

Datatypes in Lambda Calculus

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1 Booleans

1.1 Definition

```
T ≡ λ t f . t
F ≡ λ t f . f
```

1.2 Not

```
not ≡ λ b . λ t f . b f t
not ≡ λ b . b (λ t f . f) (λ t f . t)
not ≡ λ b . b F T
¬a ↔ not a
```

1.3 And

```
and ≡ λ b1 b2 . b1 b2 F
a & b ↔ and a b
```

1.4 Or

```
or ≡ λ b1 b2 . b1 T b2
a | b ↔ and a b
```

2 Numbers

2.1 Definition

```
nat ::= 0
      | Succ nat
```

```
0 ≡ λ fs vz . vz
succ ≡ λ n . λ fs vz . fs (n fs vz)
```

Church numerals

```
0 ≡ λ s z . z
1 ≡ λ s z . s z
2 ≡ λ s z . s (s z)
...
```

2.2 Is zero?

```
is_zero? ≡ λ n . n (λ α . F) T
```

2.3 Plus

```
plus ≡ λ m n . m succ n
a + b ~> plus a b
```

2.4 Times

```
times ≡ λ m n . m (plus n) 0
a * b ~> times a b
```

3 Pairs

3.1 Definition

```
pair ::= P * *
```

```
mk_pair ≡ λ a b . λ f_p . f_p a b  
<a, b> ~⇒ mk_pair a b  
π1 ≡ λ p . p (λ a b . a)  
π2 ≡ λ p . p (λ a b . b)
```

3.2 Swap

```
swap ≡ λ p . <π2 p, π1 p>  
swap ≡ λ p . mk_pair (π2 p) (π1 p)
```

3.3 Sliding Window

$$\langle a, b \rangle \xrightarrow{f} \langle b, a + b \rangle$$

```
fib ≡ λ n . π1 (n f <0, 1>)
```

$$\langle a, b \rangle \xrightarrow{f} \langle \text{succ } a, a * b \rangle$$

```
fac ≡ λ n . π2 (n f <1, 1>)
```

$$\langle a, b \rangle \xrightarrow{f} \langle b, a \rangle$$

```
is_even? ≡ λ n . π1 (n f <T, F>)
```

4 Lists

4.1 Definition

```
list ::= Nil
      | Cons * list
```

```
mk_cons ≡ λ hd tl . λ fc vn . fc hd (tl fc vn)
hd :: tl ~> mk_cons hd tl
nil ≡ λ fc vn . vn
```

4.2 Sum

```
sum_list ≡ λ xs . xs plus 0
```

4.3 Append

```
append ≡ λ xs ys . xs mk_cons ys
```

5 Trees

5.1 Definition

```
tree ::= Leaf *
      | Node tree tree
```

```
mk_node ≡ λ t1 t2 . λ fn fl . fn (t1 fn fl) (t2 fn fl)
mk_leaf ≡ λ n . λ fn fl . fl n
```

5.2 Sum

```
sum_tree ≡ λ t . t plus id
```

6 Exercises

1. Define the binary xor function, which returns `T` exactly when one of its arguments is `T`.
2. Prove De Morgan's laws by verifying that for all boolean terms `b1`, `b2` it holds that:

$$\begin{aligned} \text{not } (\text{and } b1 \ b2) &= \text{or } (\text{not } b1) \ (\text{not } b2) \\ \text{not } (\text{or } b1 \ b2) &= \text{and } (\text{not } b1) \ (\text{not } b2) \end{aligned}$$

3. We extend boolean logic with another value called *unknown*. For obvious reasons this is called three-valued logic. Create an ML-like datatype describing this type, and derive lambda calculus terms for the constructors of this new datatype (e.g `T`, `F` and `U`).
4. The operations `not`, `and` and `or` for three-valued logic can be described by the following truth tables:

A	¬A
F	T
U	U
T	F

A	B	A ∧ B
F	F	F
F	U	F
F	T	F
U	F	F
U	U	U
U	T	U
T	F	F
T	U	U
T	T	T

A	B	A ∨ B
F	F	F
F	U	U
F	T	T
U	F	U
U	U	U
U	T	T
T	F	T
T	U	T
T	T	T

Implement these functions as lambda calculus terms. Optionally, prove that the De Morgan identities are still satisfied.

5. Implement a function `power`, which computes n^m when given natural numbers `n` and `m` as church numerals.

6. Implement the function $\uparrow\uparrow$, which computes

$$\underbrace{n^{n^{\dots^n}}}_{m \text{ times}}$$

This notation is due to the great mathematician and computer scientist Donald Knuth, who also wrote n^m as $n \uparrow m$, because of the difficulties of using the n^m -notation to describe continued exponentiation.

7. Implement a function that returns $\lfloor \frac{n}{2} \rfloor$ when given a natural number n encoded as a church numeral.

Hint: Combine what we have learned from pairs, booleans and sliding window.

8. Implement a function `car` which returns the first element of a given list. You can assume the list is not empty.

Similarly implement a function `cdr`, which drops the first element of the list.

9. Implement a function `list→pair` which converts a list `xs` into nested pairs representing the same list. You can assume that $|xs| \geq 2$.

10. Implement a function `flatten`, which returns a list of all the elements in a tree.

11. Implement a function `mirror`, which flips the nodes in a tree. Example:

```
mirror (mk_node
        (mk_node
         (mk_leaf 0)
         (mk_leaf 1))
        (mk_leaf 2)) =
(mk_node
 (mk_leaf 2)
 (mk_node
  (mk_leaf 1)
  (mk_leaf 0)))
```

12. The `Option` datatype is a well-known concept from functional programming:

```
Option ::= None
         | Some *
```

Implement this datatype as lambda terms, along with predicates `is_none?` (resp. `is_some?`), which return `T` if its argument has been constructed using `None` (resp. `Some`).

13. Implement a function `maybe`, that when given a default argument `d`, a function `f` and an option value `o`, returns `d` if `is_none? o = T`. Alternatively when `o = Some v` the function should return `f v`.
14. Implement a function `lift`, that when given a function `f: A → B`, returns a function `g: Option A → Option B`.

More concretely, `None` \mapsto^g `None` and `Some v` \mapsto^g `Some (f v)`.