

Algorithms and Data Structures

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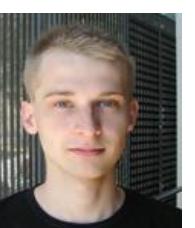
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Strict Fibonacci Heaps

Gerth Stølting Brodal

Aarhus University



George Lagogiannis

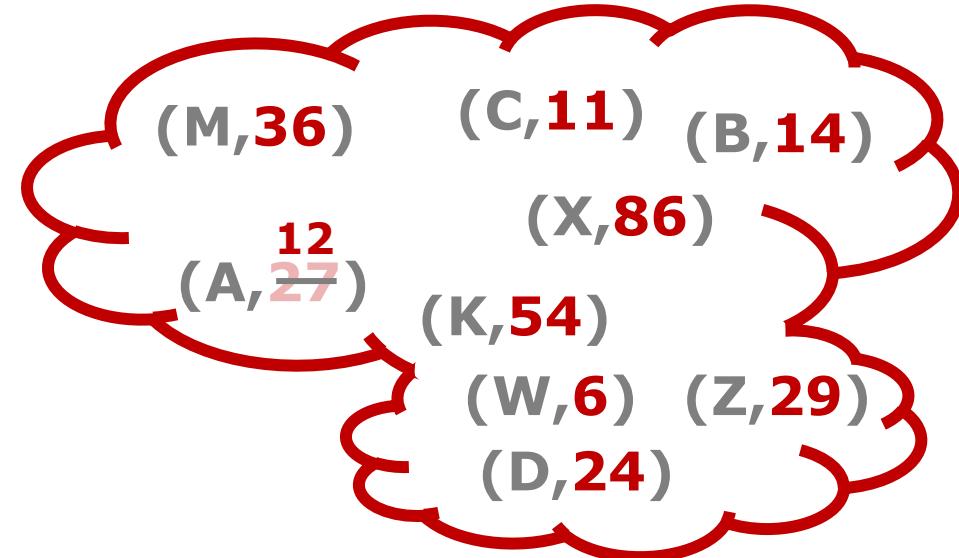
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The Problem – Priority Queues

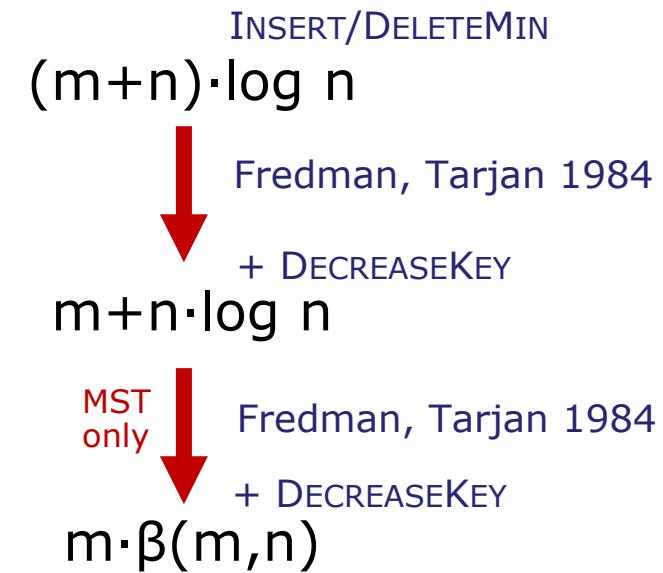
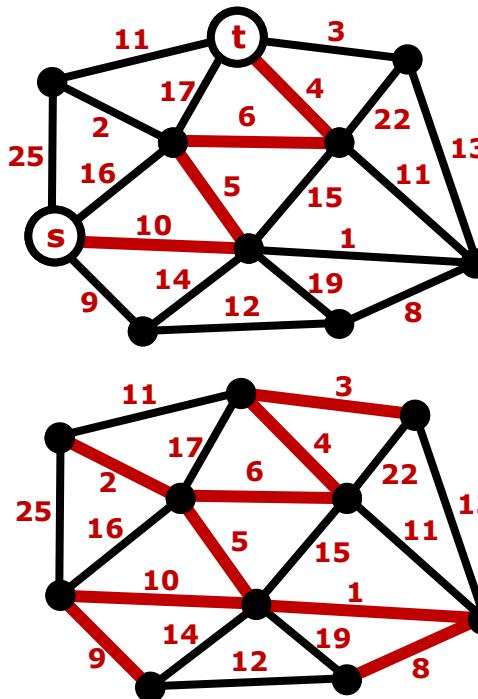
- $\text{INSERT}(\text{value}, \text{key})$
- FINDMIN
- $\text{DELETEMIN} / \text{DELETE}(\&\text{value})$
- $\text{MELD}(Q_1, Q_2)$
- $\text{DECREASEKEY}(\&\text{value}, \Delta)$



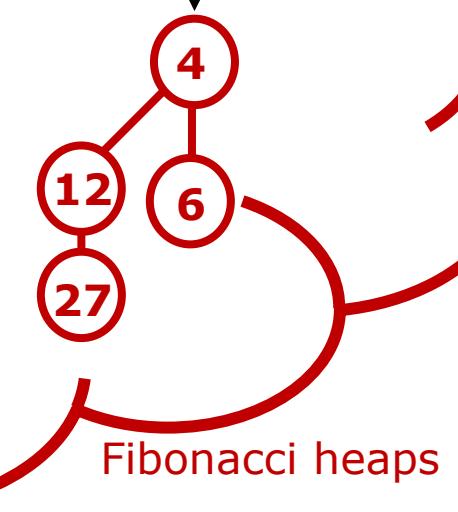
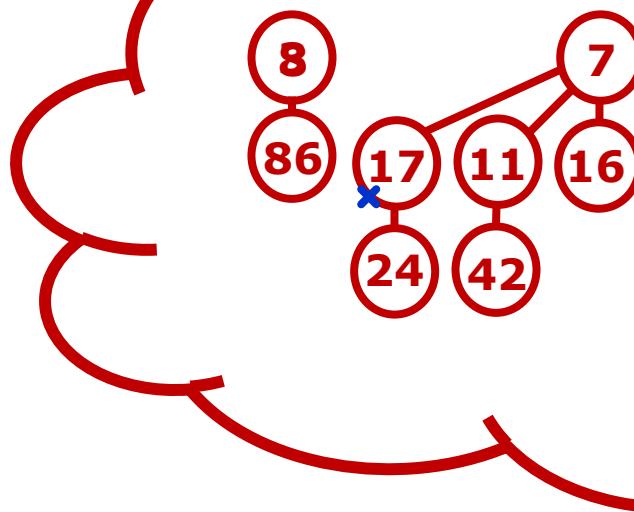
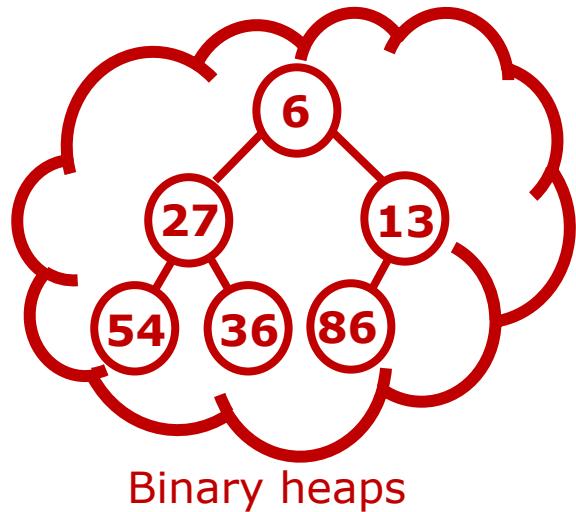
Applications

- **Shortest Path**
Dijkstra (1956)
- **Minimum Spanning Tree**
Borůvka (1926)
Jarník (1930)

(n node, m edges)



History

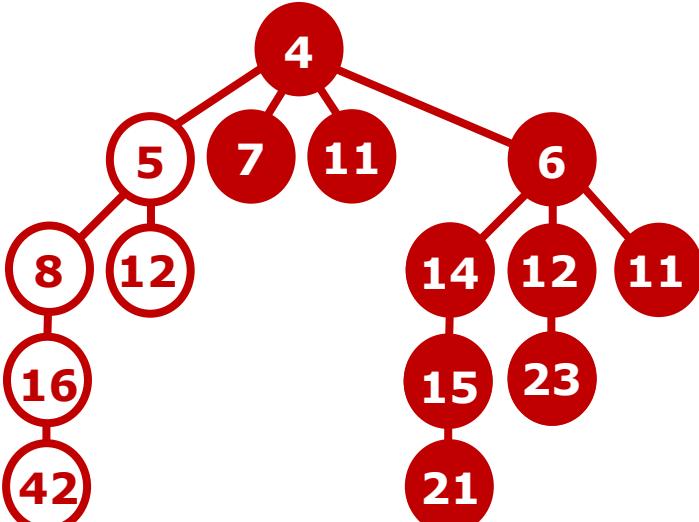
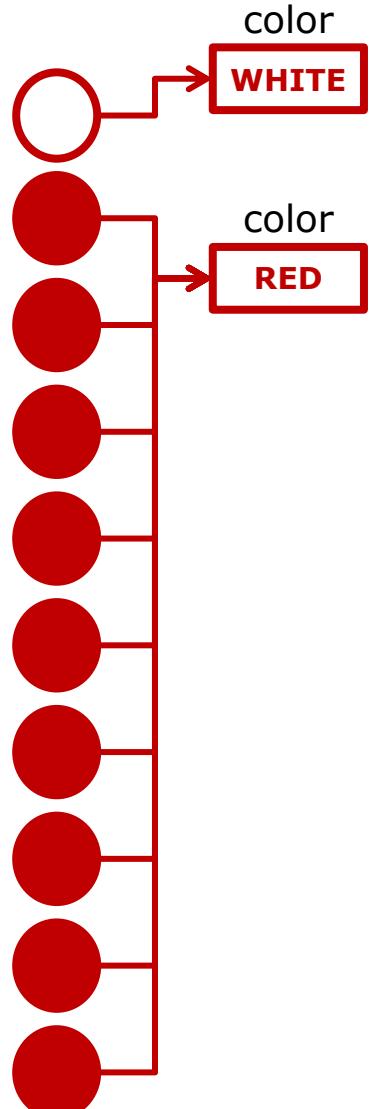


	Binary heaps	Binomial queues	Fibonacci heaps	Run-relaxed heaps	one tree		
	Williams 1964	Vuillemin 1978	Fredman Tarjan 1984	Driscoll et al. 1988	Brodal 1995	Brodal 1996	STOC 2012
Insert	$\log n$	$\log n$	1	1	1	1	1
FindMin	1	1	1	1	1	1	1
Delete	$\log n$	$\log n$	$\log n$	n	$\log n$	$\log n$	$\log n$
Meld	-	$\log n$	1	1	$\log n$	1	1
DecreaseKey	$\log n$	$\log n$	$\log n$	n	1	$\log n$	1

Amortized complexity (Tarjan 1983)

Arrays Pointers only

Ideas



MELD \Rightarrow smallest tree **red** + link

Invariants

1. **white** nodes share one color record
2. **white** children left of **red**
3. root is **red**

Intuition

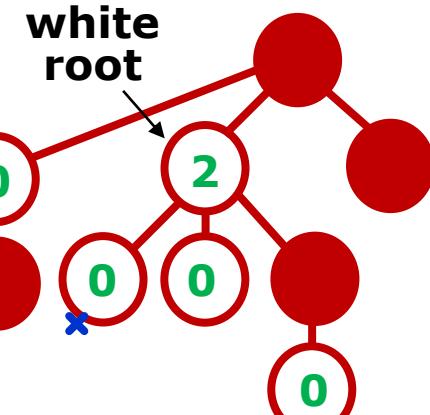
1. Only maintain structure for **white** nodes
2. **Red** non-root nodes never increase degree
3. DELETE: $O(1)$ **red** nodes \rightarrow **white**

Definitions

1. **white** node **rank** = #**white** children
2. DECREASEKEY: cut + **mark** parent (if white)

Invariants

1. i'th rightmost **white** child of a **white** node:
$$\text{rank} + \#\text{marks} \geq i - 1$$
2. $\#\text{marks} + \#\text{white roots} = O(\log n)$

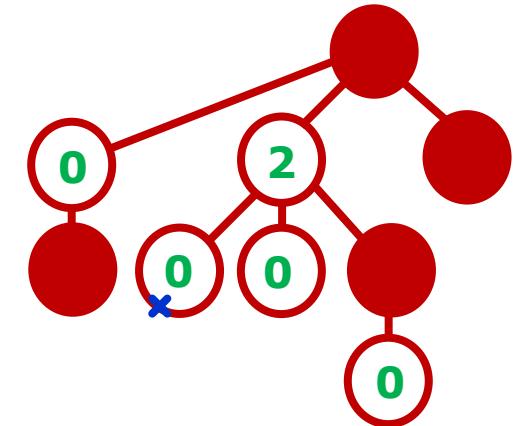


Theorem

max rank
 $O(\log n)$

Priority Queue Operations

- FINDMIN = return root
- INSERT = make single **red** node + MELD
- DELETE = DECREASEKEY to $-\infty$ + DELETEMIN
- MELD = color smaller tree **red** + link + O(1) transformations
root degree +1
- DECREASEKEY = cut + link with root + O(1) transformations
root degree +1; white root (+1); marks (+1)
- DELETEMIN = cut root + find new root + O(log n) link
root degree +O(log n); white root +O(log n)
+ O(log n) transformations



Invariant

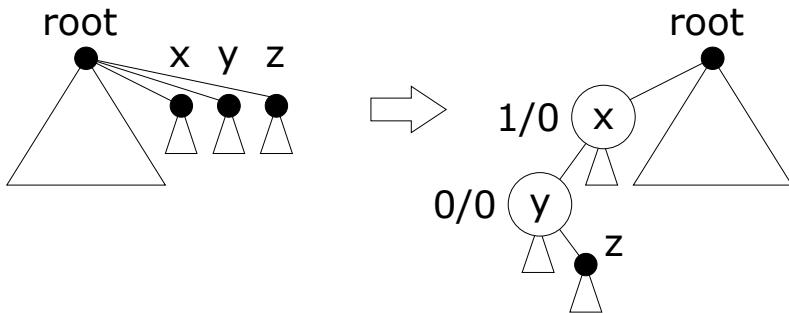
$$R = \alpha \cdot \log (3/2 \cdot \# \text{white nodes} + \# \text{red nodes}) + \beta$$

- degree **unmarked white** nodes $\leq R$
- degree **red** and **marked white** nodes $\leq R - 1$

Transformations

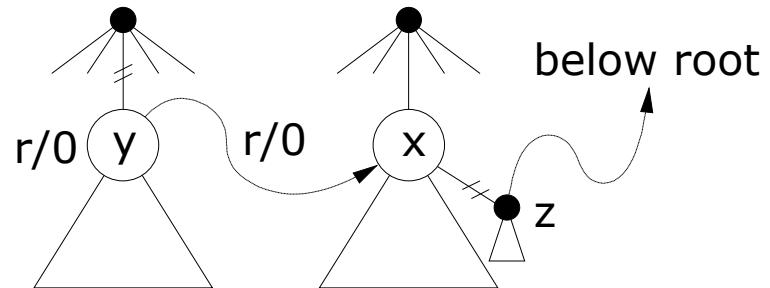
Root degree reduction

Converts two red nodes to white; reduces the root degree by two; creates one new white root



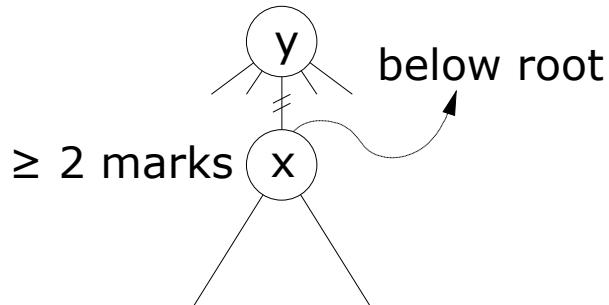
White root reduction

Two white roots of rank r is replaced by one rank $r+1$; increases root degree by one



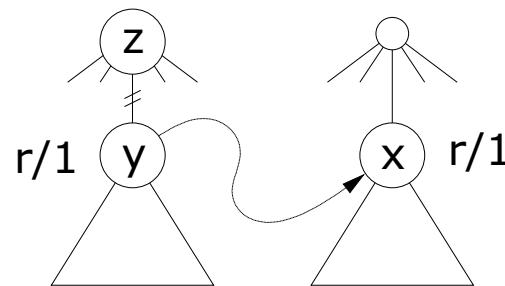
One node mark reduction

White node with ≥ 2 marks becomes a white root (unmarked); increases the root degree by one



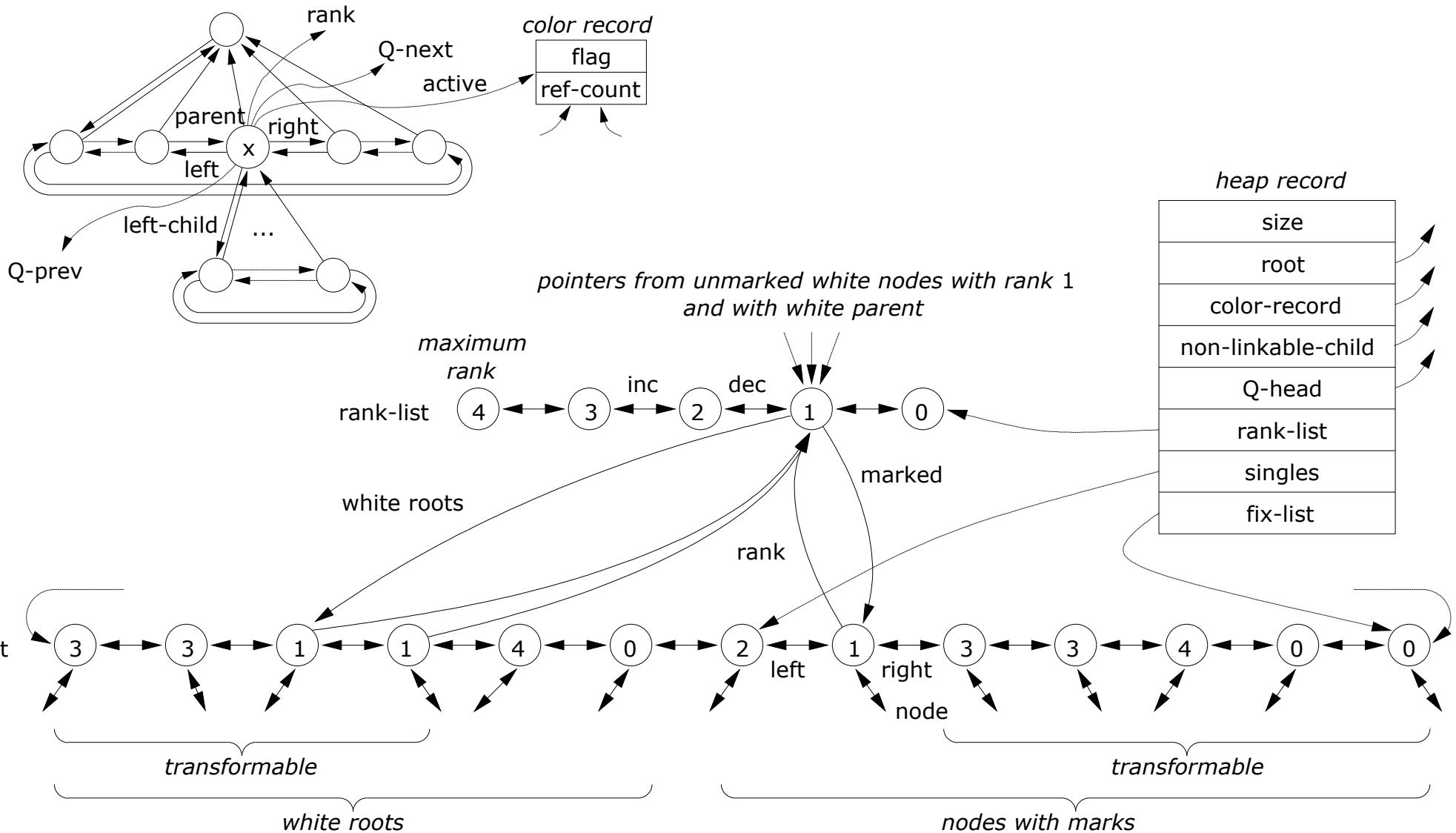
Two node mark reduction

Two nodes of equal rank r with 1 mark become unmarked; one parent one more mark



$$r / m = \text{rang} / \#\text{marks}$$

Representation



	Williams 1964	Vuillemin 1978	Fredman Tarjan 1984	Driscoll et al. 1988	Brodal 1995	Brodal 1996	STOC 2012	
Insert	$\log n$	$\log n$	1	1	1	1	1	1
FindMin	1	1	1	1	1	1	1	1
Delete	$\log n$	$\log n$	$\log n$	n	$\log n$	$\log n$	$\log n$	$\log n$
Meld	-	$\log n$	1	1	1	$\log n$	1	1
DecreaseKey	$\log n$	$\log n$	$\log n$	n	1	1	$\log n$	1

Pointers only

Thank You