An Introduction to Analysis and Verification of Software

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Goals of this talk

- To give a brief overview of the area
- Establish some terminology, including some "common knowledge"
- Propose study group projects
- Present some aspect by which the various projects may be evaluated and compared

Bugs

bug: a mismatch between **implementation** and **specification**

 If there is no spec, it's not a bug but a feature :-)



- Catching bugs after deployment can be <u>expensive</u>
- Detecting bugs vs. guaranteeing correctness?

Where does (semi-)**automated** analysis and verification have a chance?

- resource management (memory, files, allocation, locking)
- temporal properties (event ordering, concurrency, deadlocks, safety/liveness)
- datatype invariants (shapes, memory errors)
- security (integrity, confidentiality)
- numerical computations
- •

Typical bugs? See e.g. <u>bugzilla.mozilla.org</u> or <u>bugzilla.kernel.org</u>

"Everything interesting about the behavior of programs is undecidable." [paraphrase of H.G. Rice, 1953]

- Approximations (perferably conservative)
- Annotations (invariants)

Verification vs. Analysis

- Verification: checking invariants
- Analysis: detecting invariants

- this is just one possible definition of the words
- in practical tools, there is a large overlap
- analysis is not necessarily "harder" than verification

Relation to Debugging

- **Debugging**: you know there is a bug, but not exactly where it is
- Verification: you hope there are no bugs, but you want to be sure
- Analysis tools can be used to enhance program understanding (e.g. "program slicing")
- Verification tools can often provide counterexamples

"Program testing can be used to show the presence of bugs, but never to show their absence." [Dijkstra, 1972]

Relation to Program Optimization

- Many of the same program analysis techniques apply
- **Moore's law:** "the performance of microprocessors doubles every 18 months"
- **Proebsting's law:** "compiler technology doubles the performance of typical programs every 18 years"
- Conclusion: analysis/verification is more fun than optimization ⁽²⁾
- or: invent new high-level languages that need new optimization techniques!

Relation to Programming Language Design

- Programming at a higher level of abstraction can reduce the possibility of errors
- Examples:
 - type systems (note: types are invariants!)
 - abstract data-types
 - object oriented encapsulation and inheritance
 - domain-specific languages

A Spectrum of Techniques

- Light-weight
 - unsound and incomplete
 - bug searching via simulation
 - simple properties, efficient and easy to use
- Medium-weight
 - sound but incomplete
 - analysis via fixed-point computation, type checking/inference
- Heavy-weight
 - sound and complete
 - verification via theorem proving and user annotations
 - complex properties, resource demanding and difficult to use

Aspects

- domain of applicability (model heap, object state, events, numerical, parallelism, hardware/software, sequential/concurrent, reactive/transformational...)
- expressive power (fixed properties vs. using logic/automata in specs)
- degree of automation, annotations, theorem prover guidance
- modularity (interprocedural, whole-program, iteration, ...)
- scalability (program size)
- efficiency (time & space), theoretical complexity
- learnability (training requirements)
- soundness/completeness, spurious errors, heuristics, transparency (e.g. type systems should be transparent)
- error tracking (precision of error messages), counterexamples
- division between program abstraction and solution computation (verification conditions)

Approaches

- Axiomatic Semantics
 - Hoare logic, weakest precondition, separation logics
 - often used for sequential transformational programs
- Abstract Interpretation
 - data-flow analysis, constraint-based analysis
- Model Checking
 - state-space exploration with modal logics
 - often used for concurrent reactive systems
- Type Systems
 - type checking and inference

Hot Research Projects and Topics

- <u>SLAM</u> model checking for boolean programs and C, abstraction refinement
- Java PathFinder model checking for Java
- <u>Bandera</u> model checking for Java
- <u>Blast</u> model checking for C, abstraction refinement
- <u>SPIN</u> LTL model checking
- <u>ESP</u> control-flow + data-flow analysis on C/C++
- <u>3-Valued Logic Analysis Engine TVLA</u> shape analysis
- <u>BANE/Banshee</u> constraint-based program analysis
- <u>Abstract Interpretation</u> foundation for program analysis
- <u>Behave!</u> checking behavioral properties using type systems and Pi-calculus
- <u>Vault</u> type system for a safe version of C
- <u>Cyclone</u> type system for a safe version of C
- <u>CQual</u> type qualifiers for C
- <u>PVS</u> verification system based on theorem proving with higher-order logic
- <u>ESC/Java</u> theorem proving for light-weight properties
- <u>Proof-Carrying Code</u> safe execution of untrusted code
- <u>Meta-Level Compilation / Metal</u> finding bugs in system code, e.g. Linux kernel