

# Exact and Approximate Algorithms for Finding $k$ -Shortest Paths with Limited Overlap

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## MOTIVATION

Finding multiple short yet different routes between two locations in a road network is a problem with various real-world applications:

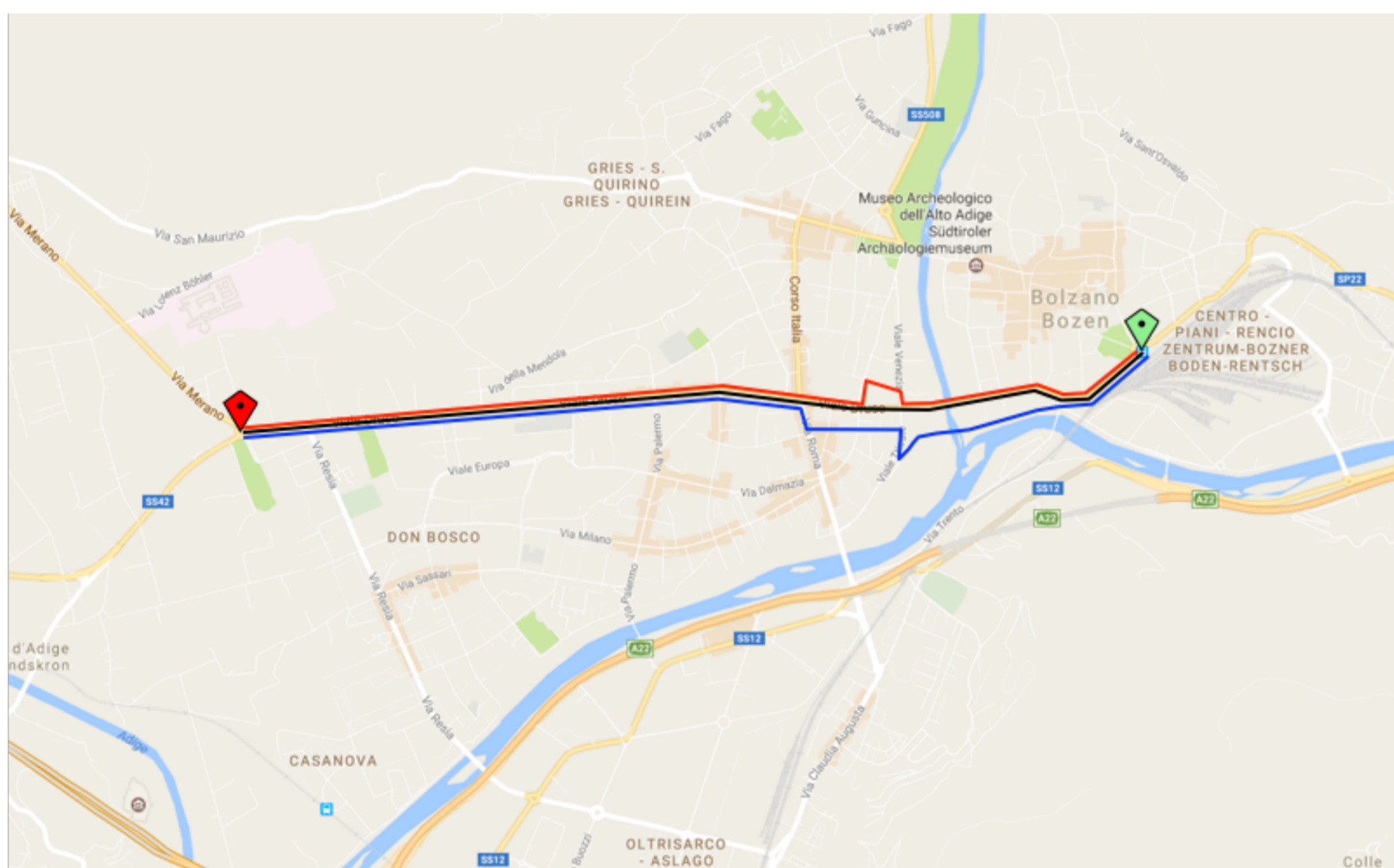
- ✓ Commercial Route Planners
- ✓ Evacuation planning
- ✓ Humanitarian aid

## $k$ -SPwLO PROBLEM

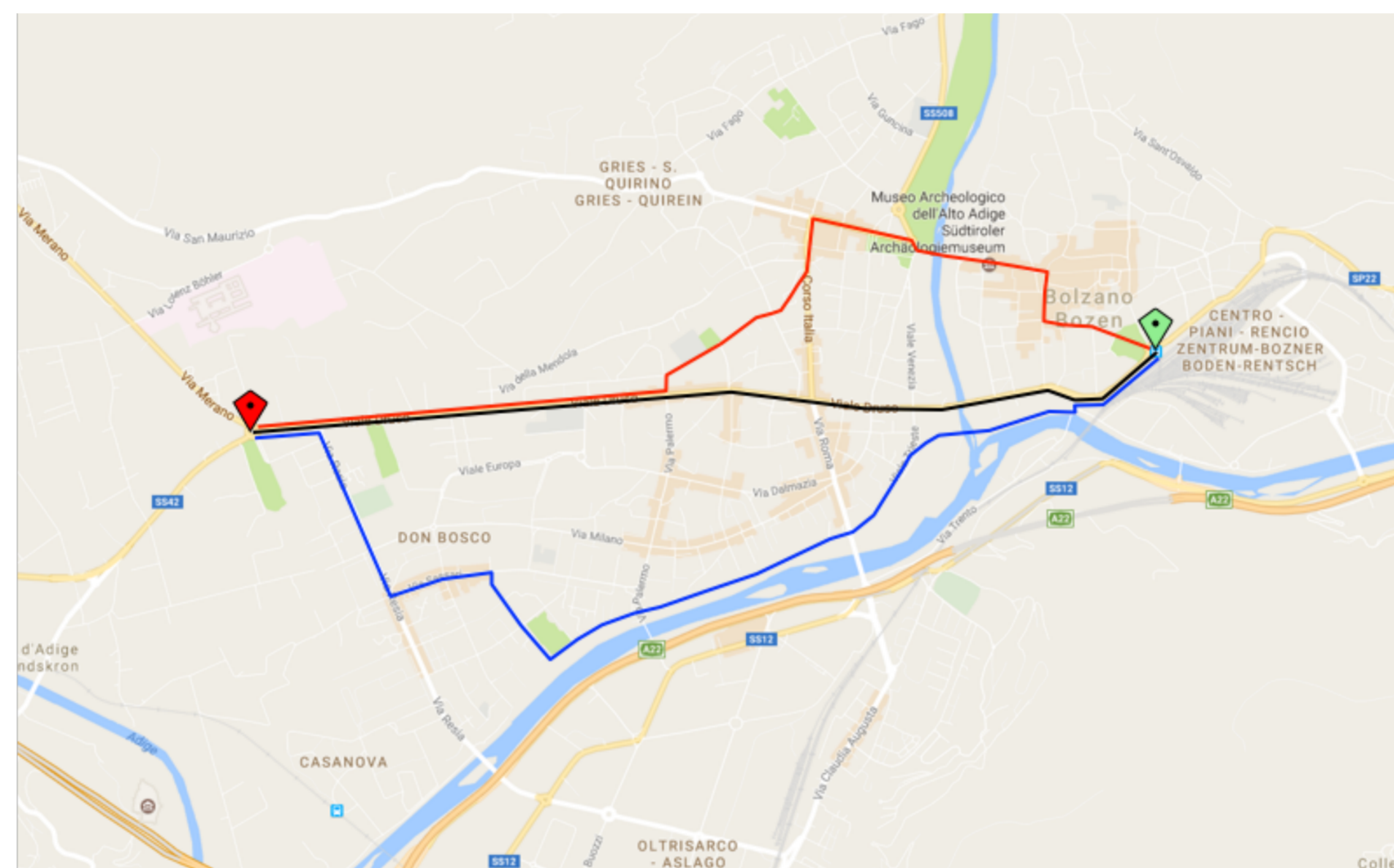
Given a source  $s$  and a target  $t$ , the  $k$ -SPwLO is a set of  $k$  paths from  $s$  to  $t$ , sorted by length in increasing order, such that:

- (a) the set includes the shortest path  $p_0(s \rightarrow t)$ ,
- (b) every path is dissimilar to its predecessors w.r.t. a similarity threshold  $\theta$ ,
- (c) all  $k$  paths are as short as possible.

## EXAMPLE



K-SP: Short but very similar paths



$k$ -SPwLO: Longer but dissimilar paths

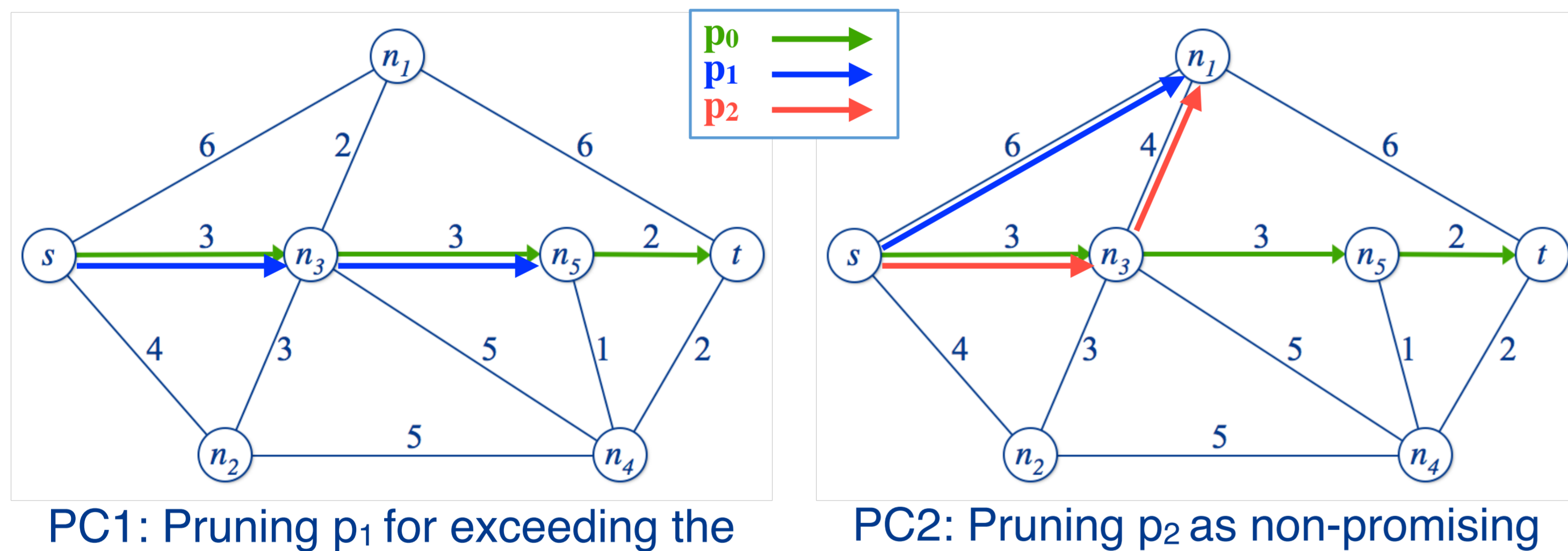
Result Set 1 - K-Shortest Paths			
$p$	$\ell$	$Sim(p, p_0)$	$Sim(p, p_1)$
Shortest path	3.4 km	-	-
1st alternative	3.6 km	95%	-
2nd alternative	3.7 km	80%	76%

Result Set 2 - $k$ SPwLO ( $\theta = 50\%$ )			
$p$	$\ell$	$Sim(p, p_0)$	$Sim(p, p_1)$
Shortest path	3.4 km	-	-
1st alternative	4.1 km	47%	-
2nd alternative	4.5 km	20%	6%

## EXACT ALGORITHM

### MultiPass

- ✓ Traverses the road network  $k-1$  times
- ✓ Restarts the expansion after each alternative path is found
- ✓ Employs two powerful pruning criteria



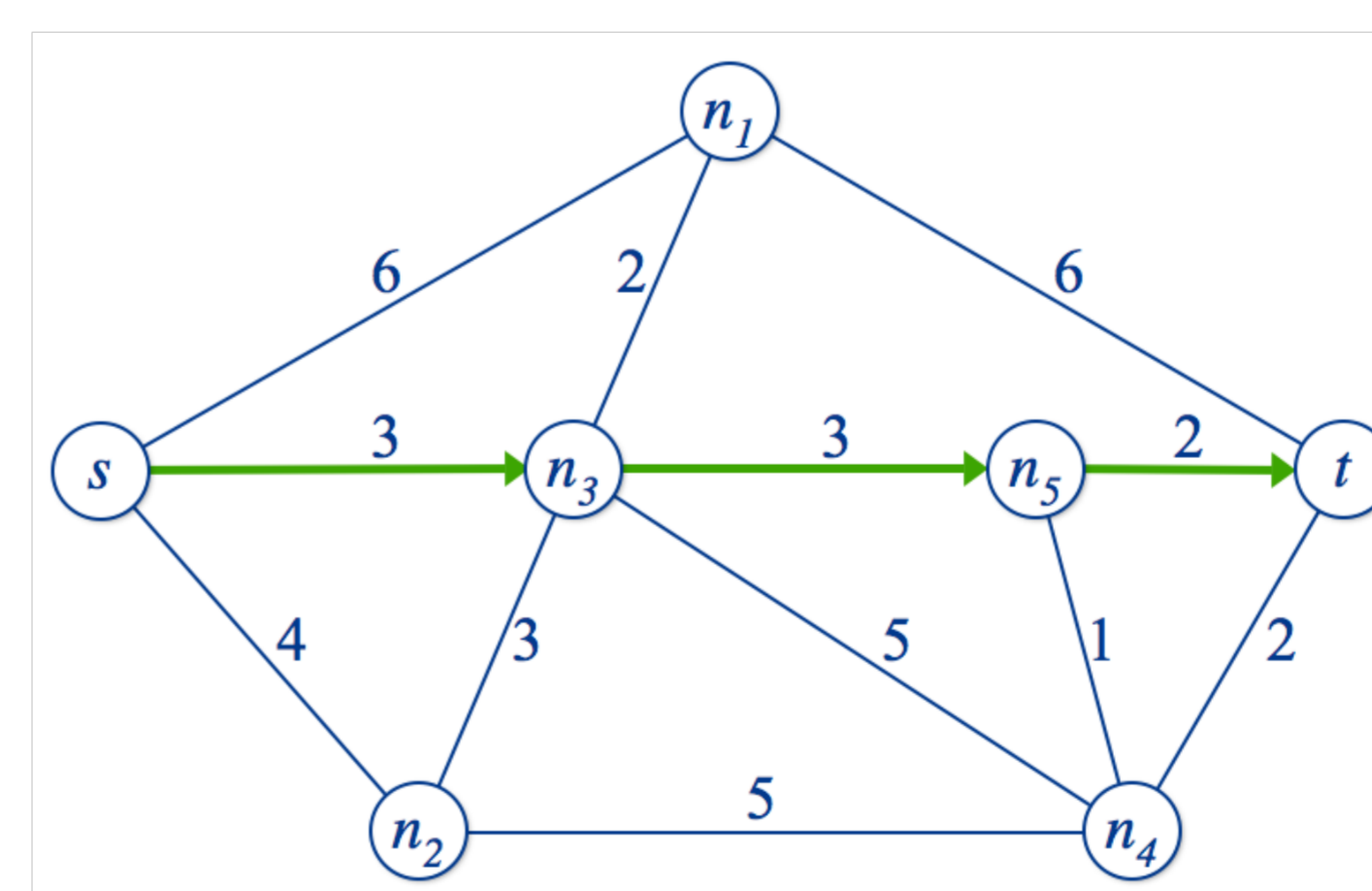
## APPROXIMATE ALGORITHMS

### OnePass+

- ✓ Traverses the road network once
- ✓ Prunes paths with both PC1 and PC2
- ✓ Does not guarantee that the exact solution will be found

### ESX (Edge Subset Exclusion)

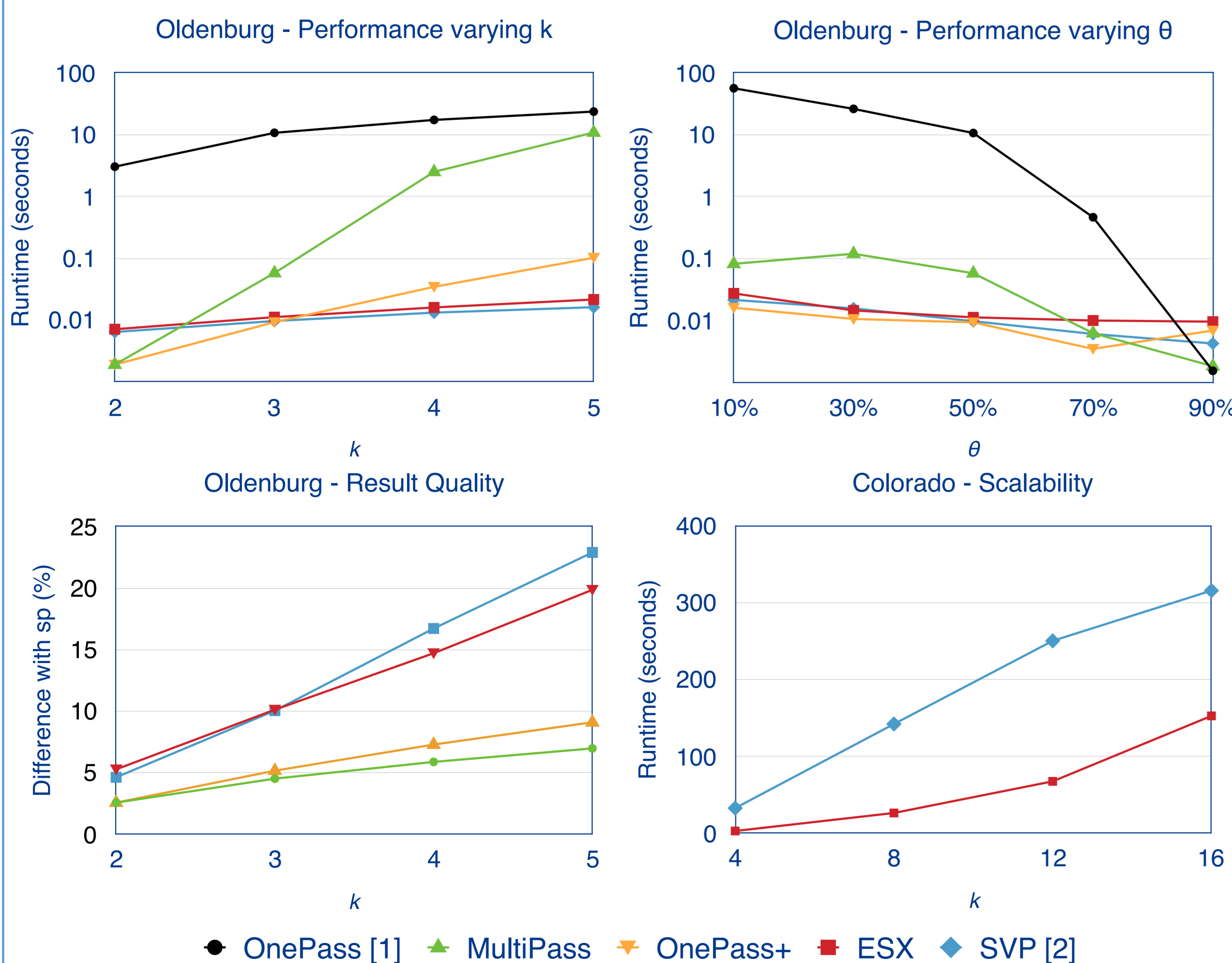
- ✓ Iteratively removes edges from the road network that lie on some already computed alternative path
- ✓ Computes the shortest path on the updated graph
- ✓ Continues until a sufficiently dissimilar path is found



Sample execution of ESX ( $\theta=50\%$ )

1. Remove edge  $(n_5, t)$   
 $p_1 = \langle (s, n_3)(n_3, n_5)(n_5, n_4)(n_4, t) \rangle$   
**Path  $p_1$  is not an alternative**
2. Remove edge  $(n_3, n_5)$   
 $p_2 = \langle (s, n_3)(n_3, n_5)(n_5, n_4)(n_4, t) \rangle$   
**Path  $p_2$  is an alternative**

## EXPERIMENTAL EVALUATION



## SUMMARY

**MultiPass:** computes the optimal result but is practical only for small road networks

**OnePass+:** good approximation and practical for larger road networks than MultiPass

**ESX:** less accurate but practical even for large road networks and large values of  $k$

[1] T. Chondrogiannis, P. Bouros, J. Gamper, U. Leser. Alternative routing:  $k$ -shortest paths with limited overlap. *In Proc. of the 23rd ACM SIGSPATIAL International Conference on Advances in Geographic Information Systems*, pages 68:1–68:4, 2015.

[2] I. Abraham, D. Delling, A.V. Goldberg, R.F. Werneck. Alternative routes in road networks. *Journal of Experimental Algorithmics*, 18:1–17, 2013.