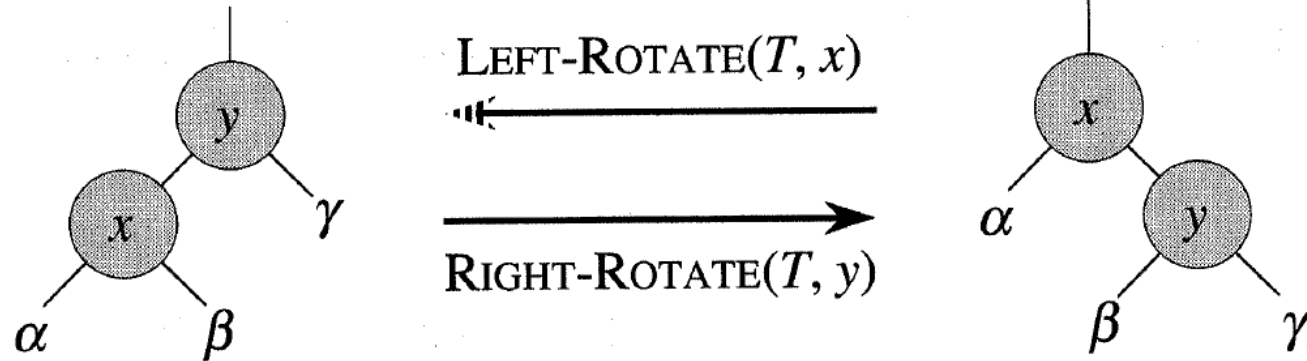


**Splittes** i to knuder og midterste nøgle flyttes et niveau op



# Opdateringer af Rød-Sorte Træer

# Rotationer



$\text{LEFT-ROTATE}(T, x)$

```
1   $y = x.\text{right}$            // set  $y$ 
2   $x.\text{right} = y.\text{left}$        // turn  $y$ 's left subtree into  $x$ 's right subtree
3  if  $y.\text{left} \neq T.\text{nil}$ 
4       $y.\text{left}.p = x$ 
5   $y.p = x.p$                  // link  $x$ 's parent to  $y$ 
6  if  $x.p == T.\text{nil}$ 
7       $T.\text{root} = y$ 
8  elseif  $x == x.p.\text{left}$ 
9       $x.p.\text{left} = y$ 
10 else  $x.p.\text{right} = y$ 
11  $y.\text{left} = x$              // put  $x$  on  $y$ 's left
12  $x.p = y$ 
```



**Insert**

Ubalanceret  
indsættelse

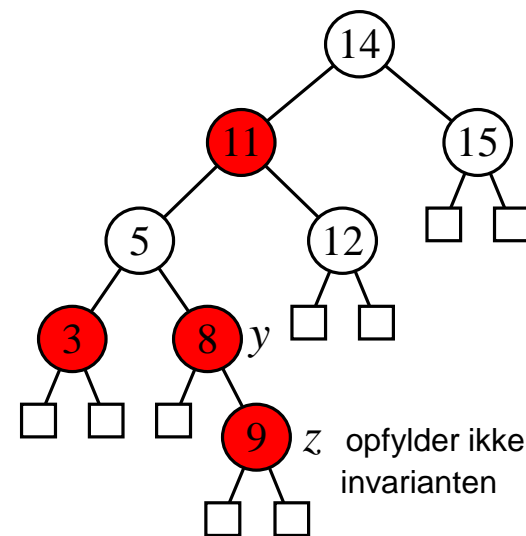
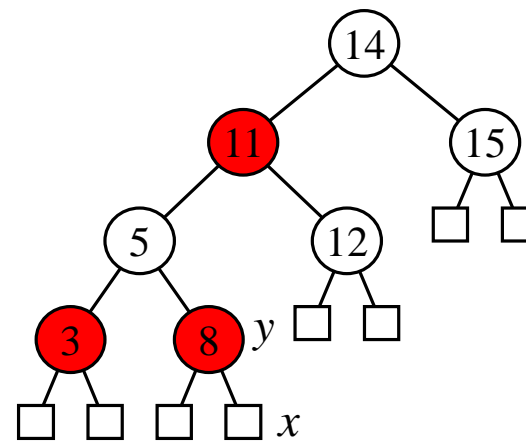
søg

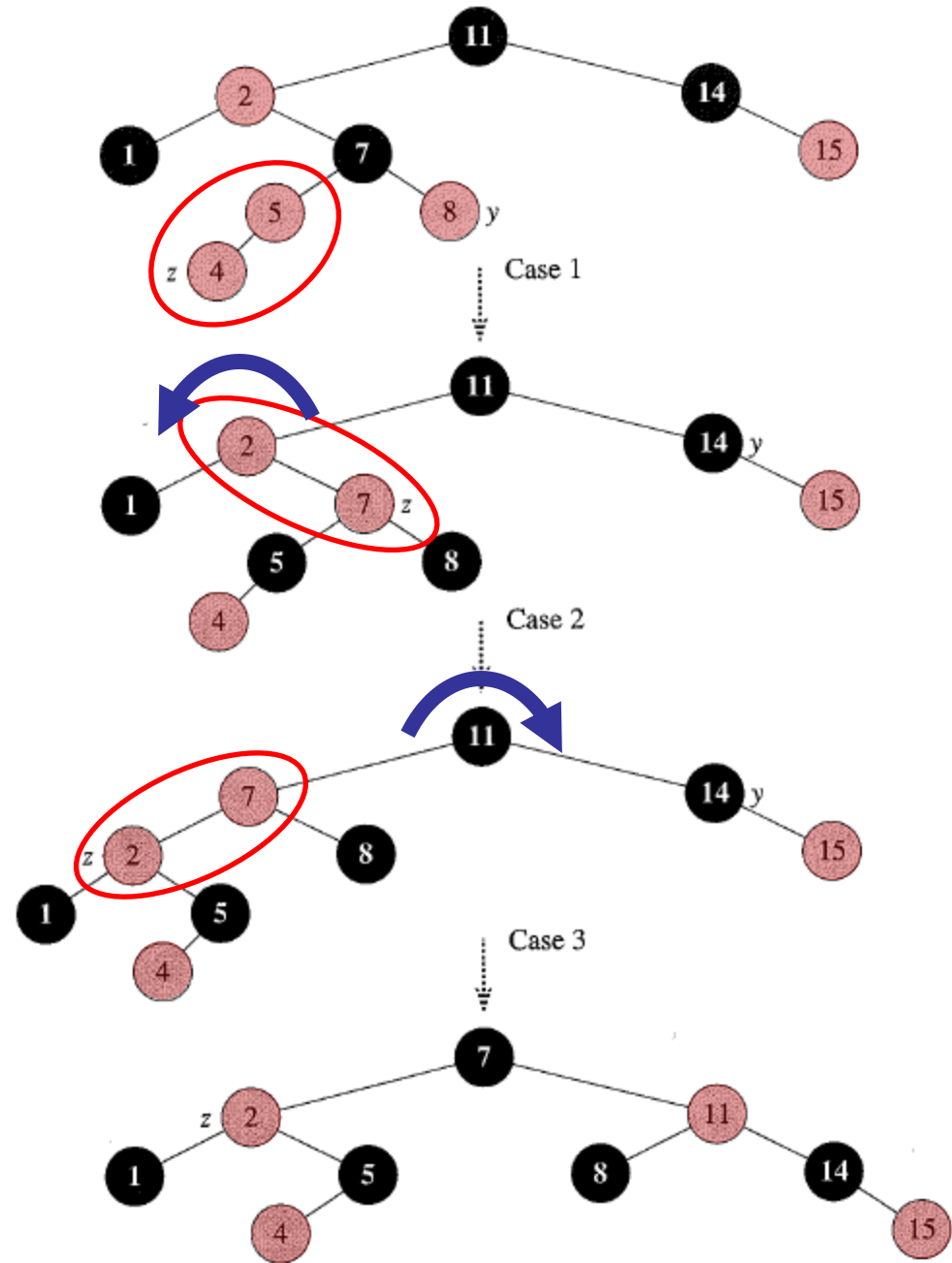
opret z

## RB-INSERT( $T, z$ )

```
1   $y = T.nil$ 
2   $x = T.root$ 
3  while  $x \neq T.nil$ 
4       $y = x$ 
5      if  $z.key < x.key$ 
6           $x = x.left$ 
7      else  $x = x.right$ 
8   $z.p = y$ 
9  if  $y == T.nil$ 
10      $T.root = z$ 
11  elseif  $z.key < y.key$ 
12      $y.left = z$ 
13  else  $y.right = z$ 
14   $z.left = T.nil$ 
15   $z.right = T.nil$ 
16   $z.color = RED$ 
17  RB-INSERT-FIXUP( $T, z$ )
```

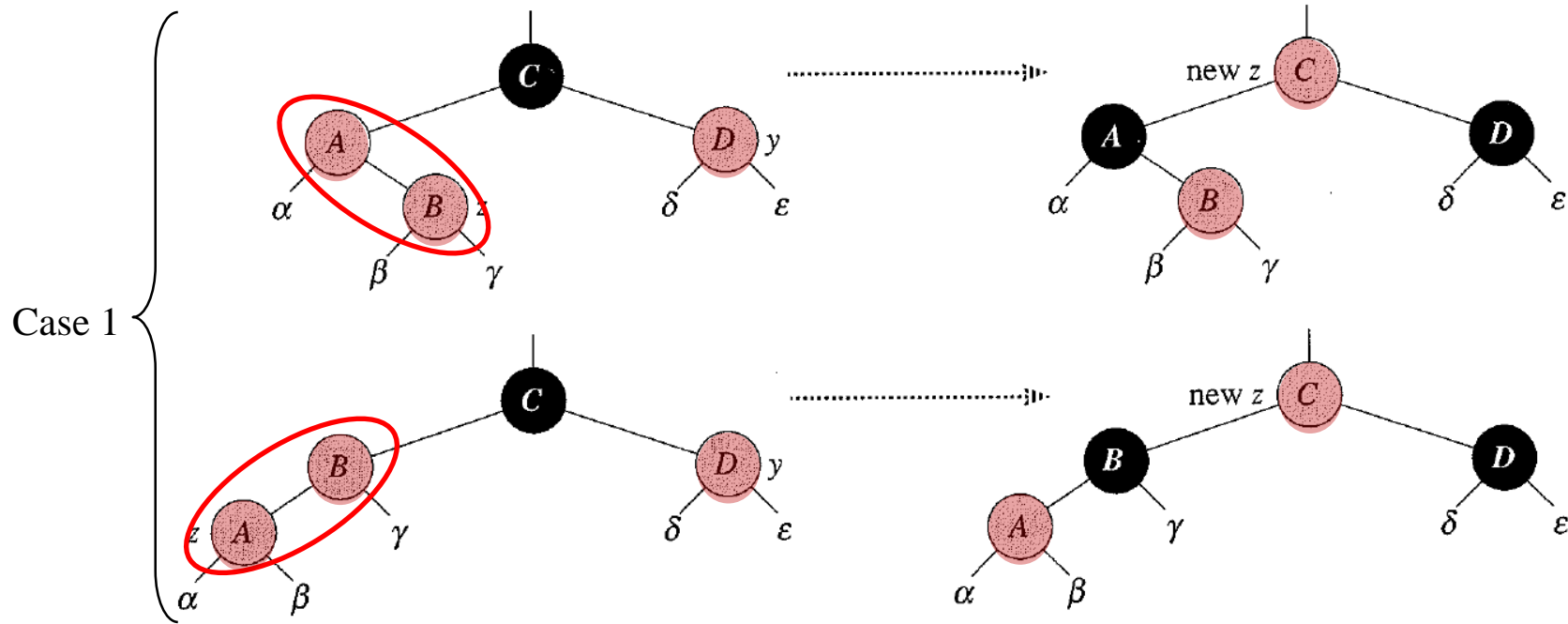
## Insert(9)



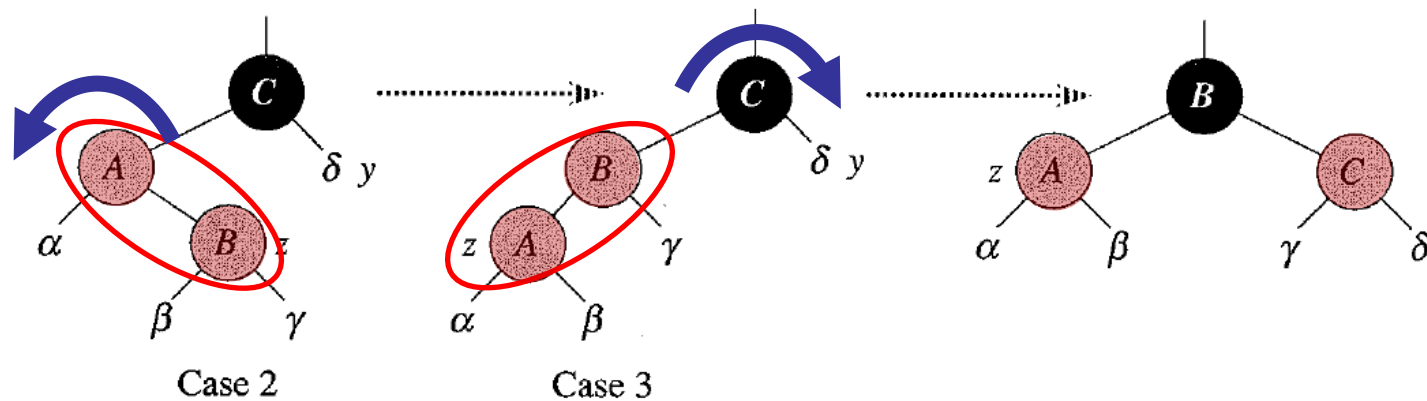


# Indsættelse i rød-sort træer: rebalancering

Omfarv  $C$   
og dens to  
røde børn



Afslut

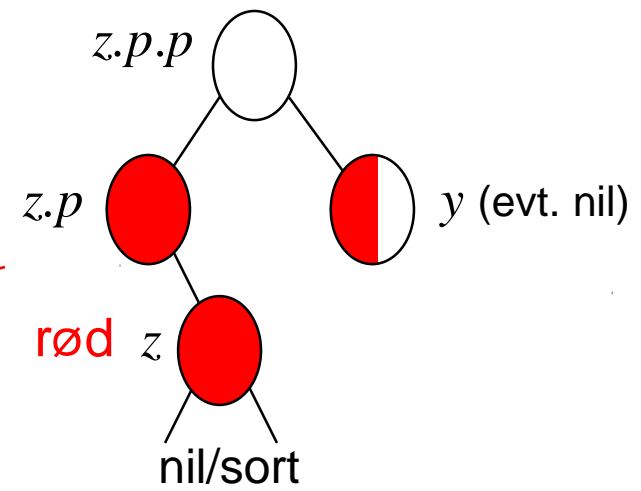


## RB-INSERT-FIXUP( $T, z$ )

```

1  while  $z.p.color == RED$ 
2      if  $z.p == z.p.p.left$ 
3           $y = z.p.p.right$ 
4          if  $y.color == RED$ 
5               $z.p.color = BLACK$  // case 1
6               $y.color = BLACK$  // case 1
7               $z.p.p.color = RED$  // case 1
8               $z = z.p.p$  // case 1
9          else if  $z == z.p.right$ 
10              $z = z.p$  // case 2
11             LEFT-ROTATE( $T, z$ ) // case 2
12              $z.p.color = BLACK$  // case 3
13              $z.p.p.color = RED$  // case 3
14             RIGHT-ROTATE( $T, z.p.p$ ) // case 3
15         else (same as then clause
                with “right” and “left” exchanged)
16      $T.root.color = BLACK$ 

```



# Worst-case antal rotationer ved indsættelse i et rød-sort træ ?



a) 1

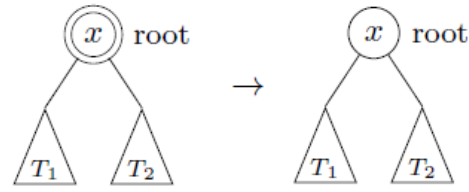
b) 2

c)  $\log n$

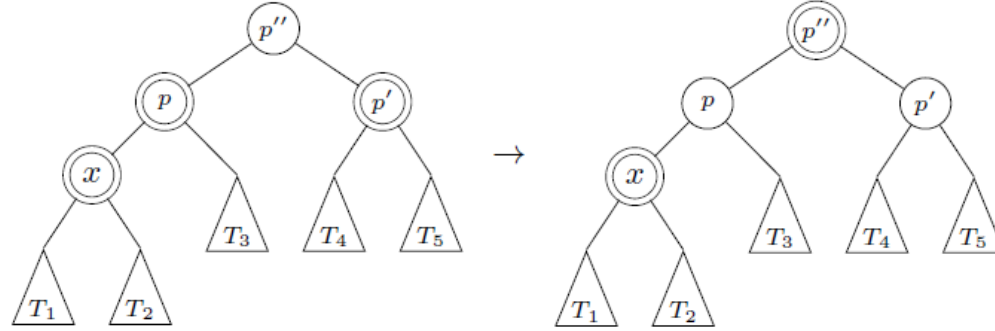
d)  $2 \cdot \log n$

e) Ved ikke

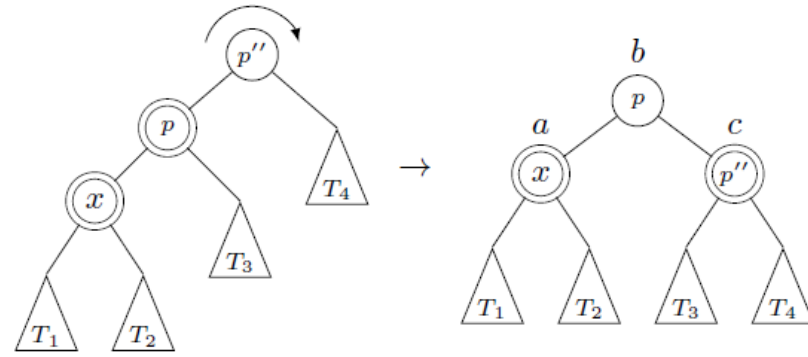
[B]



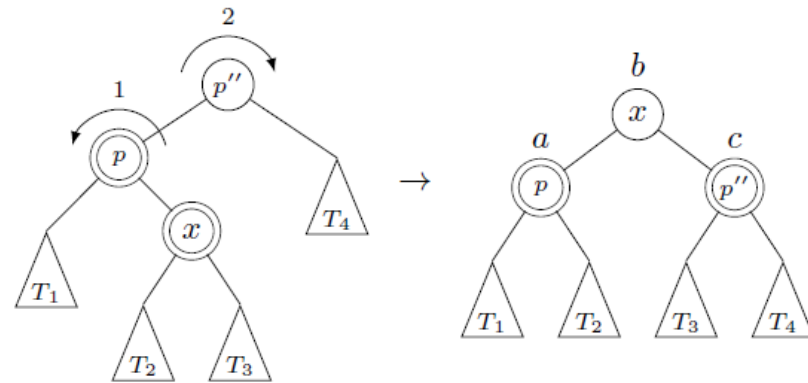
case I



case II



case III



**Delete**



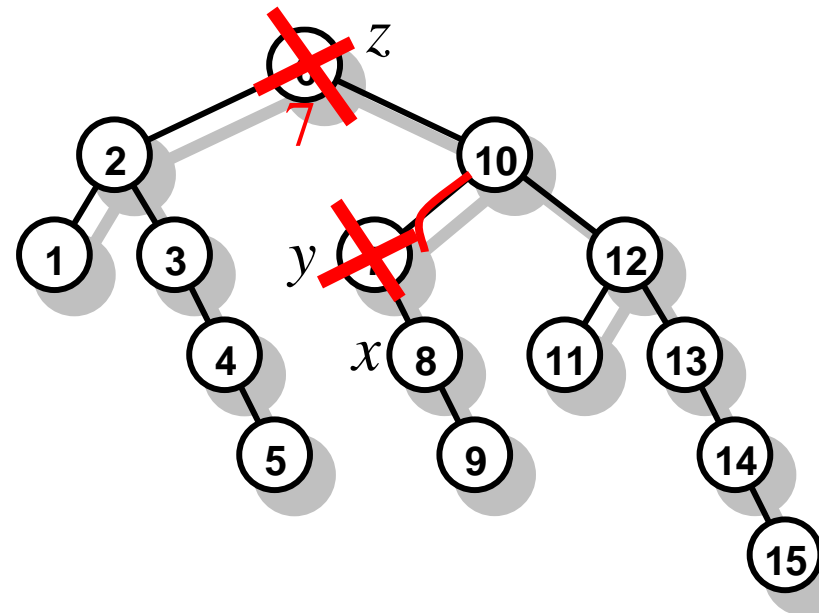
RB-DELETE( $T, z$ )

Ubalanceret slettelse

```

1   $y = z$ 
2   $y\text{-original-color} = y.\text{color}$ 
3  if  $z.\text{left} == T.\text{nil}$ 
4       $x = z.\text{right}$ 
5      RB-TRANSPLANT( $T, z, z.\text{right}$ )
6  elseif  $z.\text{right} == T.\text{nil}$ 
7       $x = z.\text{left}$ 
8      RB-TRANSPLANT( $T, z, z.\text{left}$ )
9  else  $y = \text{TREE-MINIMUM}(z.\text{right})$ 
10      $y\text{-original-color} = y.\text{color}$ 
11      $x = y.\text{right}$ 
12     if  $y.p == z$ 
13          $x.p = y$ 
14     else RB-TRANSPLANT( $T, y, y.\text{right}$ )
15          $y.\text{right} = z.\text{right}$ 
16          $y.\text{right}.p = y$ 
17     RB-TRANSPLANT( $T, z, y$ )
18      $y.\text{left} = z.\text{left}$ 
19      $y.\text{left}.p = y$ 
20      $y.\text{color} = z.\text{color}$ 
21 if  $y\text{-original-color} == \text{BLACK}$ 
22     RB-DELETE-FIXUP( $T, x$ )

```

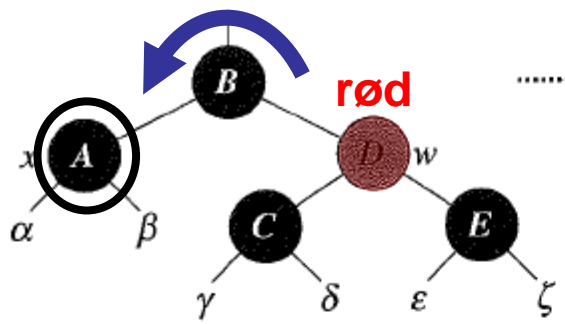
 **$x$  og  $y.\text{right}$  kan være  $T.\text{nil}$  !!!**RB-TRANSPLANT( $T, u, v$ )

```

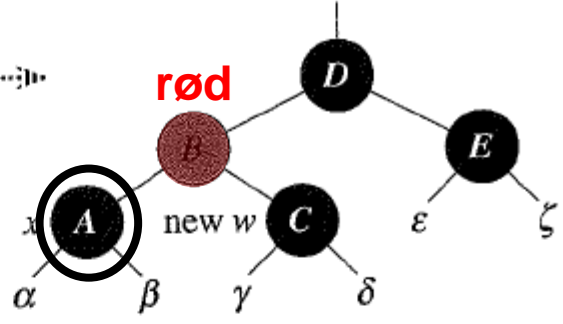
1  if  $u.p == T.\text{nil}$ 
2       $T.\text{root} = v$ 
3  elseif  $u == u.p.\text{left}$ 
4       $u.p.\text{left} = v$ 
5  else  $u.p.\text{right} = v$ 
6   $v.p = u.p$ 

```

$x$   
"dobbel sort"

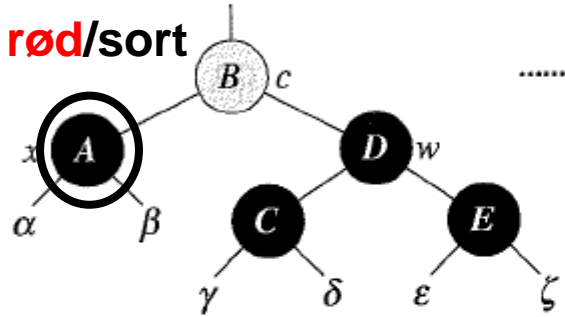


Case 1

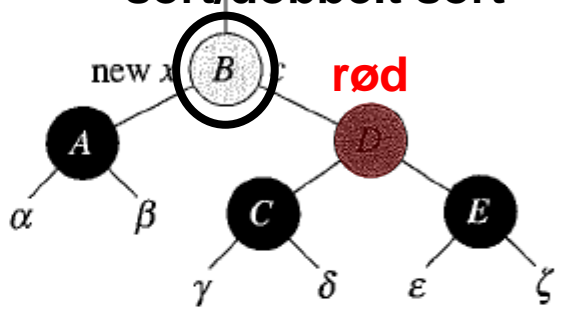


sort/dobbel sort

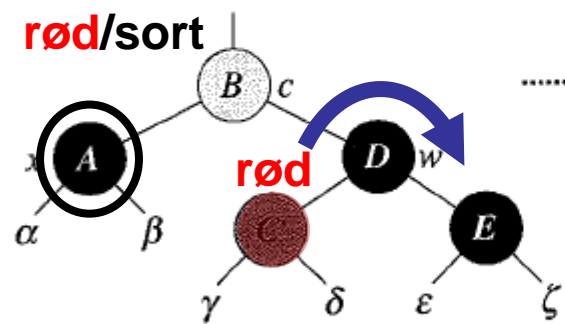
Case 1 → 2,3,4



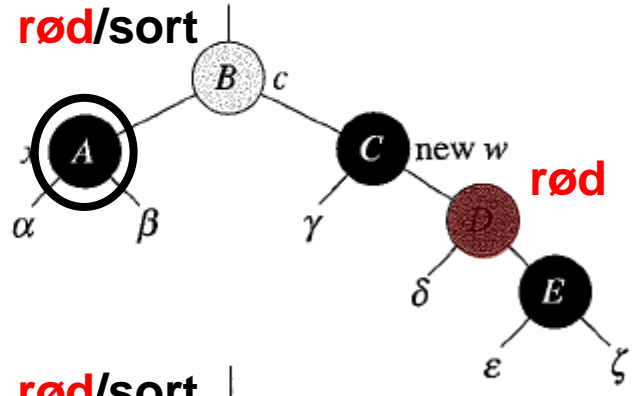
Case 2



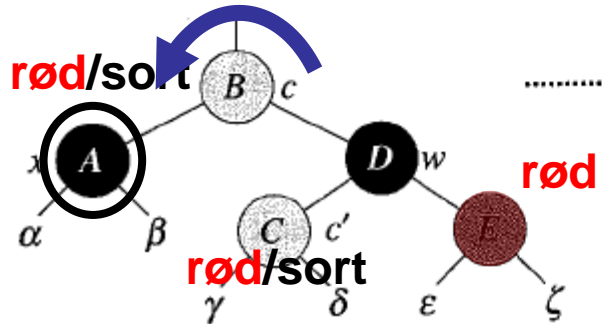
Case 2 →  
Problemet  $x$  flyttet  
tættere på roden  
eller færdig



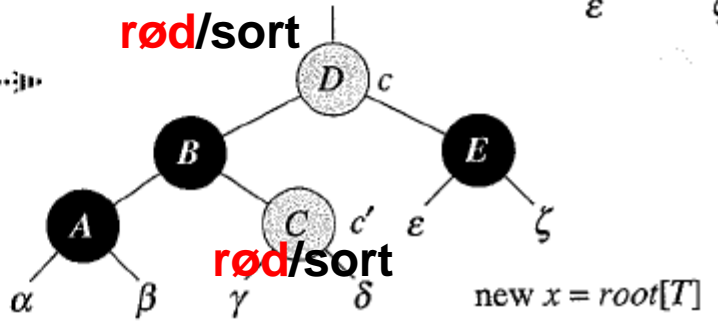
Case 3



Case 3 → 4  
(og færdig)



Case 4



Case 4 → Færdig

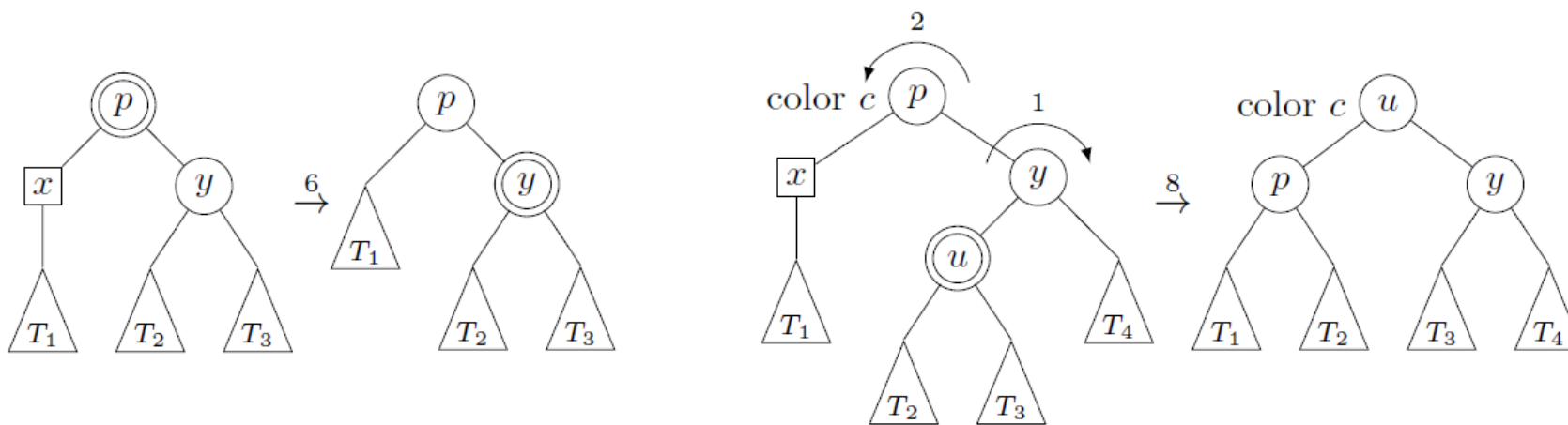
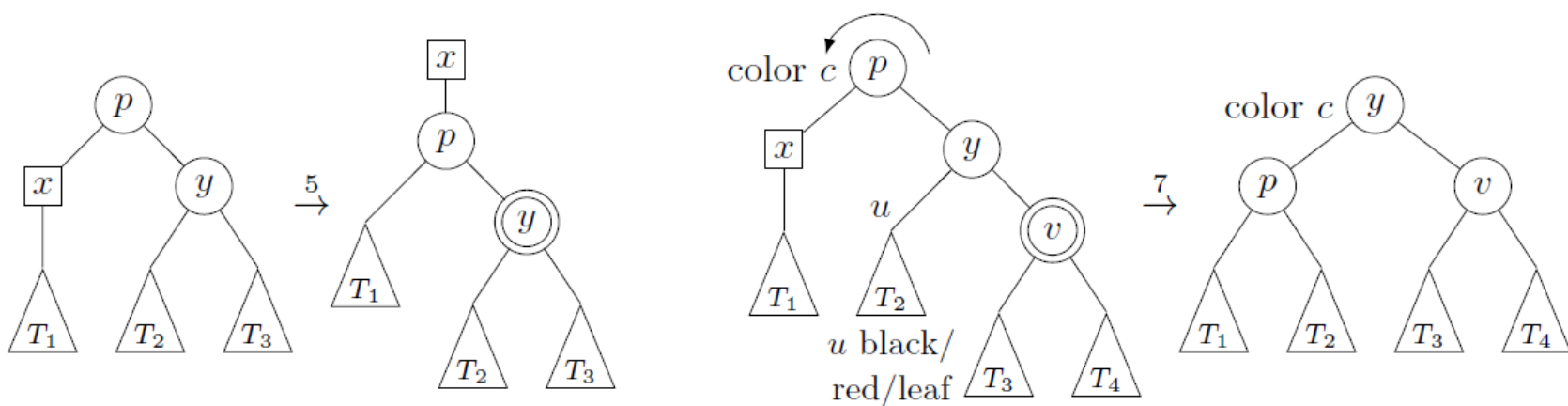
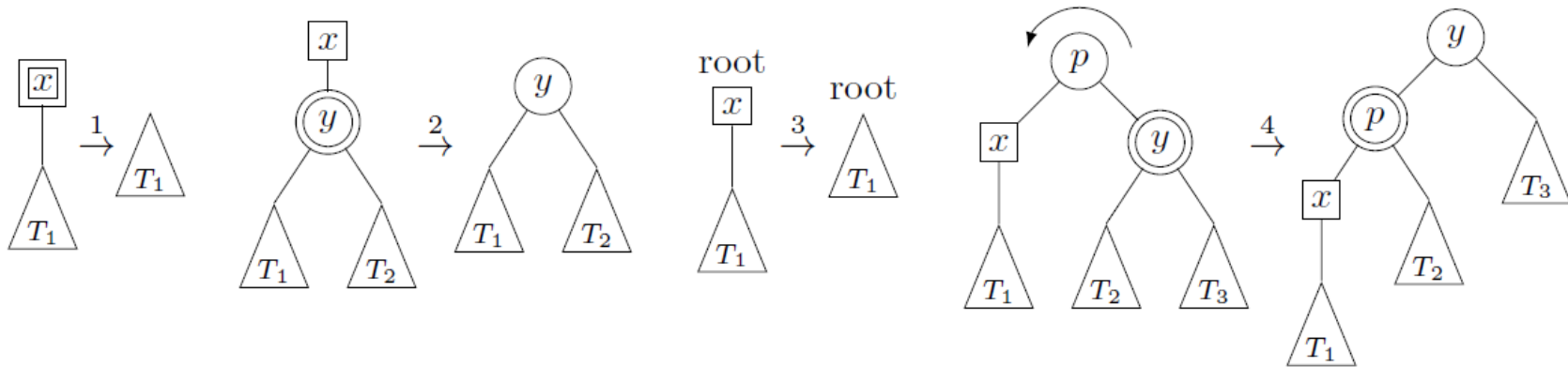
## RB-DELETE-FIXUP( $T, x$ )

```

1  while  $x \neq T.root$  and  $x.color == BLACK$ 
2      if  $x == x.p.left$ 
3           $w = x.p.right$ 
4          if  $w.color == RED$ 
5               $w.color = BLACK$  // case 1
6               $x.p.color = RED$  // case 1
7              LEFT-ROTATE( $T, x.p$ ) // case 1
8               $w = x.p.right$  // case 1
9          if  $w.left.color == BLACK$  and  $w.right.color == BLACK$ 
10              $w.color = RED$  // case 2
11              $x = x.p$  // case 2
12         else if  $w.right.color == BLACK$ 
13              $w.left.color = BLACK$  // case 3
14              $w.color = RED$  // case 3
15             RIGHT-ROTATE( $T, w$ ) // case 3
16              $w = x.p.right$  // case 3
17              $w.color = x.p.color$  // case 4
18              $x.p.color = BLACK$  // case 4
19              $w.right.color = BLACK$  // case 4
20             LEFT-ROTATE( $T, x.p$ ) // case 4
21              $x = T.root$  // case 4
22         else (same as then clause with “right” and “left” exchanged)
23      $x.color = BLACK$ 

```

[B]



# Dynamisk Ordbog :

## Rød-Sorte Træer

<b>Search(<math>S, x</math>)</b>	$O(\log n)$
<b>Insert(<math>S, x</math>)</b>	
<b>Delete(<math>S, x</math>)</b>	

# Rød-Sorte Træer i Java og C++

	Java (JDK SE 12.0.1)	C++ (GCC 9.1)
Metode	java.util.TreeMap	std::map
Dokumentation	<a href="https://docs.oracle.com/javase/10/docs/api/java/util/TreeMap.html">https://docs.oracle.com/javase/10/docs/api/java/util/TreeMap.html</a>	<a href="http://www.cplusplus.com/reference/map/map/">http://www.cplusplus.com/reference/map/map/</a>
Kildekode	Installerer JDK fra <a href="http://www.oracle.com/technetwork/java/javase/downloads/">www.oracle.com/technetwork/java/javase/downloads/</a> Fil java.base\java\util\ TreeMap.java fra C:\ProgramFiles\Java\ jdk-12.0.1\lib\src.zip	<a href="https://github.com/gcc-mirror/gcc/blob/master/libstdc++-v3/include/bits/stl_tree.h">https://github.com/gcc-mirror/gcc/blob/master/libstdc++-v3/include/bits/stl_tree.h</a>