

Computer Science Day: Mathematical Computer Science

Thomas Dueholm Hansen

June 15, 2011

The Mathematical Computer Science group

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- Kristoffer Arnsfelt Hansen
- Peter Bro Miltersen

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- Elad Verbin (also at MADALGO)
- Joshua Brody
- Dominik Scheder
- Wei Yu (also at MADALGO)
- Jie Zhang

- PhD students:

- Søren Stil Frederiksen
- Thomas Dueholm Hansen
- Rasmus Ibsen-Jensen

- Administration:

- Dorthe Haagen Nielsen



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Terminated within the last year:

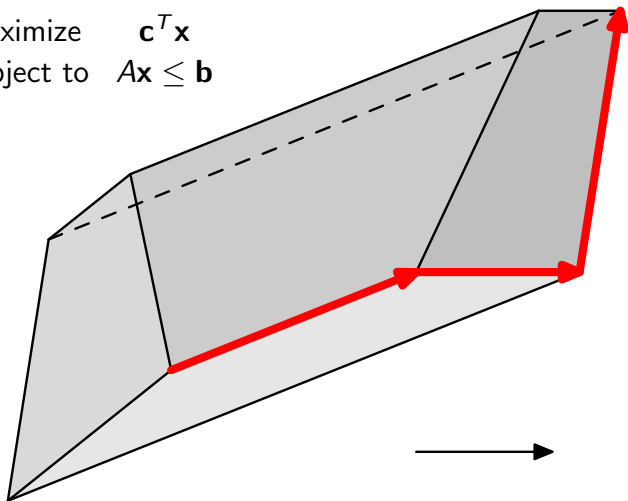
AGT Center for
Algorithmic Game Theory

COMA

- Computational Complexity Theory
- Algorithmic and Computational Game Theory
- Mathematical Programming
- Combinatorial Optimization
- Approximation Algorithms
- Computational Algebra
- Symbolic-Numerical Computations
- Real Algebraic Geometry
- Theoretical aspects of multiagent systems

Linear programming: Dantzig's simplex algorithm (1947)

maximize $\mathbf{c}^T \mathbf{x}$
subject to $A\mathbf{x} \leq \mathbf{b}$



Example: The diet problem

Buy food items in order to satisfy daily intake of energy (2000 kcal), protein (55 g) and calcium (800 mg), minimizing the cost.

Food	Serving size	Energy	Protein	Calcium	Price
Oatmeal	28 g	110	4	2	3
Chicken	100 g	205	32	12	24
Eggs	2 large	160	13	54	13
Whole milk	237 cc	160	8	285	9
Cherry pie	170 g	420	4	22	20
Pork with beans	260 g	260	14	80	19

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$$\text{Minimize} \quad 3x_1 + 24x_2 + 13x_3 + 9x_4 + 20x_5 + 19x_6$$

$$\text{subject to} \quad 110x_1 + 205x_2 + 160x_3 + 160x_4 + 420x_5 + 260x_6 \geq 2000$$

$$4x_1 + 32x_2 + 13x_3 + 8x_4 + 4x_5 + 14x_6 \geq 55$$

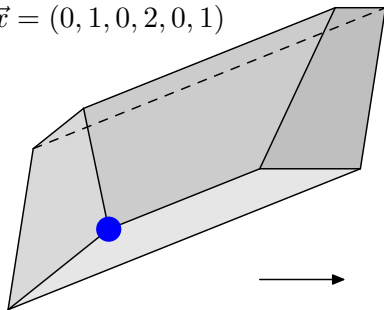
$$2x_1 + 12x_2 + 54x_3 + 285x_4 + 22x_5 + 80x_6 \geq 800$$

$$x_1, x_2, x_3, x_4, x_5, x_6 \geq 0$$

Basic feasible solutions and pivoting

$$\vec{x} = (0, 1, 0, 2, 0, 1)$$

$$\begin{array}{ll} \max & 2x_1 - 2x_3 - 2x_5 - x_6 \\ \text{s.t.} & \frac{1}{3}x_1 + x_2 - \frac{2}{3}x_3 - \frac{2}{3}x_5 = 1 \\ & x_3 + x_4 - x_6 = 1 \\ & x_5 + x_6 = 1 \\ & x_1, x_2, x_3, x_4, x_5, x_6 \geq 0 \end{array}$$

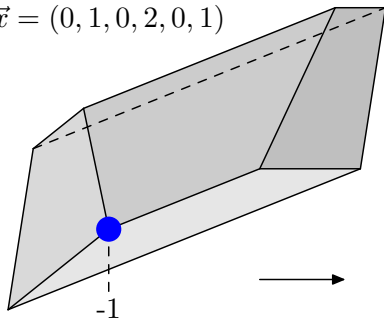


- The corners of the polytope correspond to **basic feasible solutions**: At most n , the number of equality constraints, variables are non-zero. A *basis* consists of n *basic* variables such that all non-basic variables are zero.

Basic feasible solutions and pivoting

$$\vec{x} = (0, 1, 0, 2, 0, 1)$$

$$\begin{aligned} \max \quad & -1 + 2x_1 - 2x_3 - x_5 \\ \text{s.t.} \quad & x_2 = 1 - \frac{1}{3}x_1 + \frac{2}{3}x_3 + \frac{2}{3}x_5 \\ & x_4 = 2 - x_3 - x_5 \\ & x_6 = 1 - x_5 \\ & x_1, x_2, x_3, x_4, x_5, x_6 \geq 0 \end{aligned}$$

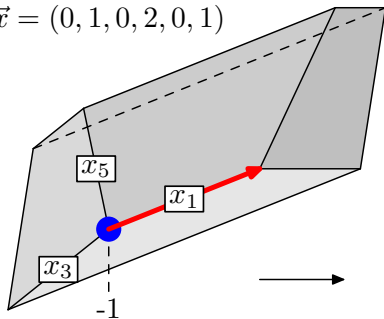


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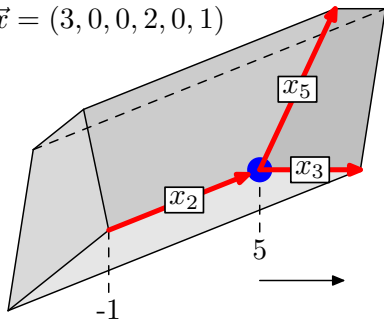


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Basic feasible solutions and pivoting

$$\begin{aligned} \max \quad & 5 - 6x_2 + 2x_3 + 3x_5 \\ \text{s.t.} \quad & x_1 = 3 - 3x_2 + 2x_3 + 2x_5 \\ & x_4 = 2 - x_3 - x_5 \\ & x_6 = 1 - x_5 \\ & x_1, x_2, x_3, x_4, x_5, x_6 \geq 0 \end{aligned}$$

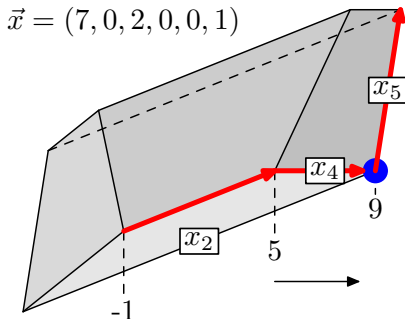
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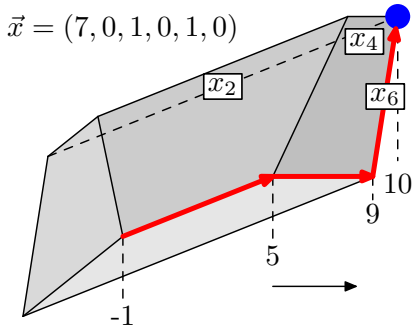
$$\begin{aligned} \max \quad & 9 - 6x_2 - 2x_4 + x_5 \\ \text{s.t.} \quad & x_1 = 7 - 3x_2 - 2x_4 \\ & x_3 = 2 - x_4 - x_5 \\ & x_6 = 1 - x_5 \\ & x_1, x_2, x_3, x_4, x_5, x_6 \geq 0 \end{aligned}$$



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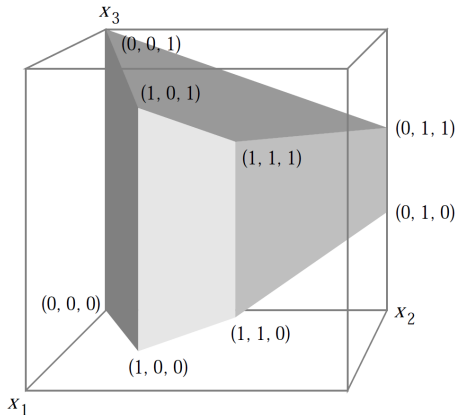
$$\begin{aligned} \max \quad & 10 - 6x_2 - 2x_4 - x_6 \\ \text{s.t.} \quad & x_1 = 7 - 3x_2 - 2x_4 \\ & x_3 = 1 - x_4 + x_6 \\ & x_5 = 1 - x_6 \\ & x_1, x_2, x_3, x_4, x_5, x_6 \geq 0 \end{aligned}$$



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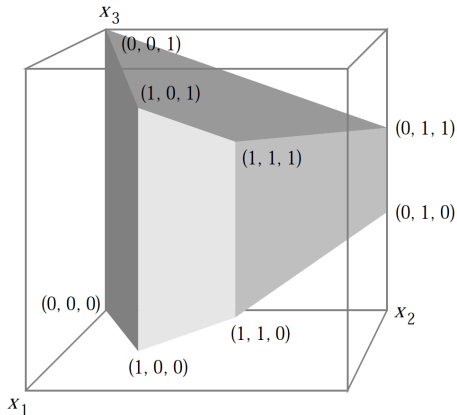
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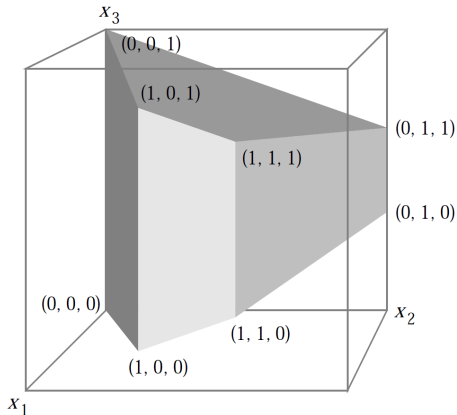
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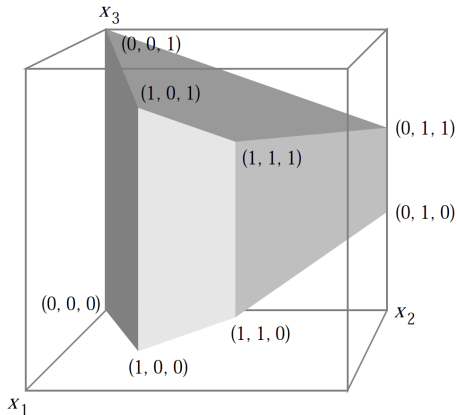
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- What about randomized
pivoting rules?
 - No good lower bounds
were known for forty
years.



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Theorem (Friedmann, Hansen and Zwick (2011))

RANDOMEDGE and RANDOMFACET may require an expected number of pivoting steps that is almost exponential in the number of constraints and variables: $2^{\Omega(n^{1/4})}$ and $2^{\Omega(\sqrt{n}/\log n)}$, respectively.

Thank you for listening!