

Purely Functional Worst Case Constant Time Catenable Sorted Lists

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Joint work with

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Catenable Sorted Lists

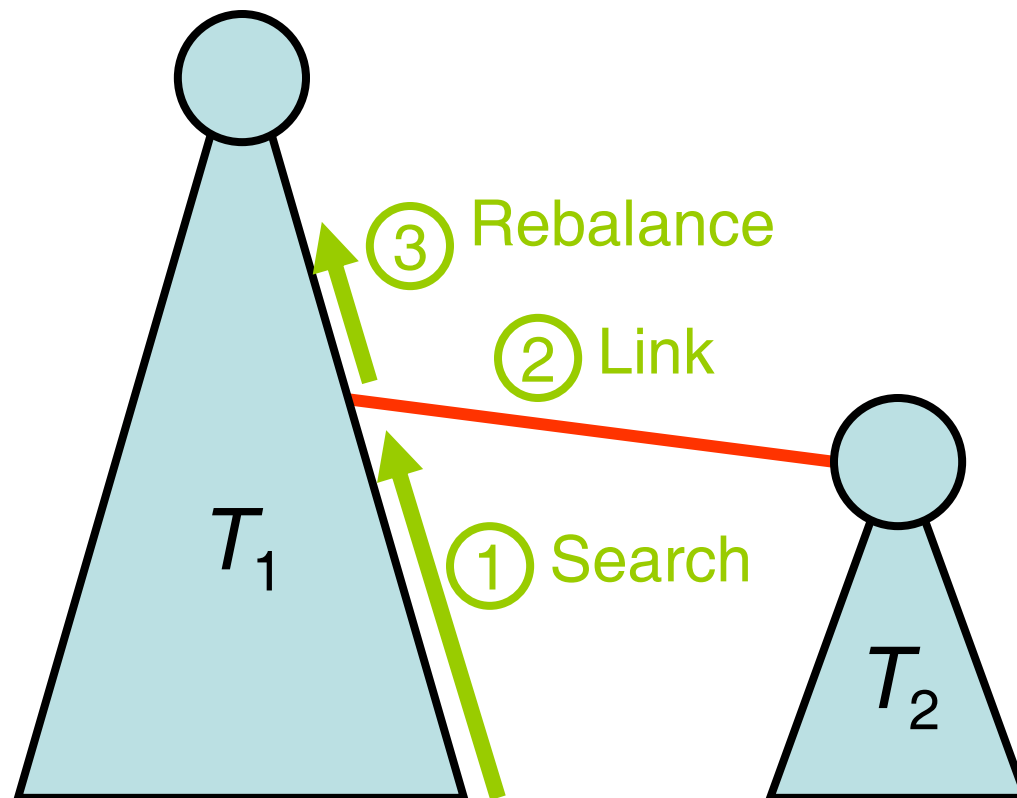
- Insert(T, x)
- Delete(T, x)
- Search(T, x)
- Join(T_1, T_2)
- Split(T, x)

This talk

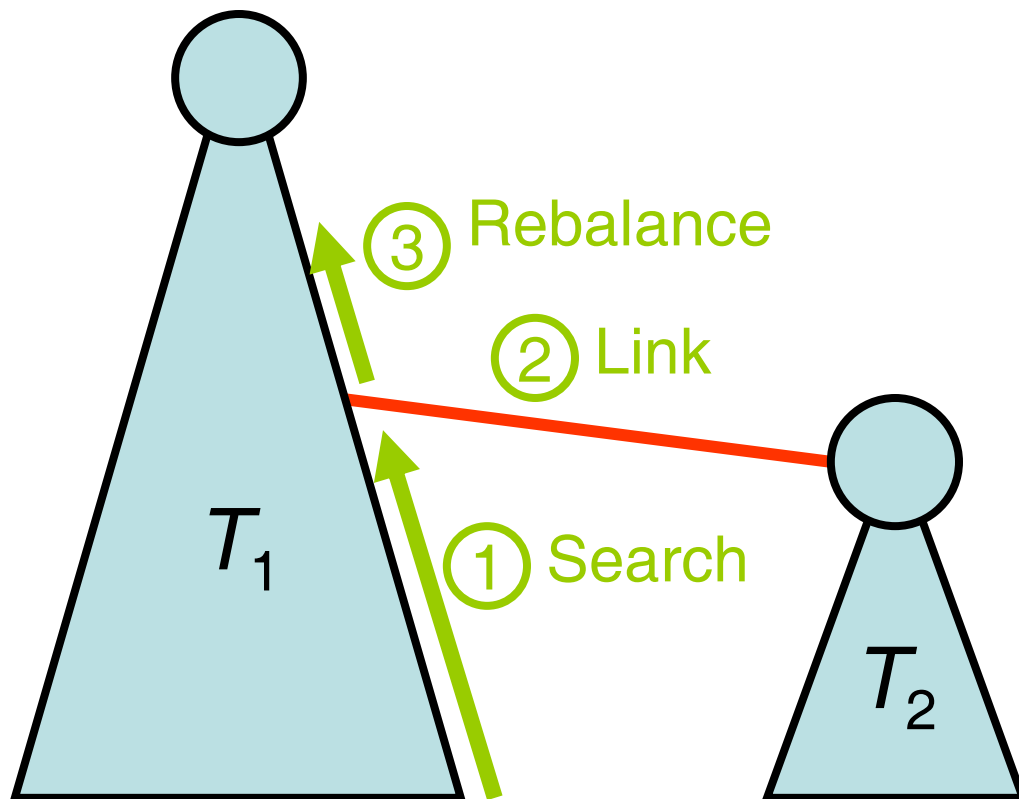
$$T_1 = 2, 5, 7, 8, 10, 11 \quad T_2 = 13, 17, 19, 34, 58, 79$$

$$\text{Join}(T_1, T_2) = 2, 5, 7, 8, 10, 11, 13, 17, 19, 34, 58, 79$$

Catenable Sorted Lists: Search Trees (2,4-trees, red-black trees...)



Catenable Sorted Lists: Search Trees (2,4-trees, red-black trees...)



- ① Worst-case $O(\log |T_2|)$
or $O(\log \log |T_2|)$
or **amortized** $O(1)$
- ② Worst-case $O(1)$
- ③ Worst-case $O(\log |T_1|)$
or **amortized** $O(1)$
or worst-case $O(1)$

Catenable Sorted Lists

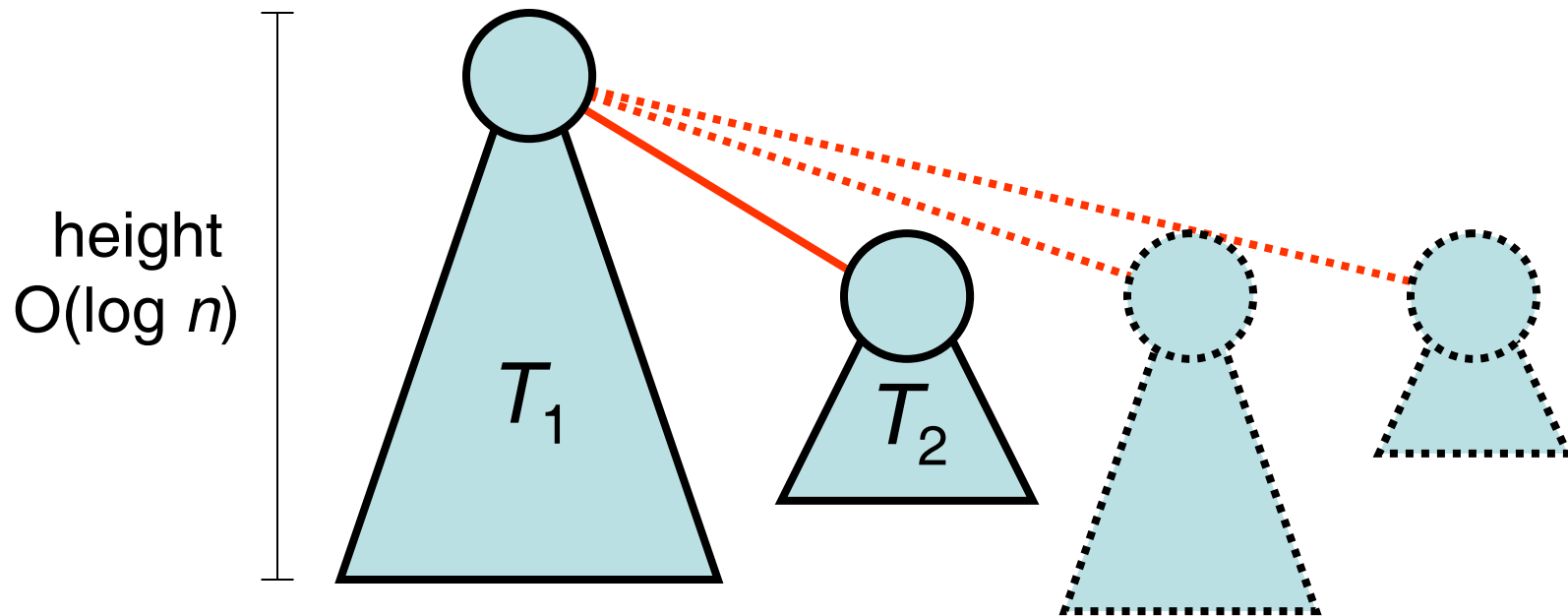
	Search trees	Kaplan Tarjan'96	This talk
Search	$\log n$	$\log n$	$\log n$
Insert/Delete	$\log n$	$\log n$	$\log n$
Join	$\log n$ 1*	$\log \log n_s$	1
Split	$\log n$	$\log n_s$	-

$n = |T|$ $n_s = \min(|T_1|, |T_2|)$ * amortized

$O(1)$ Join and $O(\log n)$ Search

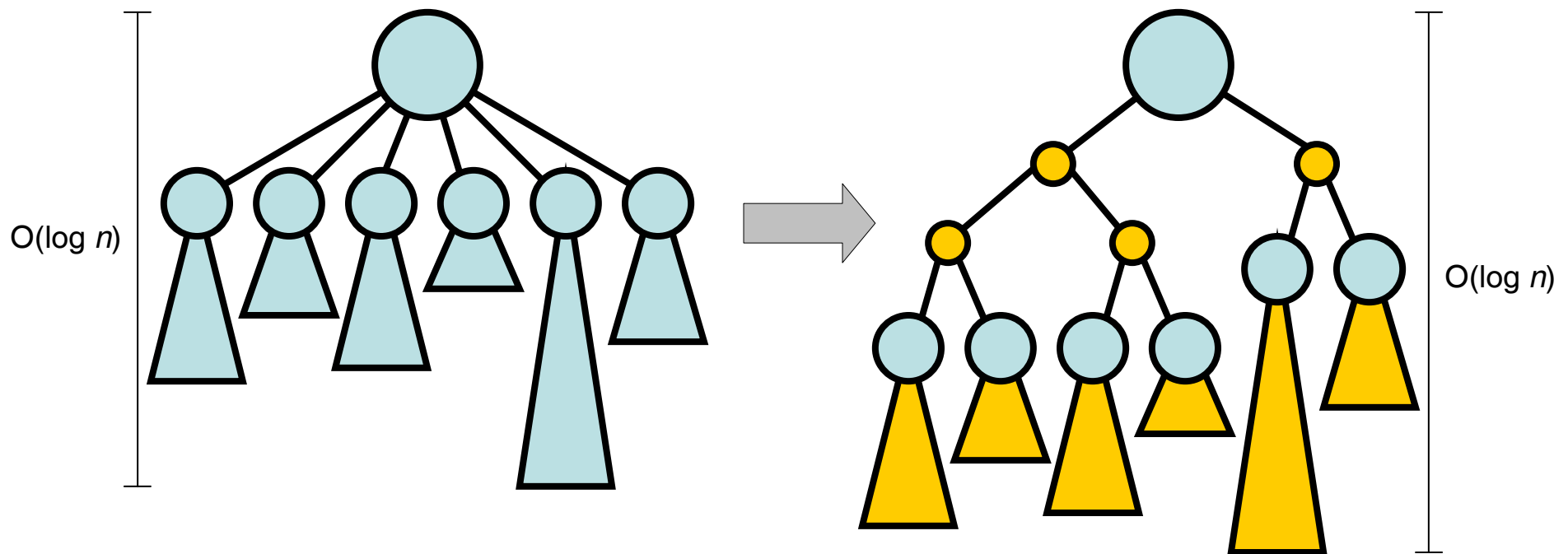
The ideas....

Idea 1: Linking By Size



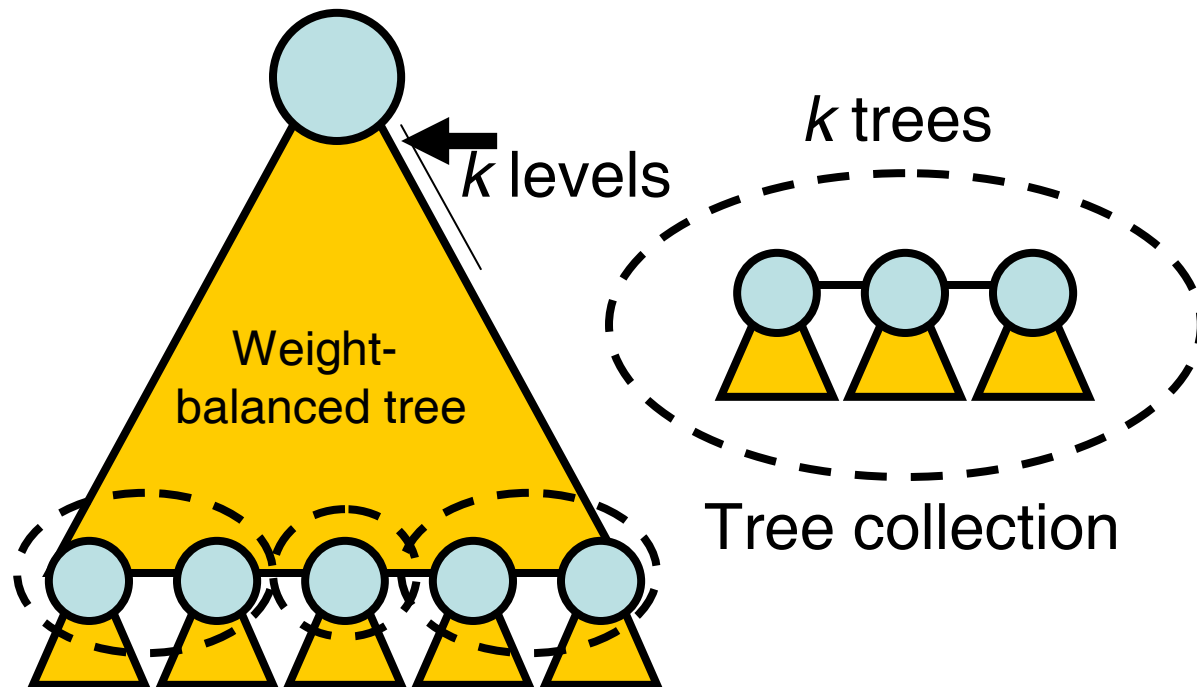
Problem: Nodes of arbitrary degree

Idea 2: Represent Nodes by Weight-Balanced Trees

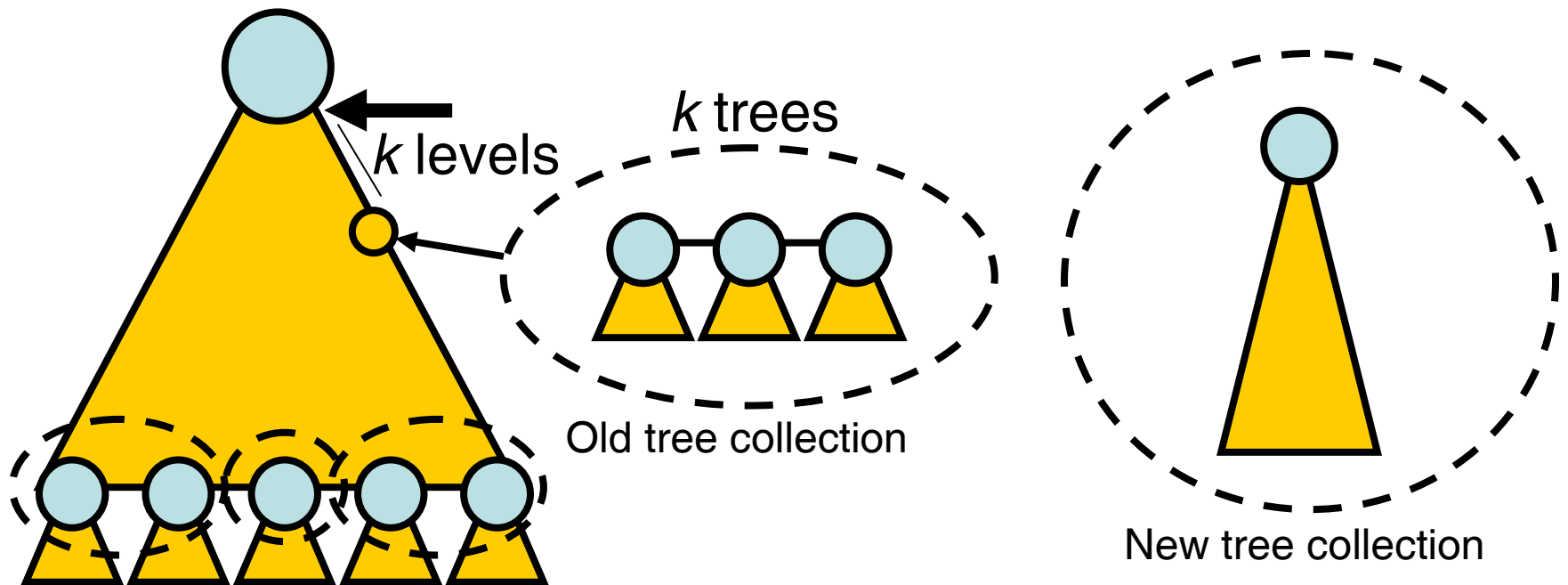


Problem: Does not support linking by size in $O(1)$ time

Idea 3: Tree-Collections and Lazy Join



Idea 3 (*cont.*): Lazy Join

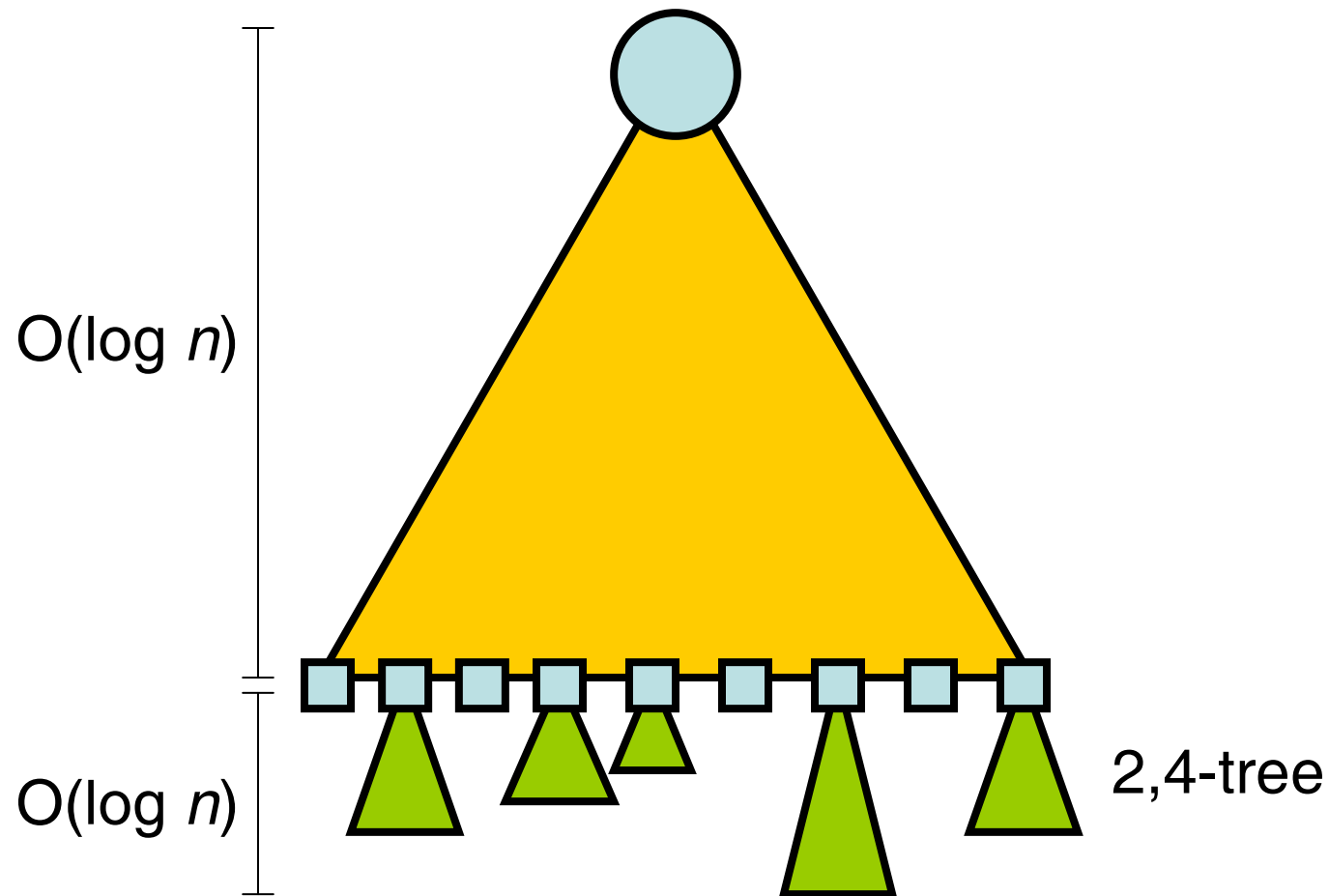


Weight balanced trees with top-down rebalancing
[Bent, Sleator, Tarjan 1985]

Catenable Sorted Lists

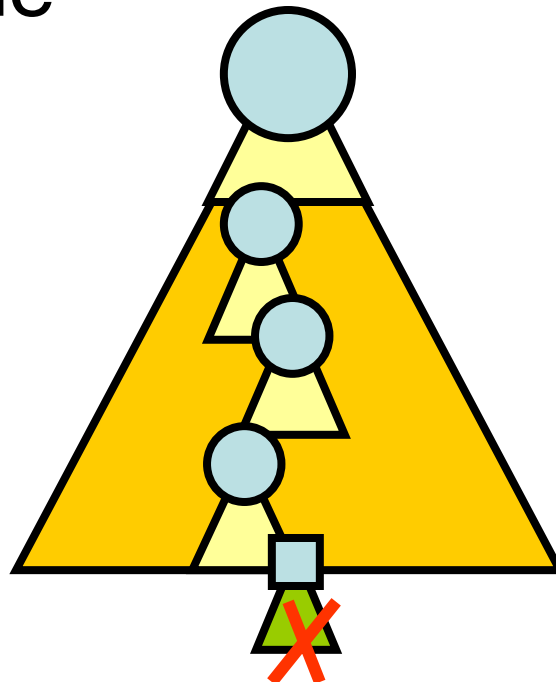
Operation	Worst-case time
Search	$O(\log n)$
Join	$O(1)$

Supporting Insertions: Buckets



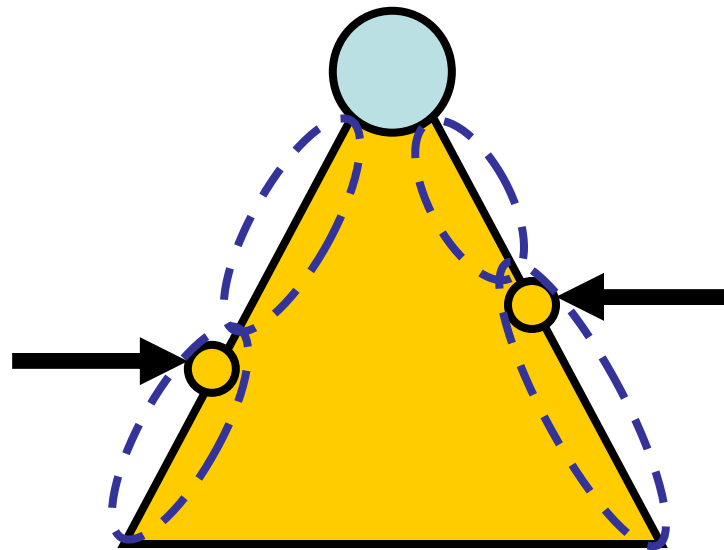
Supporting Deletions: The ideas...

- Delete leaf
- Rebalance weight-balanced trees on root-to-leaf path
- Join $O(1)$ weight-balanced trees in $O(\log n)$ time




Functional Catenable Sorte Lists

- Functional setting = no side effects allowed
- Data structures are automatically **persistent**
- **Idea**: Represent spines by functional catenable lists [Kaplan, Tarjan 1995]



Conclusion

	Search trees	This talk 
Search	$\log n$	$\log n$
Insert/Delete	$\log n$	$\log n$
Join	$\log n$ 1*	1
Split	$\log n$	-

* amortized

 purely functional