

Impartial decision making among peers

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conflict of interest in collective decision making:

my selfish interest *corrupts* the report of my subjective opinion

non corrupted information is more valuable: it produces an *impartial evaluation*

conflict of interests pervasive in collective decisions by and about peers

example: evaluate the merit of a peer's work, choose a winner among us, a ranking of us all

a necessary condition for the possibility of an impartial process:

- separate aspects of the decision related to *self interest* versus *opinions/views*

then a decision rule creates no conflict of interest if it only elicits opinions, and an agent's report **does not affect** her self interest

examples where the separation is plausible

self-interest

opinion

division of a dollar

my share

division of the remainder

award of a prize

do I win?

who wins if not me?

selecting webpages

am I in?

who is apart from me?

ranking by peers

what is my rank?

ranking of the others

- *Impartial division of a dollar*, G. de Clippel, H. Moulin and N. Tideman, *Journal of Economic Theory*, 2008.
- *Impartial award of a prize*, R. Holzman and H. Moulin, mimeo September 2010
- Sum of us: strategy-proof selection from the selectors, N. Alon, F. Fischer, A. Procaccia and M. Tennenholtz
- strategyproof and efficient allocation of private goods: Kato and Ohseto (building on the work of Hurwicz, Zhou, Serizawa and Weymark,..)

model 1: award of a prize

$$i \in N = \{1, 2, \dots, n\}$$

i 's message $m_i \in M_i$

award rule: $M_N \ni m \rightarrow f(m) \in N$

\rightarrow **Impartiality:** $f(m|{}^i m_i) = i \Leftrightarrow f(m|{}^i m'_i) = i$, for all i, m_i, m'_i

additional requirements:

- **No Discrimination:** $\forall i \exists m f(m) = i$
- **No Dummy:** $\forall i \exists m_i, m'_i, m_{-i} : f(m|{}^i m_i) \neq f(m|{}^i m'_i)$

both are (very) weak forms of symmetry among participants

note: full Anonymity impossible

Lemma (easy):

For $n \leq 3$ Impartiality \cap No Discrimination = Impartiality \cap No Dummy
= \emptyset

For $n = 4$, assume binary messages $m_i = 0, 1$

Impartiality \cap No Discrimination \cap No Dummy = $\{f^4\}$

up to relabeling agents and messages

$$f^4(\cdot, 0, 0, 0) = f^4(\cdot, 1, 1, 1) = 1; f^4(0, \cdot, 1, 0) = f^4(1, \cdot, 0, 1) = 2$$

$$f^4(1, 1, \cdot, 0) = f^4(0, 0, \cdot, 1) = 3; f^4(0, 1, 0, \cdot) = f^4(1, 0, 1, \cdot) = 4$$

for $n \geq 5$, there are many more rules

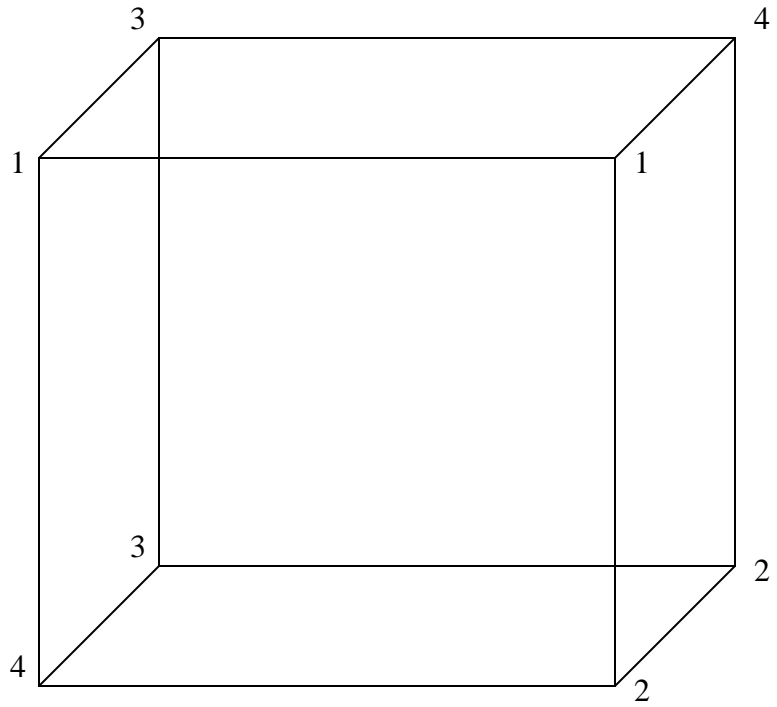
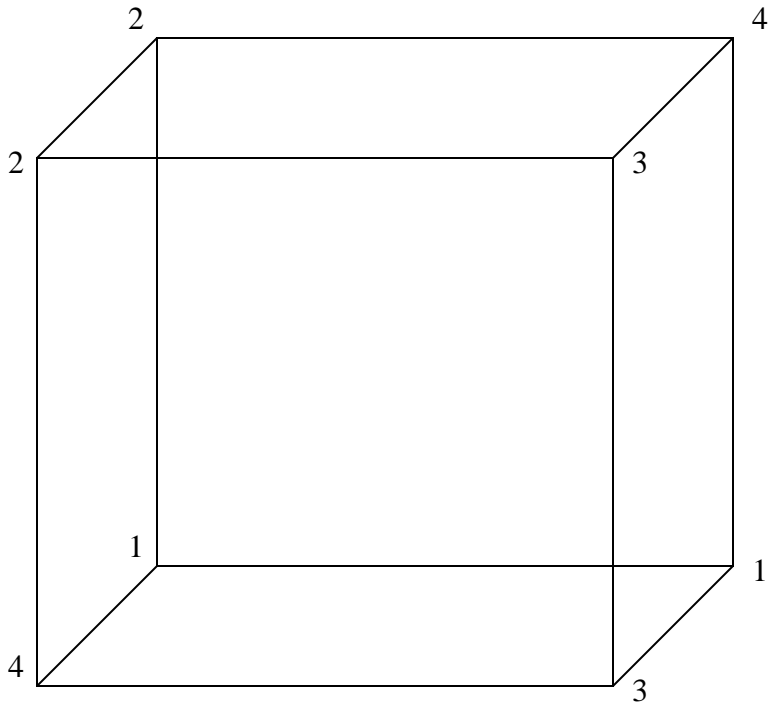


Figure 1

quota rules

everyone but the incumbent nominates someone (no self nomination)

$q > \frac{n}{2}$: *absolute* quota rule $I^{ab}(q)$: i wins if $\text{score}(i) \geq q$

$2 \leq q \leq \frac{n}{2}$ *relative* quota rule $I^r(q)$: i wins if $\text{score}(i) \geq \text{score}(j|N \setminus \{i\}) + q$ for all $j \neq i$

if no such winner, the incumbent wins

→ Impartial, No Discrimination, but the incumbent is a *dummy*

combine two of these rules

partition $N = N_1 \cup N_2$; choose q_1, q_2

step 1: run $I^{\varepsilon_1}(q_1)$ in N_1 ; stop if there is a winner

otherwise go to

step 2: N_1 vote to choose the incumbent $j \in N_2$, then run $I^{\varepsilon_2}(q_2)$ in N_2

\Rightarrow Impartial, No Discrimination, No Dummy

critique: unequal influence of N_1 versus N_2

toward a more equal distribution of the decision power:

i influences $j \stackrel{def}{\iff} \exists m \in M^N, m'_i \in M^i : f(m|{}^i m_i) = j \neq f(m|{}^i m'_i)$

Full mutual Influence: $\forall i, j \in N: i$ influences j

Full Influence \Rightarrow No Dummy and No Discrimination

nomination rules

simple and natural messages: $M_i = N \setminus \{i\}$ agent i *nominates* j

Monotonicity: $\forall i, j, i \neq j \forall m \in M_N : f(m) = j \Rightarrow f(m|_j^i) = j$

Anonymous ballots: for all $m, m' \in M_N$

$\{\forall i \mid |\{j \in M^i \mid m_j = i\}| = |\{j \in M^i \mid m'_j = i\}|\} \Rightarrow f(m) = f(m')$

Lemma (easy): *the only impartial nomination rules with anonymous ballots are the constant rules*

eschewing the impossibility: restrict the legitimate ballots $M_i \subseteq N \setminus \{i\}$

\Rightarrow *positional* nomination rules along a tree

example

order agents by *seniority*

*everyone nominates someone **more senior than himself***

the youngest nominated agent wins

- impartial, monotonic, anonymous ballots
- discriminates against the most junior
- the most senior is a dummy

Figure 2a

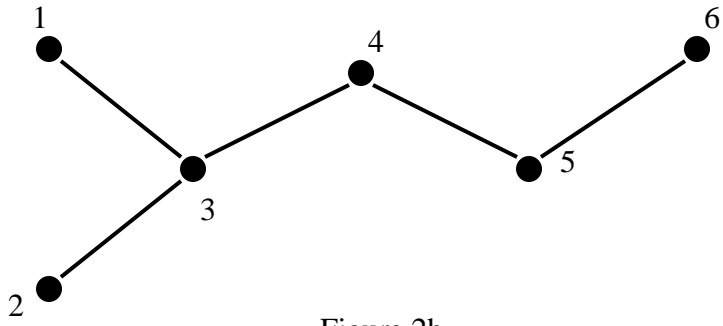
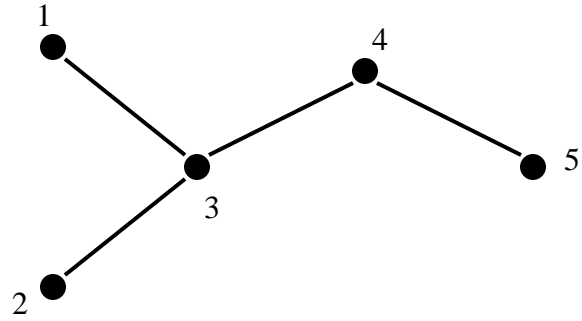


Figure 2b

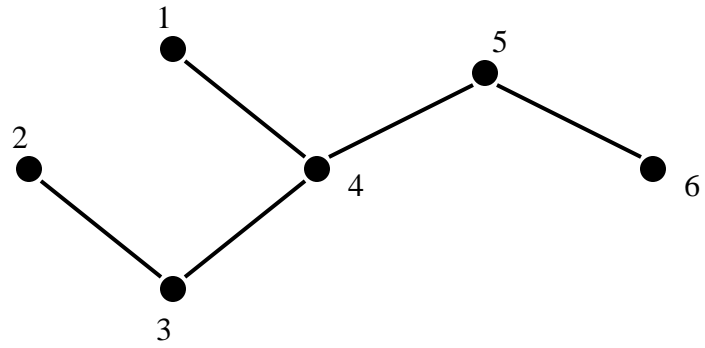


Figure 2c

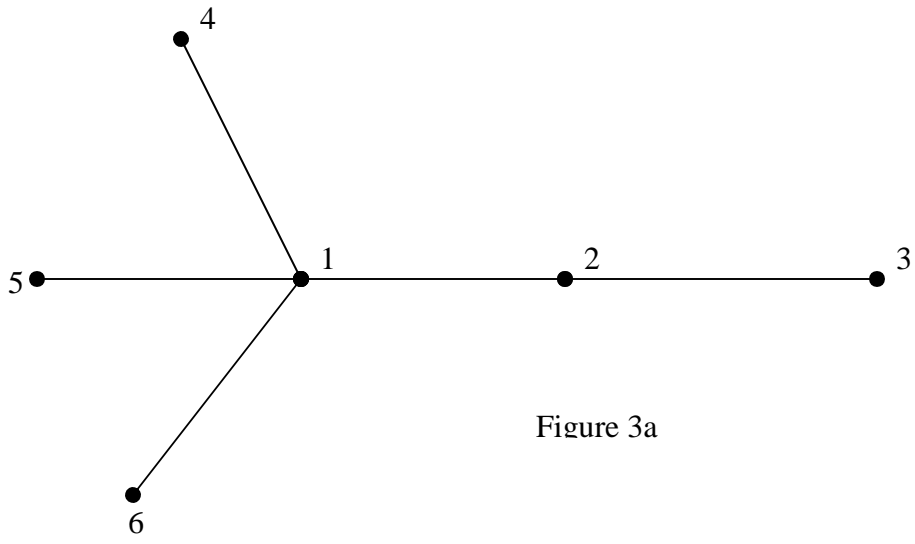


Figure 3a

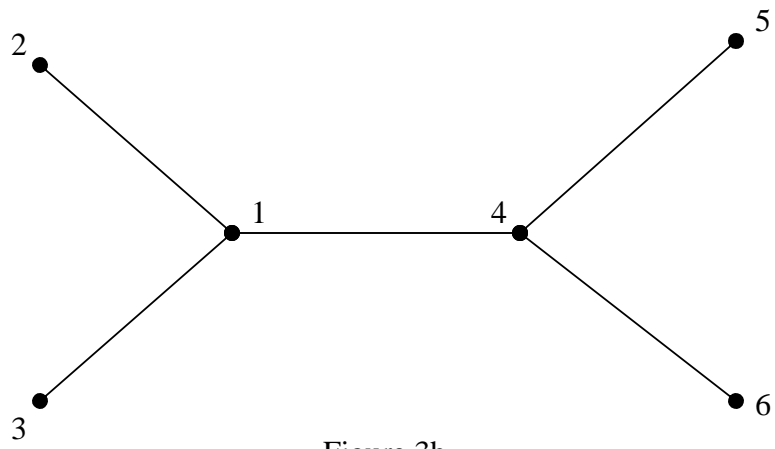


Figure 3b

the family of median nomination rules (n odd, $n \geq 5$)

the agents are the nodes of a tree Γ

Γ is neither a line nor a simple star

i^* is the median node/agent of Γ

M_i is the largest subtree rooted at j adjacent to i , away from i

M_{i^*} is *one of* the largest subtrees at j^* adjacent to i^* , away from i^*

→ *winner: the median vote*

n even: add (carefully) a fixed ballot

Theorem:

The median nomination rule on Γ is impartial, monotonic, unanimous and has anonymous ballots; and i influences $j \Leftrightarrow j \in M_i$

- Unanimity: if all $j \in N \setminus \{i\}$ such that $i \in M_j$ nominate i , then i wins

critique

- unequal influence: tradeoff: maximize $\min |M_i| \leftrightarrow \sum_N |M_i|$
- Negative Unanimity: an agent can win without receiving any nomination

Figure 4a

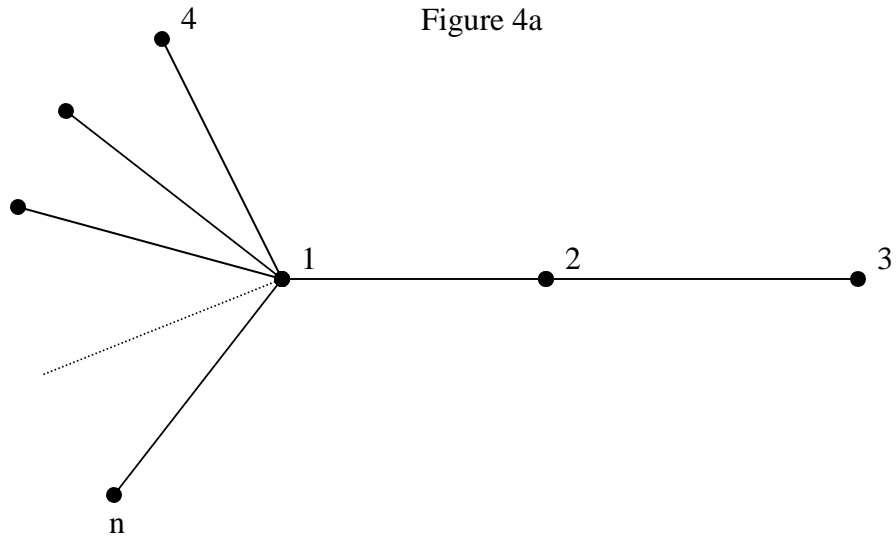
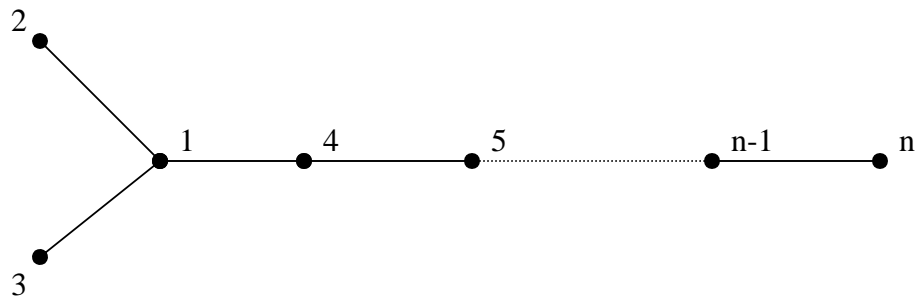


Figure 4b



Conjecture: we cannot construct an impartial, monotonic, onto nomination rule where the winner is always a nominee

voting rules

the most natural messages: $M_i = \mathcal{L}(N \setminus \{i\})$ linear ordering of other agents

- **Monotonicity:** lifting j in i 's ranking does not threaten j 's win
- **Unanimity:** $\{i = \text{top}(m_j) \text{ for all } j \in N \setminus \{i\}\} \Rightarrow i \text{ wins}$

*the family of **partition voting rules** ($n \geq 7$)*

*partition $N = \cup_{k=1}^K N_k$ in **districts** s. t. $|N_1| \geq 4$ and $|N_k| \geq 3$ for $k \geq 2$*

for each k choose a quota rule $I^{\varepsilon_k}(N_k, q_k)$, $\varepsilon_k = ab, r$

choose a default agent i^ in district 1*

two equivalent definitions: direct voting, or two steps voting

Step 1

run $I^{\varepsilon_k}(N_k, q_k)$ in each district $k \geq 2$: call i a local winner if she wins

call i^ a local winner if he wins in $I^{\varepsilon_1}(N_1, q_1)$*

call $i \in N_1 \setminus \{i^\}$ a local winner if she wins without i^* 's support*

$$\text{if } \varepsilon_1 = ab : s_i(N_1 \setminus \{i, i^*\}) \geq q_1$$

$$\text{if } \varepsilon_1 = r : s_i(N \setminus \{i, i^*\}) \geq s_j(N \setminus \{i, j\}) + q_1 \text{ for all } j \in N_1 \setminus \{i\}$$

If there is no local winner anywhere, i^ wins*

if there is a single local winner, she wins; otherwise go to

Step 2 All the non local winners use a standard voting rule to award the prize to one of the local winners.

Theorem

A partition voting rule is impartial, unanimous, and has full mutual influence. If it uses an absolute quota in district 1, or if $|N_1| = 4$, the rule is monotonic.

critique:

- Negative Unanimity: the default i^* can win and be the worst choice for all in $N \setminus \{i^*\}$

partial remedy: under *Impartial Culture* the probability that at least a local winner exists goes to 1 if the district size remains bounded while n increases. The convergence is faster if opinions are correlated!

⇒ the advantage of the default agent vanishes

⇒ the winner must be Pareto optimal

two vague open questions

- what is the *special role* of median rules among anonymous monotonic nomination rules?
- can we find impartial rules *more equitable* than the partition voting rules?

model 2: peer ranking

assign n ranks to n agents

private consumption of one's rank

$i \in N, a \in A$

$\Sigma(N, A) \ni \sigma : \text{bijection } N \rightarrow A$

i 's message $m_i \in M_i$

assignment mechanism: $M_N \ni m \rightarrow \theta(m) \in \Sigma(N, A)$

- **Impartiality:** $\theta(m|{}^i m_i)[i] = \theta(m|{}^i m'_i)[i]$, for all i, m_i, m'_i
- **Full Ranks :** for all $i \in N, a \in A$, for some $m \in M_N : \theta(m)[i] = a$
- **Full Range:** for all $\sigma \in \Sigma(N, A)$ for some $m \in M_N : \sigma = \theta(m)$

Lemma (easy):

For $n = 3$, Impartiality \cap Full Ranks = \emptyset

For $n = 4$, Impartiality \cap Full Ranks $\neq \emptyset$

$M^i = \{0, 1\}$ for all i , $A^* = \{1, 2, 3, 4\}$

$\theta^4(0, 0, 0, 0) = 1234$; $\theta^4(1, 0, 0, 0) = 1432$; $\theta^4(0, 0, 0, 1) = 1324$; $\theta^4(1, 0, 0, 1) = 1$

$\theta^4(0, 0, 1, 0) = 2134$; $\theta^4(0, 1, 1, 0) = 2143$; $\theta^4(0, 0, 1, 1) = 2314$; $\theta^4(0, 1, 1, 1) = 2$

$\theta^4(1, 1, 0, 0) = 3412$; $\theta^4(1, 1, 1, 0) = 3142$; $\theta^4(1, 1, 0, 1) = 3421$; $\theta^4(1, 1, 1, 1) = 3$

$\theta^4(0, 1, 0, 0) = 4213$; $\theta^4(0, 1, 0, 1) = 4321$; $\theta^4(1, 0, 1, 0) = 4132$; $\theta^4(1, 0, 1, 1) = 4$

fairly symmetric treatment of the agents

range is not full (15 assignments)

use $\theta^4 \rightarrow$ **an impartial mechanism with full ranks** for any n divisible by 4

fix a partition $N = N_1 \cup N_2 \cup N_3 \cup N_4$ with $|N_i| = \frac{n}{4}$ and an order τ of A

play θ^4 with agents in N_i jointly playing the 1st coordinate 0 or 1

N_i gets rank/object 1 \Rightarrow the first $|N_i|$ ranks in τ go to N_i ; etc..

agents in $N \setminus N_i$ jointly choose the assignment of these $|N_i|$ ranks inside N_i

construct an **impartial mechanism with full range**

→ *separating family* in $A : \mathcal{S} \subset 2^A$ such that

for all $a, b \in A, a \neq b$, there exists $S \in \mathcal{S} : a \in S, b \notin S$

→ separating family of size k : for all $S \in \mathcal{S} : |S| = k$

Lemma:

For $n = |A| \geq 6$, we can find three **pairwise disjoint** separating families in A , all of identical size.

For $n \leq 5$, we can find at most two such disjoint families.

	\mathcal{S}_1	\mathcal{S}_2	\mathcal{S}_3
$A = \{a, b, c, d, e, f\}$	abc	abd	abe
	bcd	bce	bcf
	cde	cdf	acd
	def	ade	bde
	$ae f$	$be f$	$ce f$
	$ab f$	$ac f$	$ad f$

$|A| \geq 7, A = \{1, 2, \dots, n\} \Rightarrow$ for $1 \leq t < \frac{n}{2}$ $\mathcal{S}_t = \{(a, a + t) | a \in A\}$ are separating and pairwise disjoint

choose three "leaders" agents 1, 2, 3

step 1: the leaders choose impartially three ranks for themselves

key: all assignments of $\{1, 2, 3\}$ to A are in the range

step 2: the leaders choose $i \in N \setminus \{1, 2, 3\}$ and assign her one of the free ranks;

agent i chooses $j \in N \setminus \{1, 2, 3, i\}$ and assign him one of the free ranks;

etc...

step 1 explained:

choose three separating families $\mathcal{S}_i, i = 1, 2, 3$, of identical size, pairwise disjoint

each leader chooses $S_i \in \mathcal{S}_i$; given $(S_1, S_2, S_3) \in \mathcal{S}_1 \times \mathcal{S}_2 \times \mathcal{S}_3$

assign 1 to a rank in $S_3 \cap S_2^c \neq \emptyset$

assign 2 to a rank in $S_1 \cap S_3^c \neq \emptyset$

assign 3 to a rank in $S_2 \cap S_1^c \neq \emptyset$

break ties in $S_3 \cap S_2^c$ by an onto vote of leaders 2 and 3

break ties in $S_1 \cap S_3^c$ by an onto vote of leaders 1 and 3

break ties in $S_2 \cap S_1^c$ by an onto vote of leaders 1 and 2

- many variants in step 2
- critique: the three leaders influence the rest of the agents, but not vice versa

Mutual Influence:

$$\forall i, j \in N \exists m_i, m'_i \in M^i, m_{-i} \in M^{N \setminus i} : \theta(m|_i m_i)[j] \neq \theta(m|_i m'_i)[j]$$

we can find an impartial assignment mechanism with full range, satisfying Mutual Influence

its definition is more complex

Open question: in the ranking interpretation (as opposed to assignment), the natural message space is $M_i = \mathcal{L}(N \setminus \{i\})$. Can we achieve the same properties in that format? and Unanimity?