

#JoinUs

EUROVISION
SONG CONTEST
COPENHAGEN 2014

Redefining the Eurovision voting
process using linear ECC

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Joint work with Bernardo David, Irene Giacomelli, Jesper
Buus Nielsen Ignacio Cascudo Pueyo

The artists perform



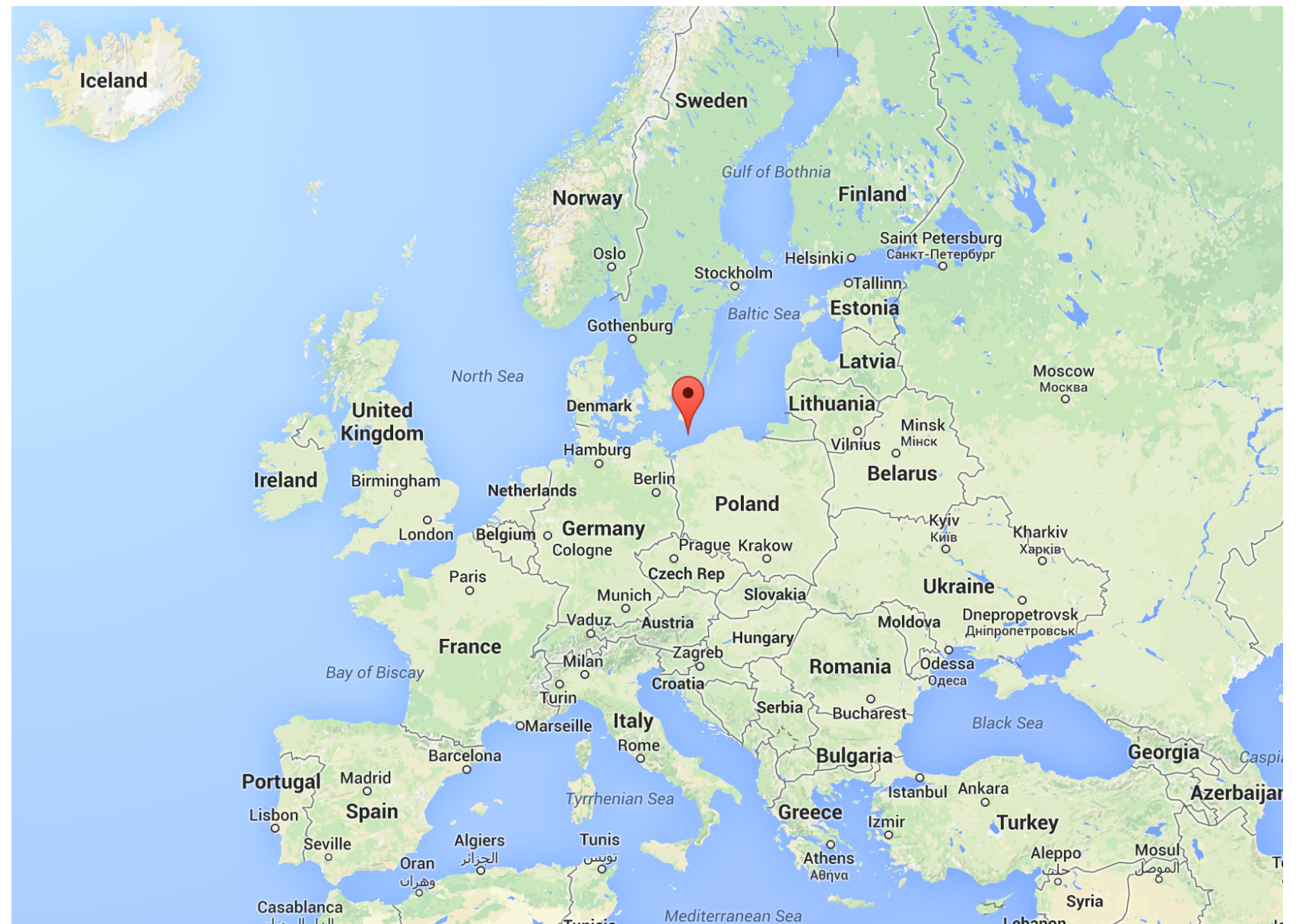
Each in their own way...



The voting process

EBU

OPERATING EUROVISION AND EURORADIO



The voting process



The voting process



A concrete cheating strategy

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- Mole within the EBU who leaks all votes as they come in to the UK
- The UK (as well as their friends) adjust their votes accordingly to maximize their winning probability



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- *In addition our solution is additively homomorphic, making it possible to anonymize the entire voting process
- But might ruin the dramatics of the voting count up

Our Scheme

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S

s_1
s_2
s_3
s_4
\dots
\dots
s_{2n-1}
s_{2n}

R

$\left. \begin{array}{c} s_1 \\ s_2 \end{array} \right\} \binom{2}{1}\text{-OT}$
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- Means R learns n seeds, based on choice bits in OTs
- Use s_i as seed for PRG and use output y_i as OTP

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	\mathbf{c}_2
+	\mathbf{z}

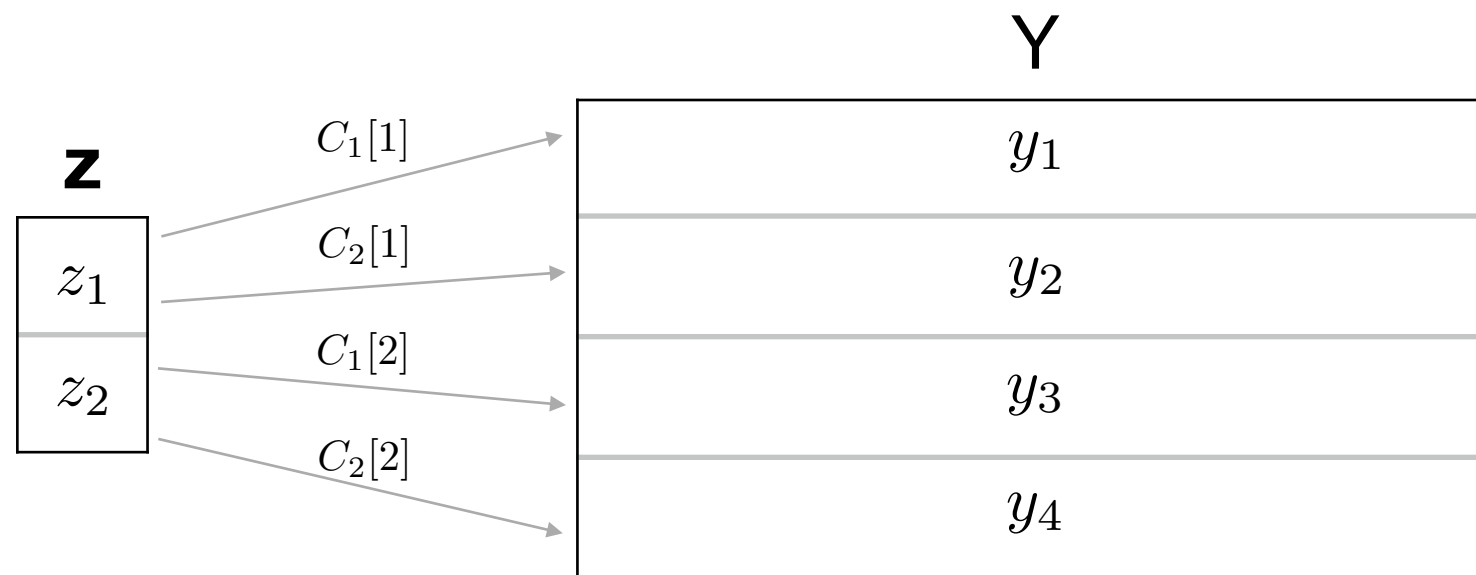
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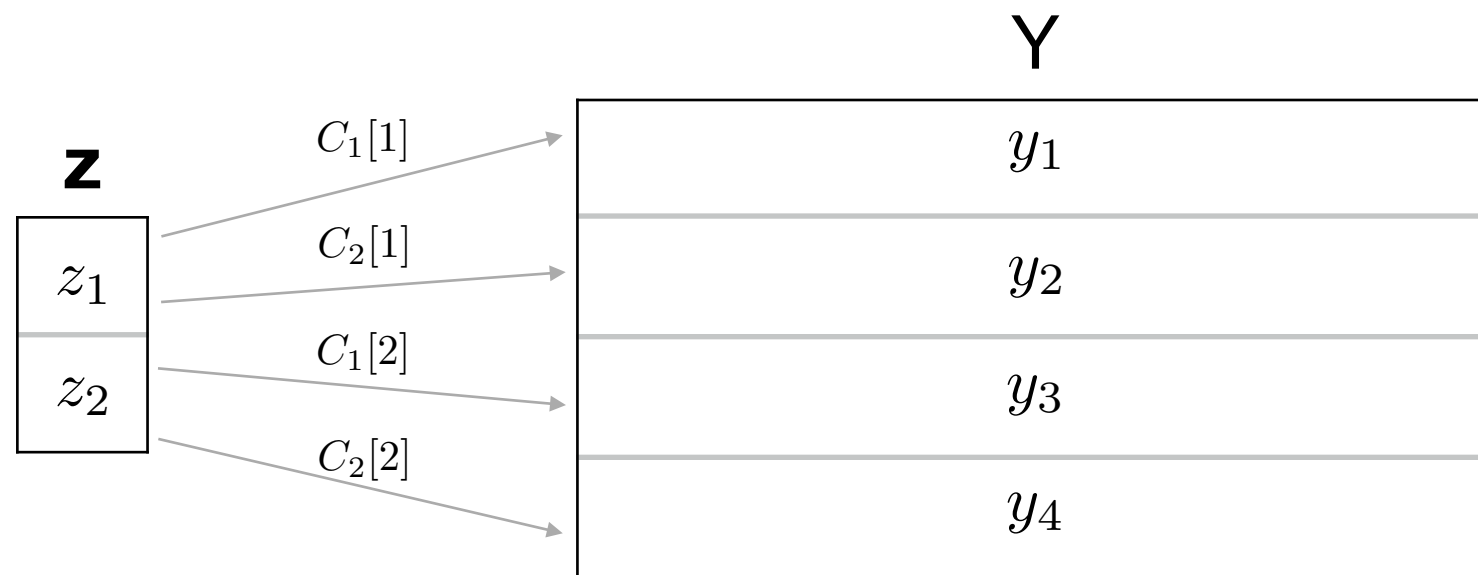
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- Entry - and pair-wise pad \mathbf{c}_1 and \mathbf{c}_2 using the y_i 's. Send the padded vectors to R

In a picture

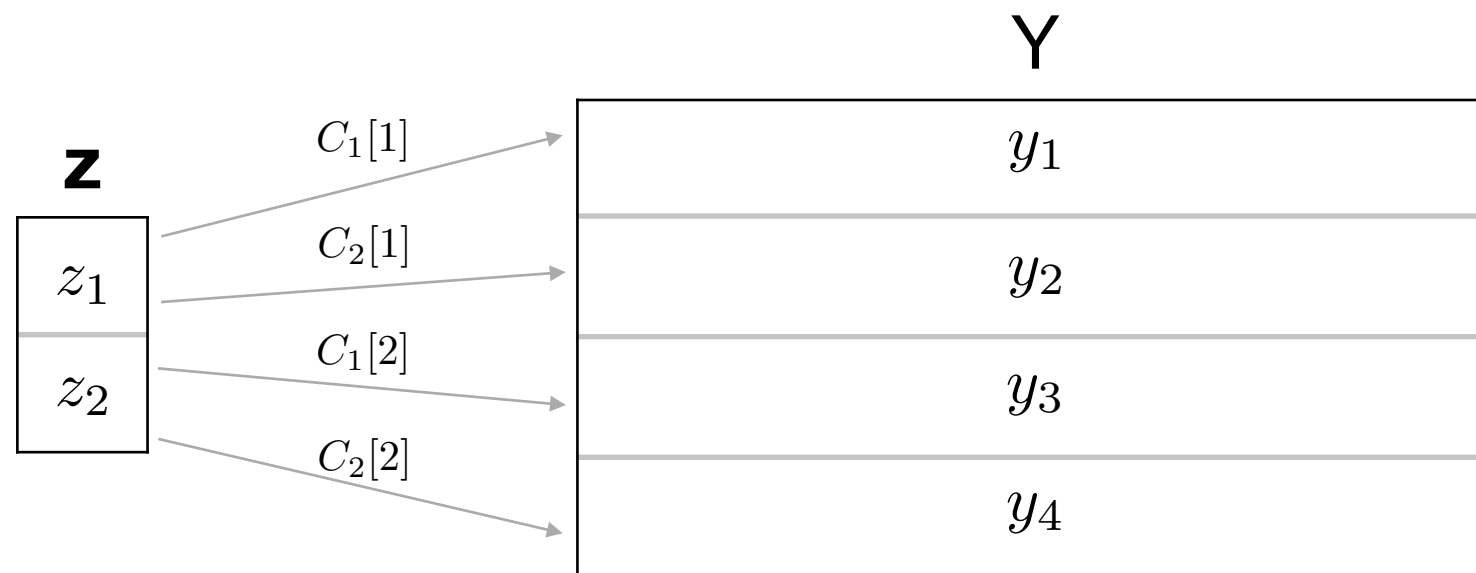


In a picture



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- Receiver learns half the shares, i.e. perfectly hiding
- Sender unaware which shares are “watched”. Need to change d shares in order to change code-word.
=> Probability of cheating 2^{-d} (more or less)